



UNICAMP



**Quantum Contextuality on Quantum Measurements
and Beyond (QCQMB) 2021**

Emergence of noncontextuality and classical limits

Roberto Dobal Baldijão*

With: Rafael Wagner, Cristhiano Duarte, Bárbara Amaral, Marcelo Terra Cunha

Arxiv 2104.05734

With: Marcelo Terra Cunha

Phys. Rev A 102, 052226(2020)/Arxiv 1811.00615

*rdbaldi@ifi.unicamp.br

The problem

Obs 1

Contextuality is a **nonclassical** feature

The problem

Obs 1

Contextuality is a **nonclassical** feature

Obs 2

Our everyday experience is **classical**

The problem

Obs 1

Contextuality is a **nonclassical** feature

Obs 2

Our everyday experience is **classical**

Obs 3

Quantum Theory is contextual

The problem

Contextuality is a **nonclassical** feature

Our everyday experience is **classical**

Quantum theory is contextual

The problem

Contextuality is a **nonclassical** feature

Our everyday experience is **classical**

Quantum theory is contextual



Classical limits must kill quantum contextuality

The problem

Contextuality is a **nonclassical** feature

Our everyday experience is **classical**

Quantum theory is contextual

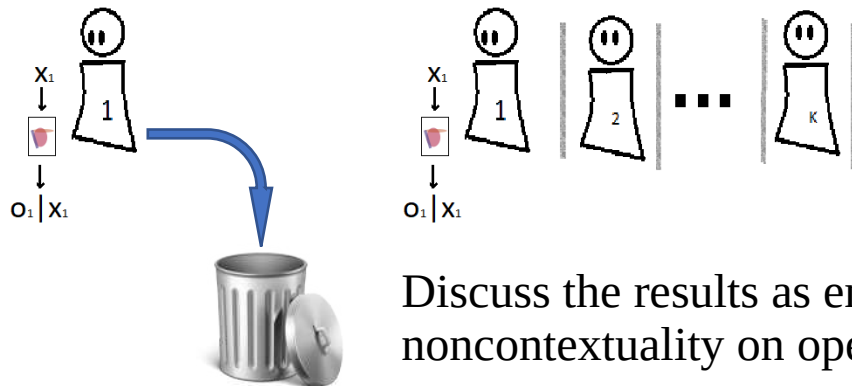


Classical limits must kill
quantum contextuality

Does contextuality fade out under
classical limits?

Part I

We analyze the best quantum realization for the odd N-cycle inequalities in **multiple** observers setup

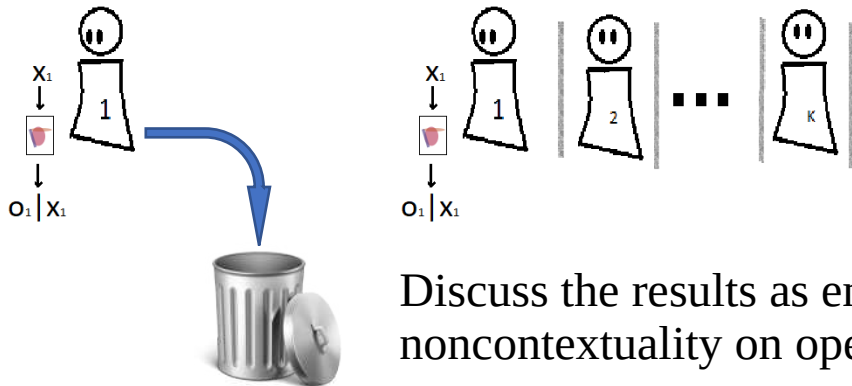


Discuss the results as emergence of noncontextuality on open quantum systems

Outline

Part I

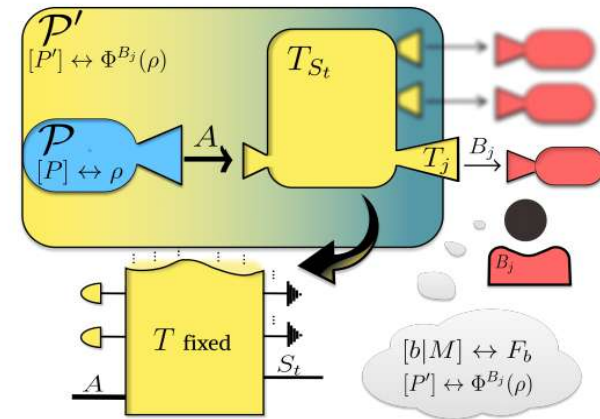
We analyze the best quantum realization for the odd N-cycle inequalities in **multiple** observers setup



Discuss the results as emergence of noncontextuality on open quantum systems

Part II

Noncontextuality emerges under quantum Darwinism processes



Part I

Part I requests

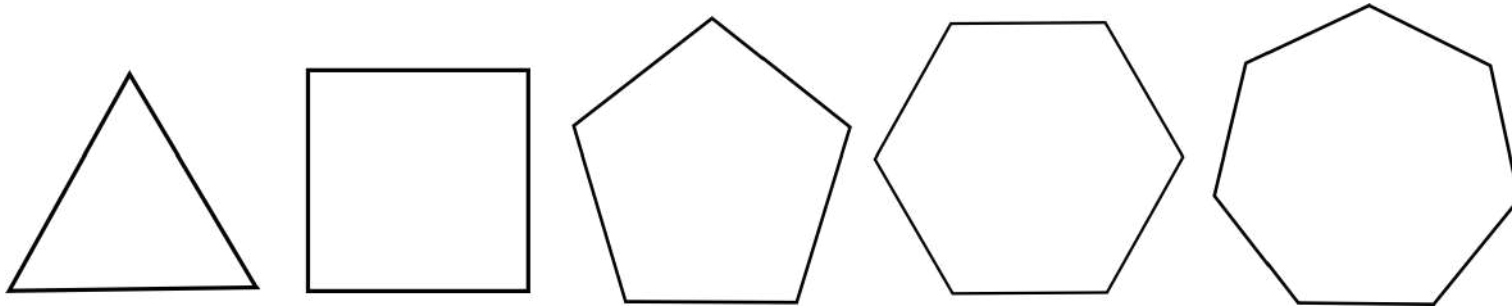
We need to set our ground:

N-cycle inequalities

Contextuality: N-cycle scenarios

- Simple compatibility scenarios.

$$\Gamma \equiv \{X, \mathcal{C}, O\}$$

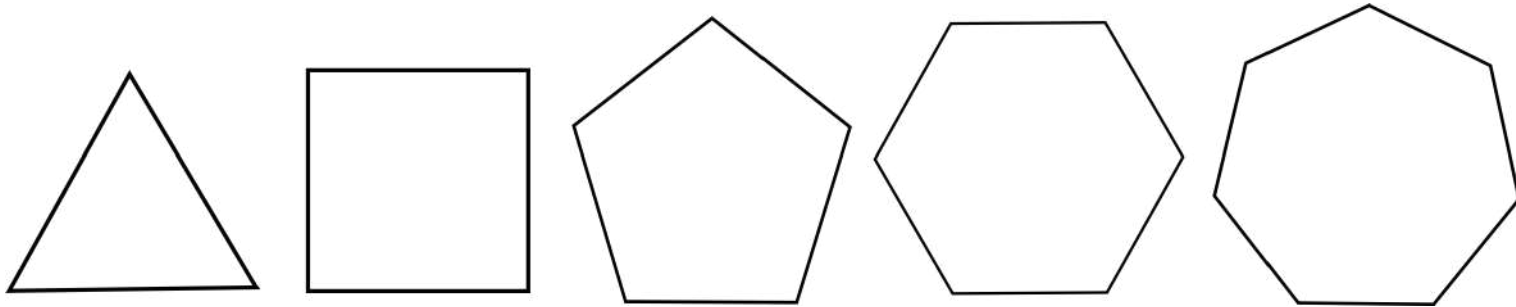


$$\mathcal{C} = \{\{A_0, A_1\}, \{A_1, A_2\}, \dots, \{A_{N-1}, A_0\}\}$$

Contextuality: N-cycle scenarios

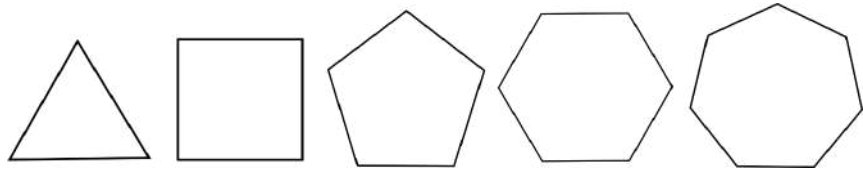
- Simple compatibility scenarios.

$$\Gamma \equiv \{X, \mathcal{C}, O\}$$



Scenarios with odd N cannot be split into 2 parties!

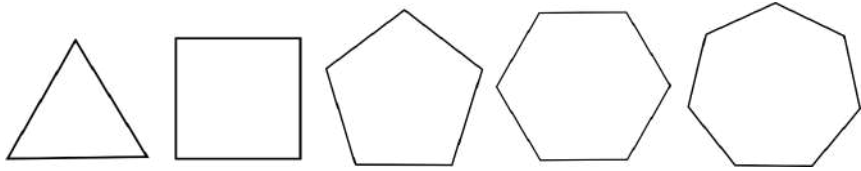
Contextuality: N-cycle scenarios



Odd N $o_i \in \{-1, 1\}$

$(+1, +1)$ $(+1, -1)$ $(-1, +1)$ $(-1, -1)$

Contextuality: N-cycle scenarios



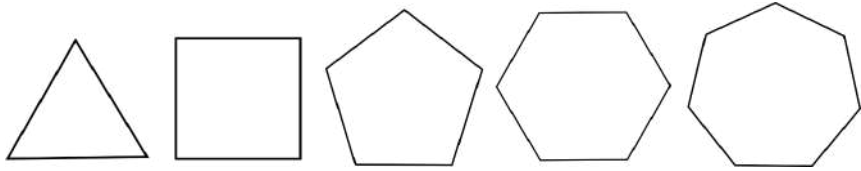
Odd N $o_i \in \{-1, 1\}$

$(+1, +1)$ $(+1, -1)$ $(-1, +1)$ $(-1, -1)$

Inequality

$$\sum_{i=0}^{N-1} \langle A_i A_{i+1} \rangle \geq -N + 2$$

Contextuality: N-cycle scenarios



Odd N $o_i \in \{-1, 1\}$

$(+1, +1)$ $(+1, -1)$ $(-1, +1)$ ~~$(-1, -1)$~~

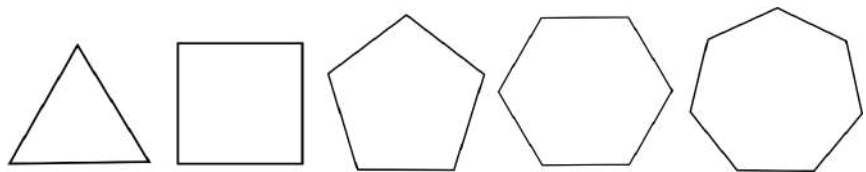
$$\sum_{i=0}^{N-1} \langle A_i A_{i+1} \rangle \geq -N + 2$$

Extra assumption

Exclusiveness:

$$p(-1, -1 | i, i \pm 1) = 0$$

Contextuality: N-cycle scenarios



Odd N $o_i \in \{-1, 1\}$

$(+1, +1)$ $(+1, -1)$ $(-1, +1)$ ~~$(-1, -1)$~~

$$\sum_{i=0}^{N-1} \langle A_i A_{i+1} \rangle \geq -N + 2$$

Exclusiveness

$$\sum_i p(1, 1|i, i+1) \geq 1$$

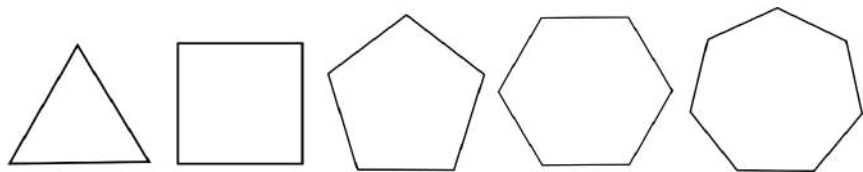
$$\sum_i p(-1|i) \leq \frac{N-1}{2}$$

Extra assumption

Exclusiveness:

$$p(-1, -1|i, i \pm 1) = 0$$

Contextuality: N-cycle scenarios



Odd N $o_i \in \{-1, 1\}$

$$\sum_{i=0}^{N-1} \langle A_i A_{i+1} \rangle \geq -N + 2$$

Exclusiveness

$$\sum_i p(1, 1 | i, i+1) \geq 1$$

$$\sum_i p(-1 | i) \leq \frac{N-1}{2}$$

Extra assumption

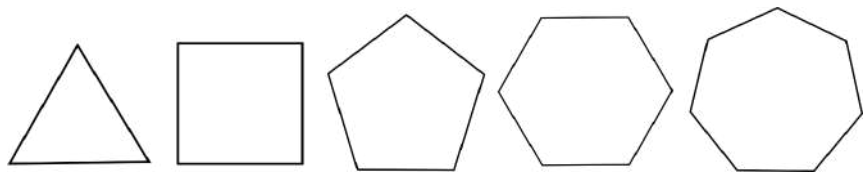
Exclusiveness:

$$p(-1, -1 | i, i \pm 1) = 0$$

b_i

a_i

Contextuality: N-cycle scenarios



Odd N $o_i \in \{-1, 1\}$

$$\sum_{i=0}^{N-1} \langle A_i A_{i+1} \rangle \geq -N + 2$$

Exclusiveness

$$\sum_i p(b_i) \geq 1$$

$$\sum_i p(a_i) \leq \frac{N-1}{2}$$

Extra assumption

Exclusiveness:

$$p(-1, -1 | i, i \pm 1) = 0$$

Contextuality: N-cycle scenarios

$$\sum_i p(b_i) \geq 1$$

$$\sum_i p(a_i) \leq \frac{N-1}{2}$$

Measurement protocols

1. $\mathcal{M}_i = \{a_i, b_i, a_{i+1}\}_i$

2. $\mathcal{M}_i^a = \{a_i, \neg a_i\}_i$

3. $\mathcal{M}_i^b = \{b_i, \neg b_i\}_i,$

Contextuality: N-cycle scenarios

$$\sum_i p(b_i) \geq 1$$

$$\sum_i p(a_i) \leq \frac{N-1}{2}$$

Measurement protocols

1. $\mathcal{M}_i = \{a_i, b_i, a_{i+1}\}_i$

2. $\mathcal{M}_i^a = \{a_i, \neg a_i\}_i$

3. $\mathcal{M}_i^b = \{b_i, \neg b_i\}_i,$

Usually
equivalent!

Quantum realizations

$$\sum_i p(b_i) \geq 1$$

$$\sum_i p(a_i) \leq \frac{N-1}{2}$$

- Hilbert space of dimension 3

$$A_i = \mathcal{I} - 2|a_i\rangle\langle a_i|$$

Quantum realizations

$$\sum_i p(b_i) \geq 1$$

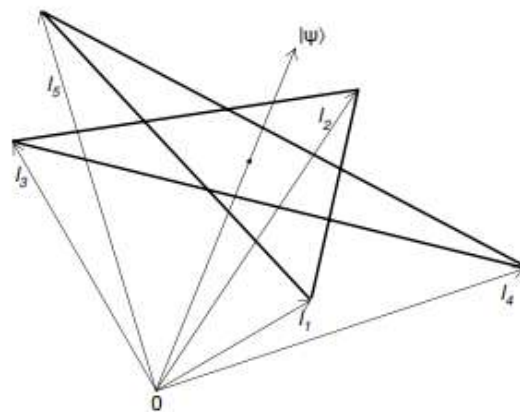
$$\sum_i p(a_i) \leq \frac{N-1}{2}$$

- Hilbert space of dimension 3

$$A_i = \mathcal{I} - 2|a_i\rangle\langle a_i|$$

$$\langle a_i | a_{i+1} \rangle = 0$$

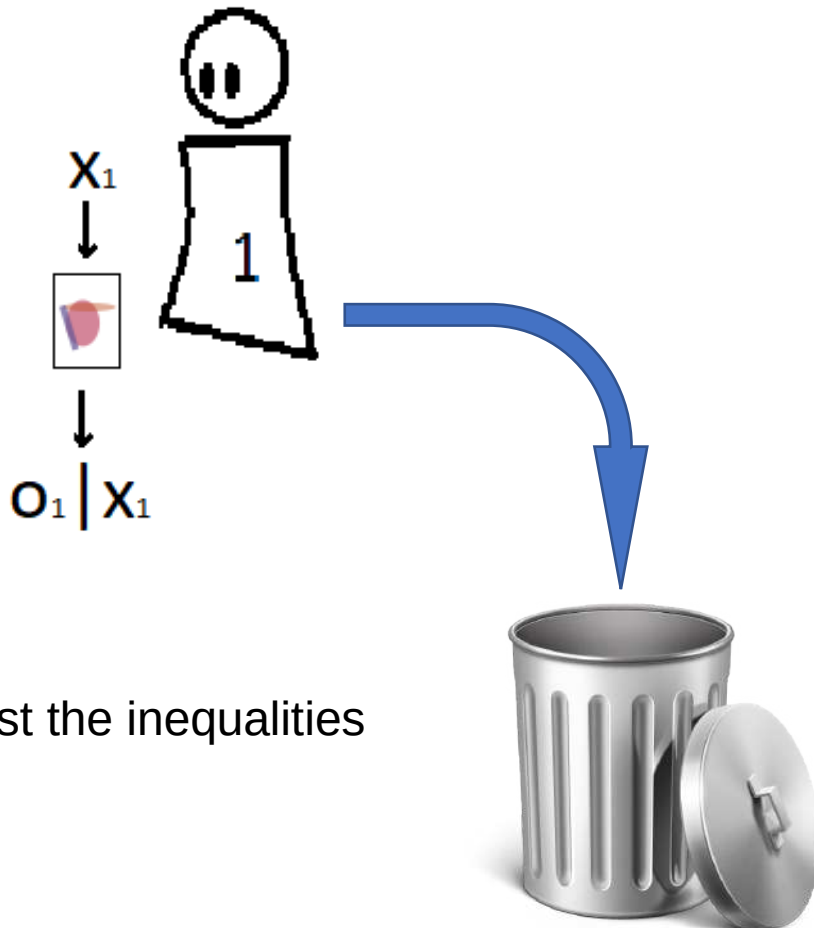
Exclusivity ✓



Contextuality: N-cycle scenarios

$$\sum_i p(b_i) \geq 1$$

$$\sum_i p(a_i) \leq \frac{N-1}{2}$$

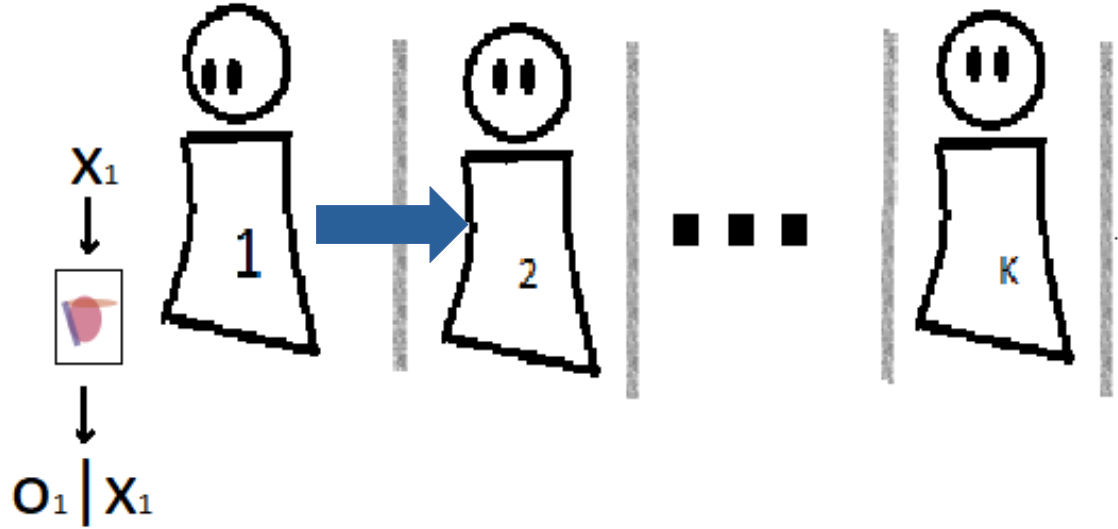


Usual way to test the inequalities

Sequential observers

$$\sum_i p(b_i) \geq 1$$

$$\sum_i p(a_i) \leq \frac{N-1}{2}$$

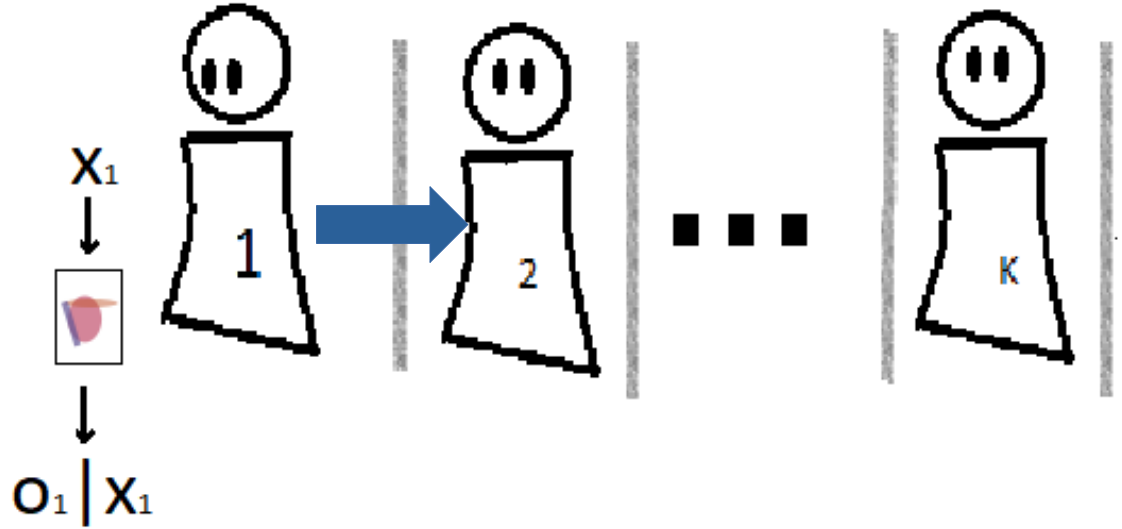


Our multiobserver setup proposal

Sequential observers

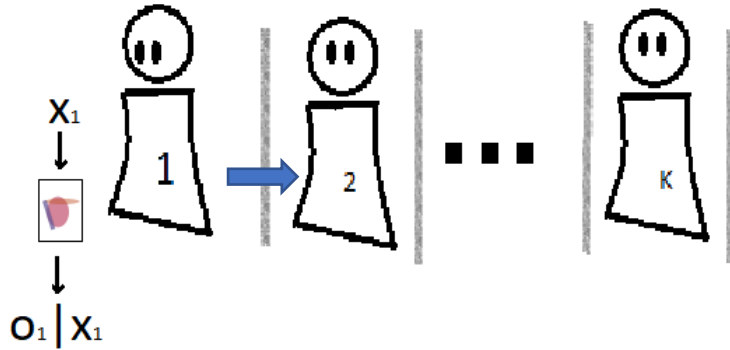
$$\sum_i p(b_i) \geq 1$$

$$\sum_i p(a_i) \leq \frac{N-1}{2}$$



Best quantum realization

Sequential observers



RDB, Marcelo Terra Cunha,
**Phys. Rev A 102, 052226(2020)/Arxiv
1811.00615**

For the odd N -cycle

- Which observers still find violations?
 - Does this depend on N ?
 - Does this depend on the measurement protocol?

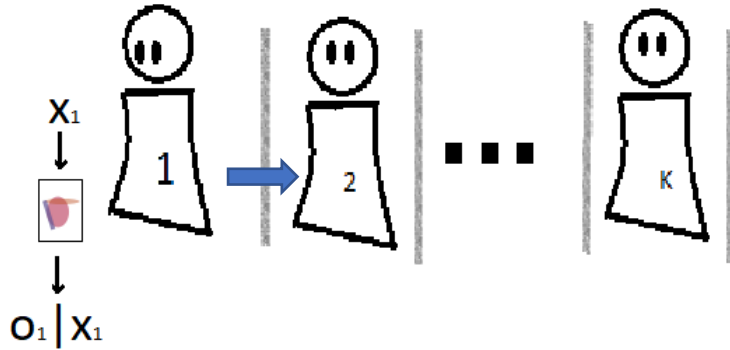
- Already analyzed, in different setups, for non-locality and steering, but lacking for Contextuality.

S. Sasmal et al, Phys. Rev. A98, 012305 (2018)

R. Silva et al, Phys. Rev. Lett.114, 250401 (2015)

D. Das et al, Phys. Rev. A99, 022305 (2019).

Sequential observers



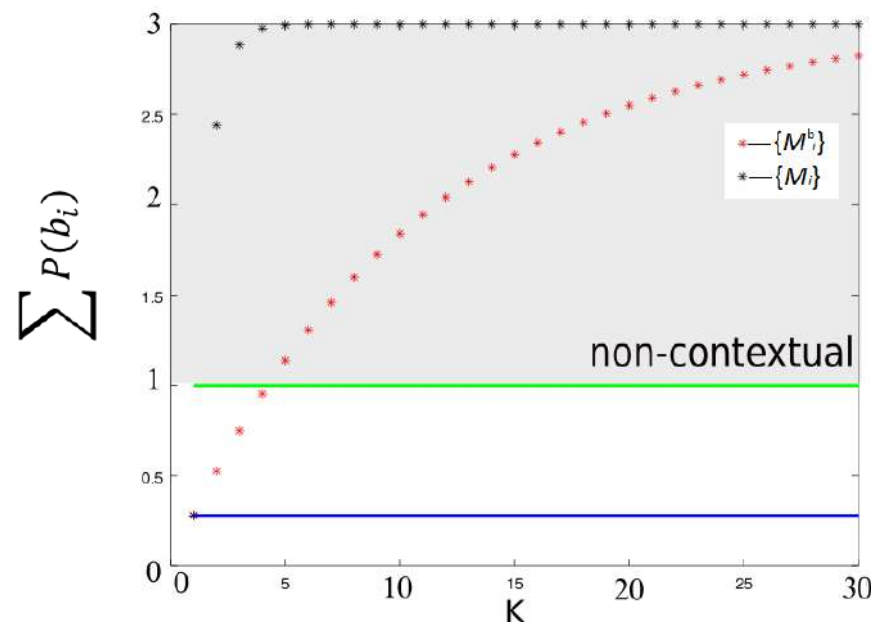
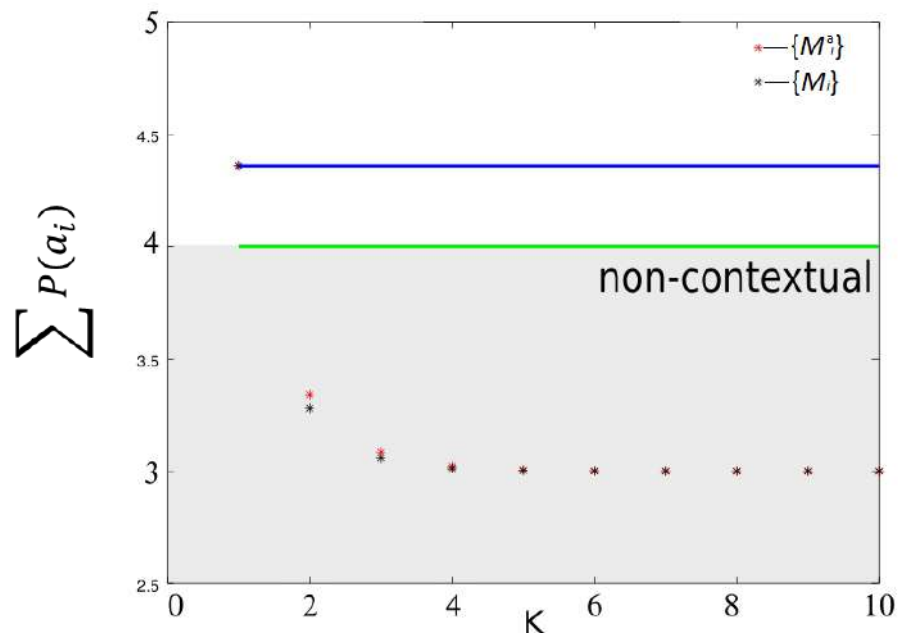
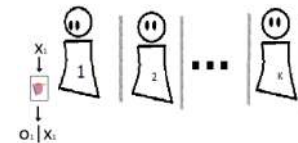
Here: we focus on the message we take regarding emergence of noncontextuality in open quantum systems

RDB, Marcelo Terra Cunha,
**Phys. Rev A 102, 052226(2020)/Arxiv
1811.00615**

For the odd N-cycle

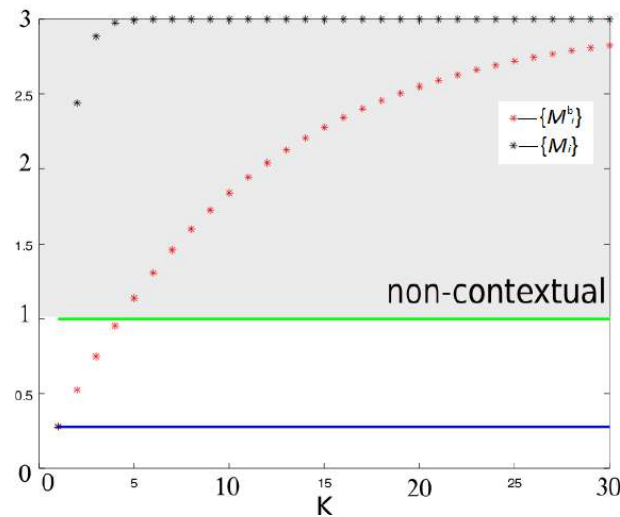
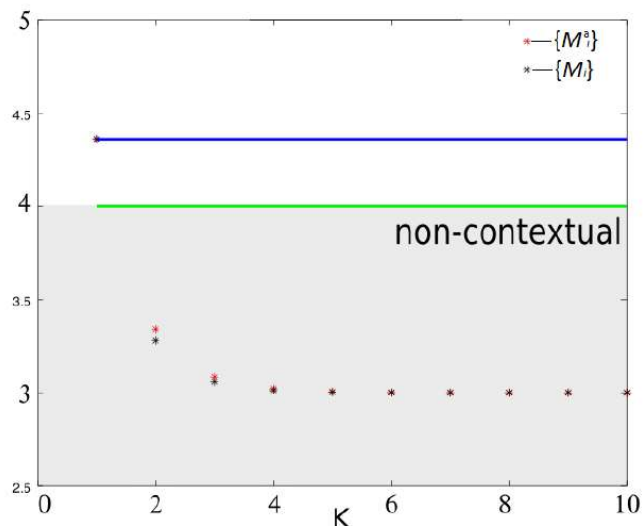
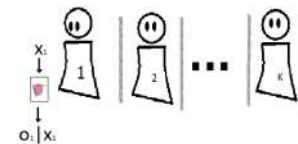
Results

The typical results (N=9)



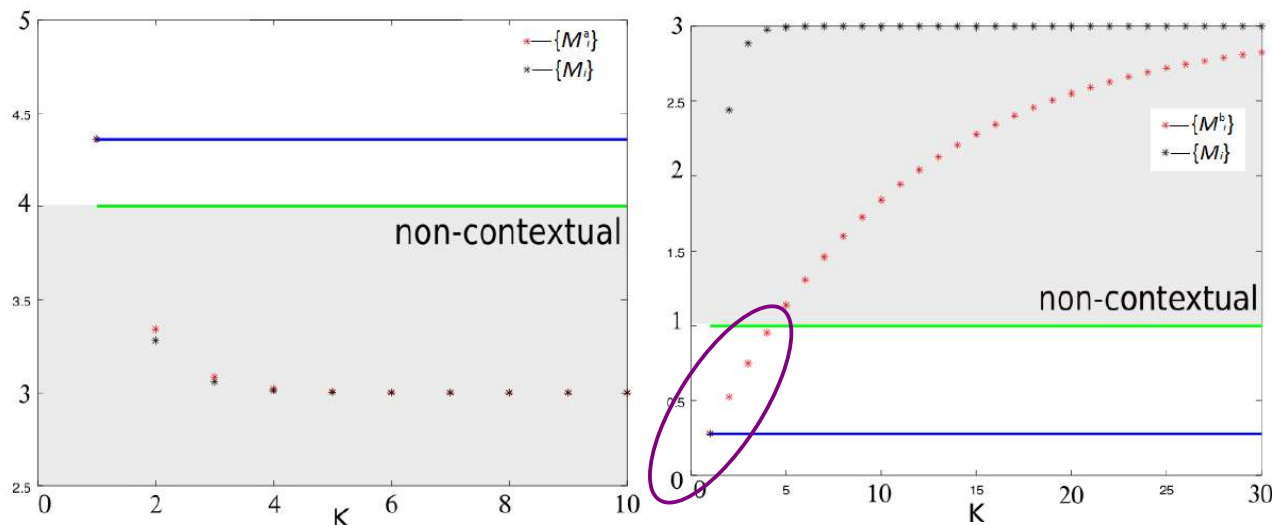
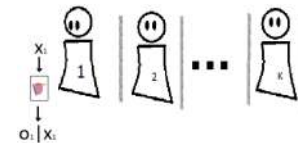
Results

The typical results (N=9)



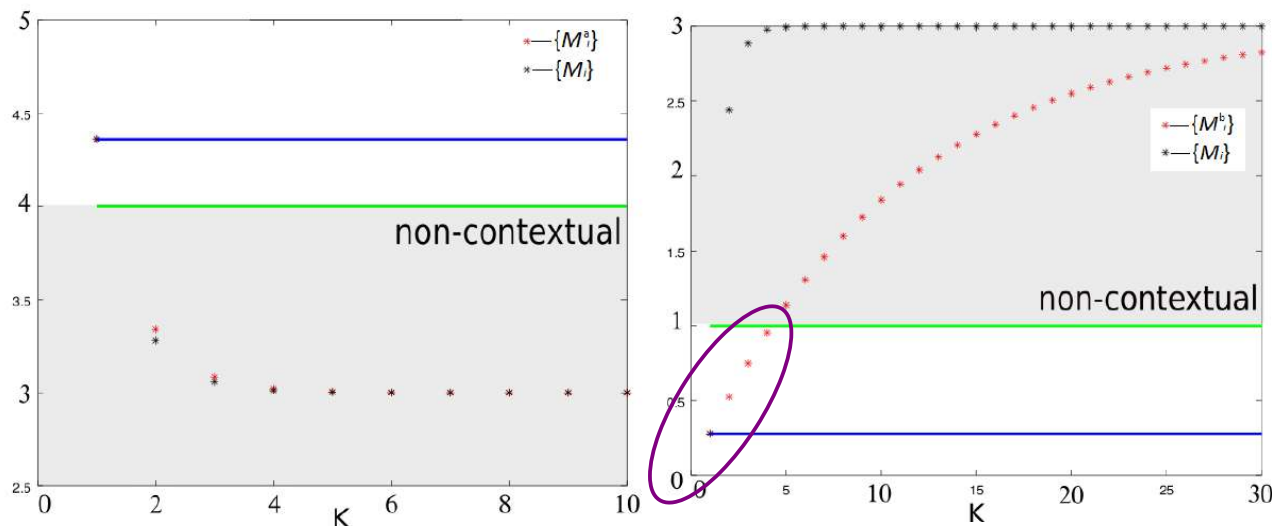
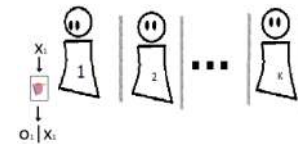
- Asymptotic limit without violation!
- Violation disappears quickly
- Protocols behave differently

Results



- Asymptotic limit without violation!
- Violation disappears quickly
- Protocols behave differently

Results



- Asymptotic limit without violation!
- Violation disappears quickly
- Protocols behave differently

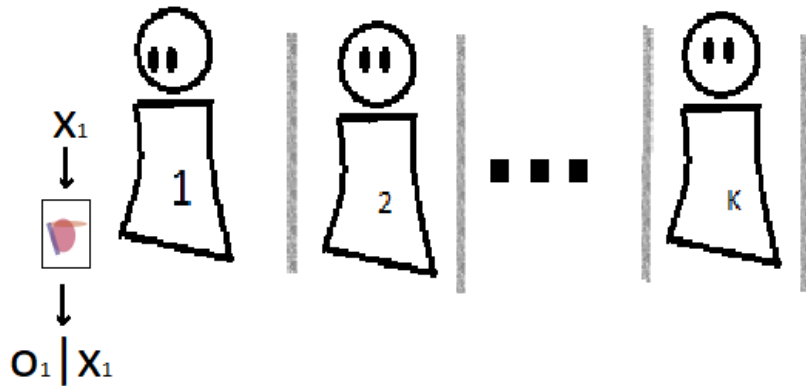
Results

- The behaviour showed for $N=9$ is qualitatively the same for all N .
- No-violation is always attained for large enough K , tending to the asymptotic limit.

Interpreting as a classical limit

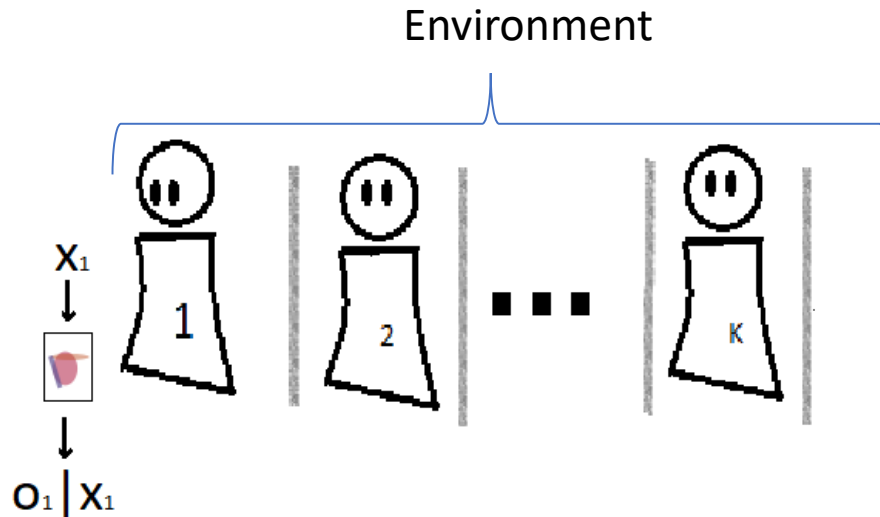
Interpreting as a classical limit

- Multiplayers setup as collisional models and multisystem environment



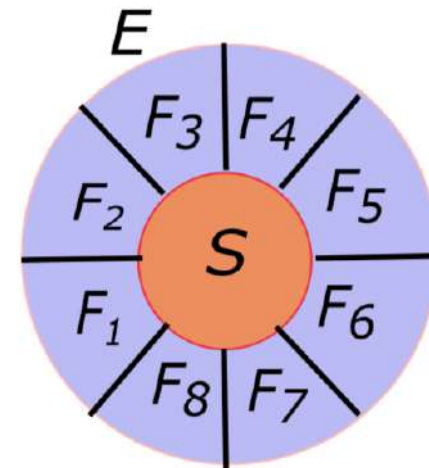
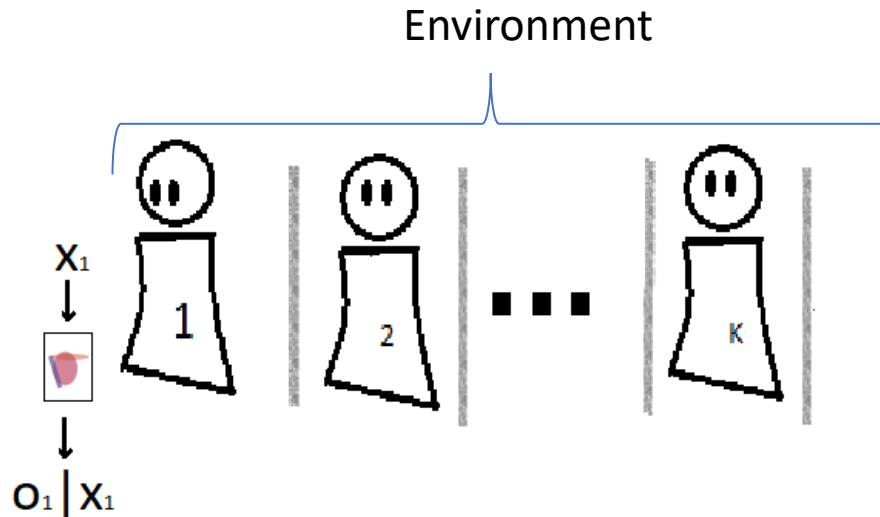
Interpreting as a classical limit

- Multiplayers setup as collisional models and multisystem environment



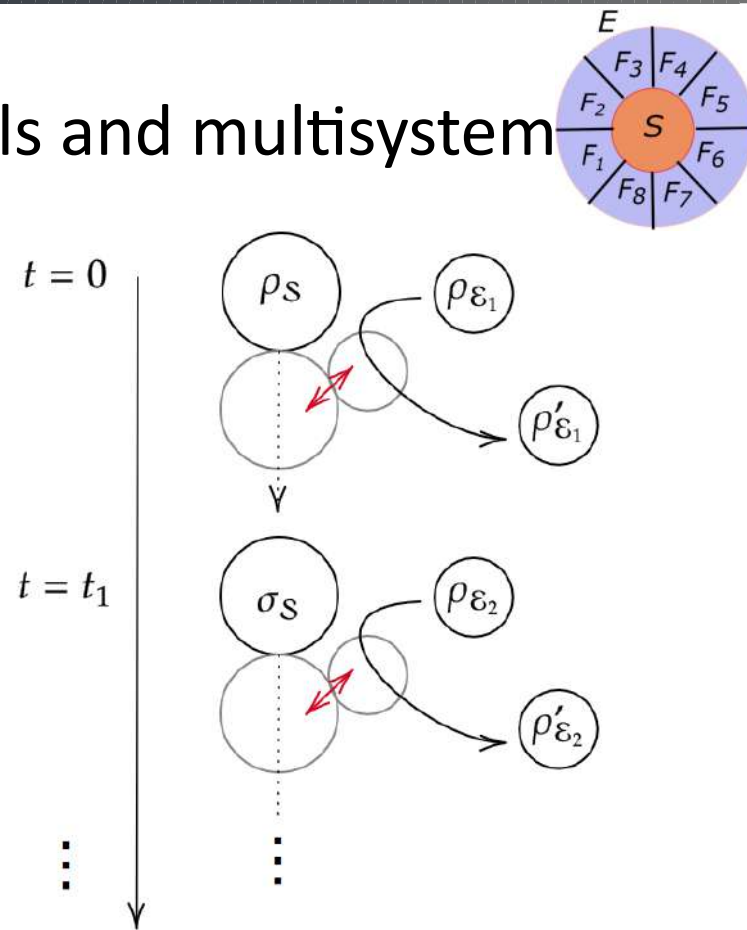
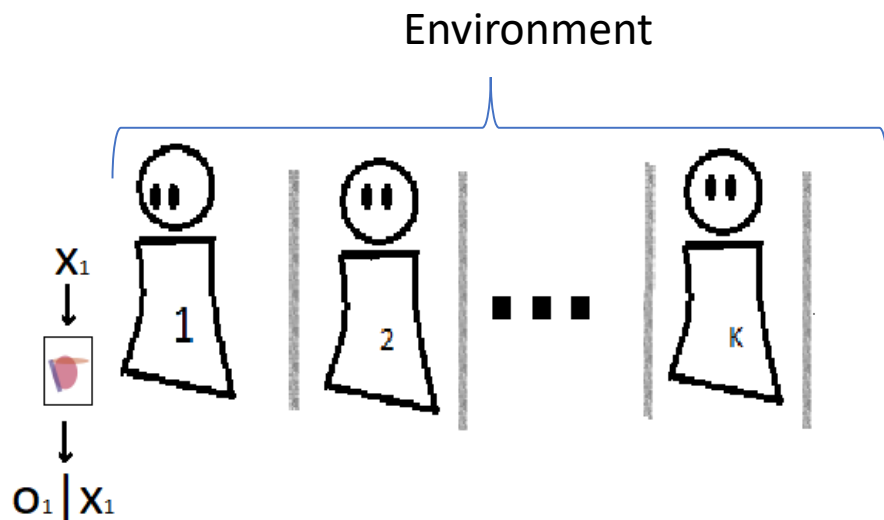
Interpreting as a classical limit

- Multiplayers setup as collisional models and multisystem environment



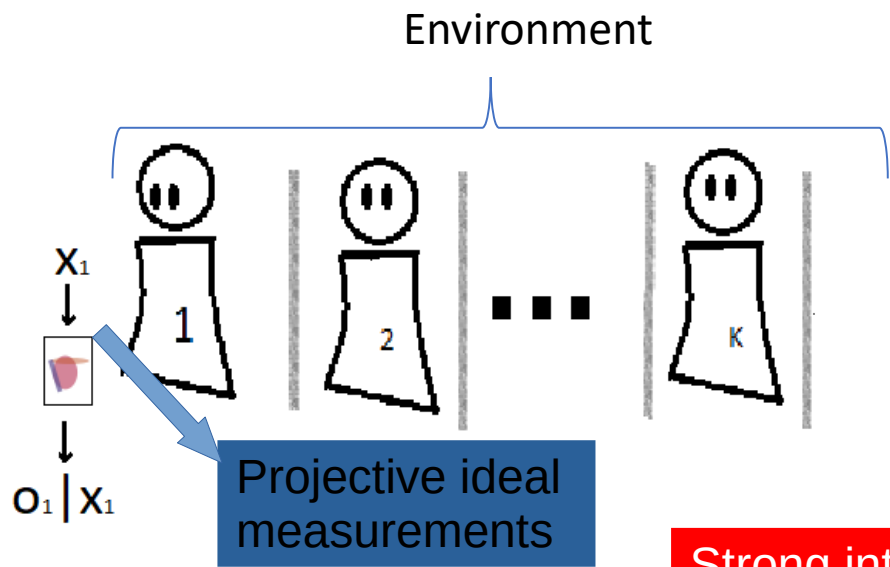
Interpreting as a classical limit

- Multiplayers setup as collisional models and multisystem environment

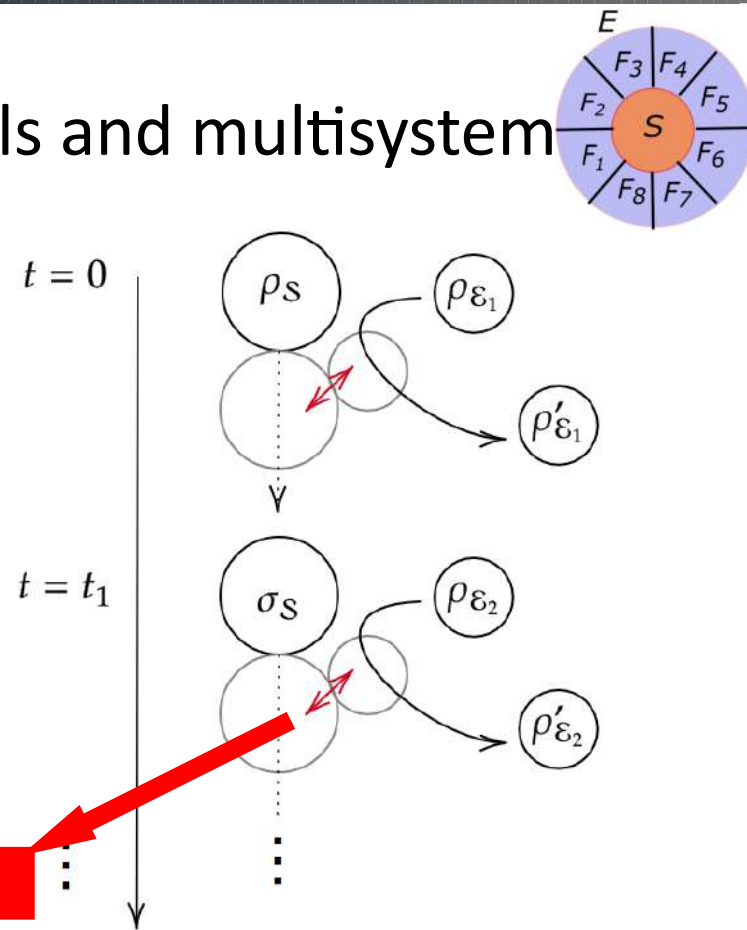


Interpreting as a classical limit

- Multiplayers setup as collisional models and multisystem environment



Strong interactions



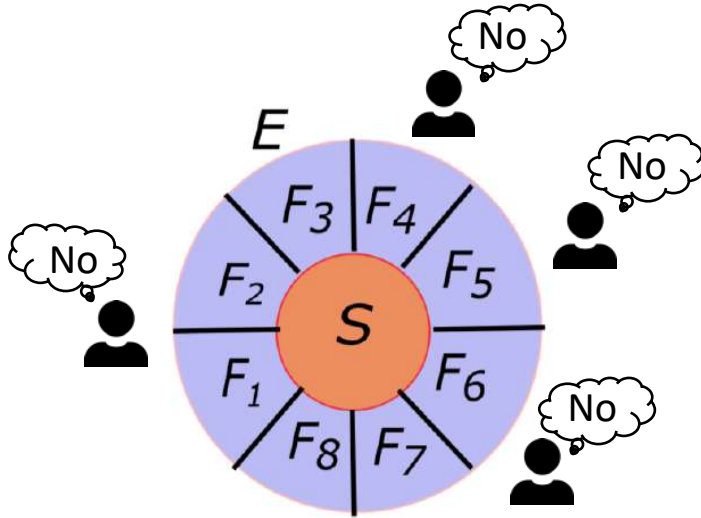
Interpreting as a classical limit

Our results
tell us...

Interpreting as a classical limit

Our results tell us...

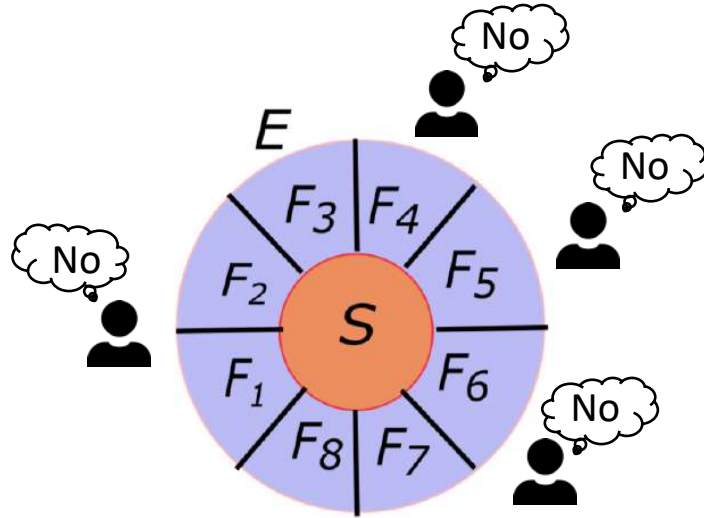
Can you witness contextuality?



Interpreting as a classical limit

Our results
tell us...

Can you witness
contextuality?



Dynamics makes contextuality disappear, for all N and large enough 'environment'.

Interpreting as a classical limit

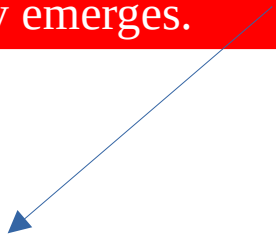
Take home message:

In this very idealized and special environment, designed to test contextuality, noncontextuality emerges.

Interpreting as a classical limit

Take home message:

In this very idealized and **special environment**, designed to test contextuality, noncontextuality emerges.



Does noncontextuality emerge in more generic classical limit processes?

Interpreting as a classical limit

Take home message:

In this very idealized and **special environment**, designed to test contextuality, noncontextuality emerges.

Suggestion:

Environment made of several (independent) subsystems

We look into information about a central system stored in these 'subenvironments'

Interpreting as a classical limit

Take home message:

In this very idealized and **special environment**, designed to test contextuality, noncontextuality emerges.

Suggestion: 

Environment made of several (independent) subsystems

We look into information about a central system stored in these 'subenvironments'

Quantum Darwinism!

Part II

Part II requests

We need to set our ground:

Quantum Darwinism (Brandão, Piani and Horodecki)

F. Brandão, M. Piani, P. Horodecki, Nat. Comm. 6 7908(2015)

Quantum Darwinism – BPH's approach

Quantum Darwinism – BPH's approach

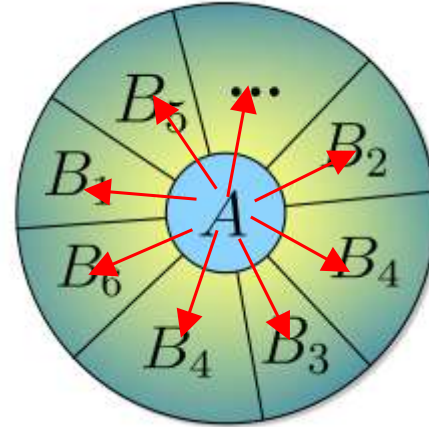


Quantum Darwinism – BPH's approach



A

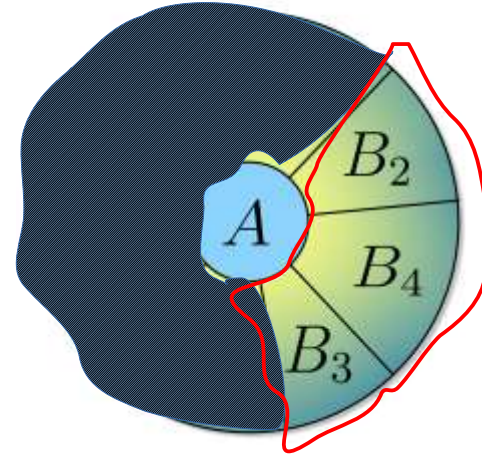
Quantum Darwinism – BPH's approach



Quantum Darwinism – BPH's approach



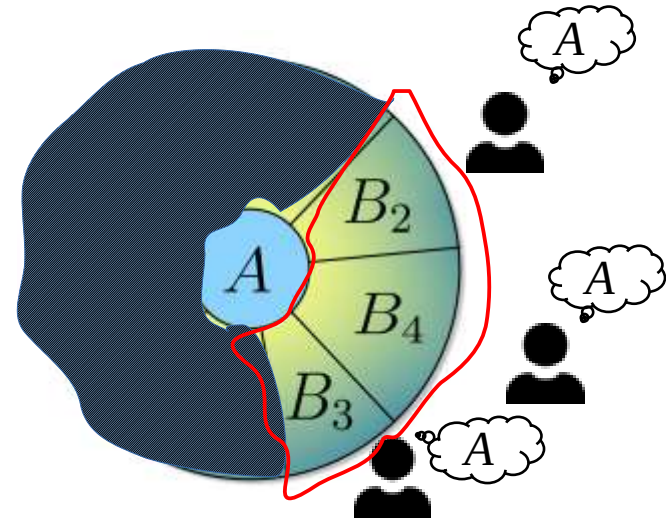
Environment as a witness dynamics



Quantum Darwinism – BPH's approach



Environment as a witness dynamics



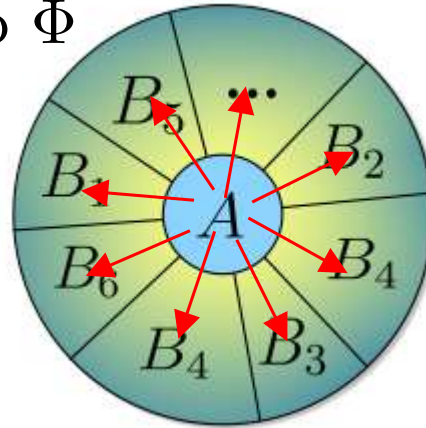
QD aims to explain why they see the same thing

Quantum Darwinism – BPH's approach

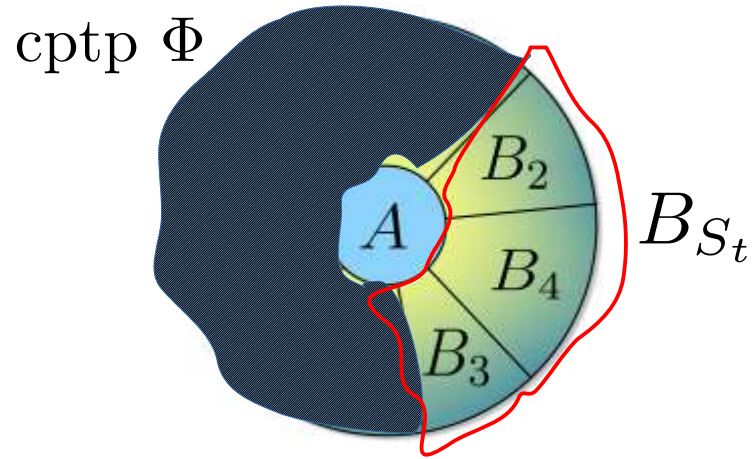
Let us formalize a bit...

Quantum Darwinism – BPH's approach

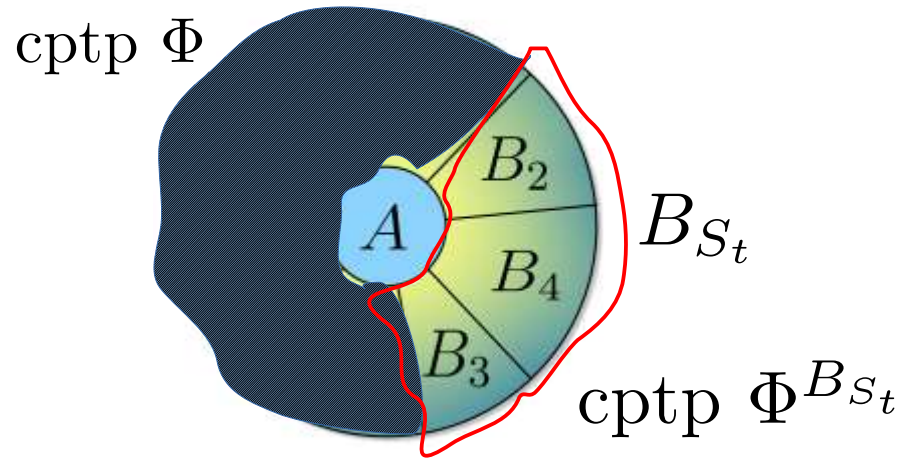
cptp Φ



Quantum Darwinism – BPH's approach

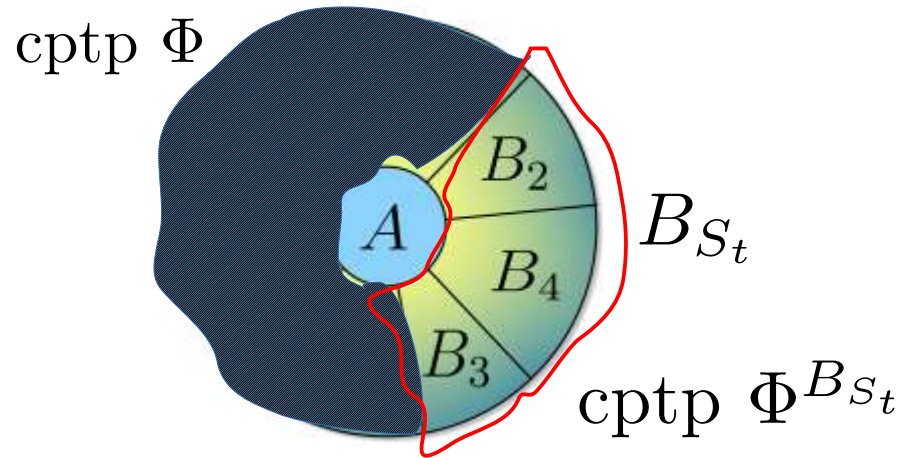


Quantum Darwinism – BPH's approach



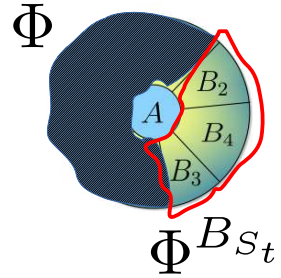
Quantum Darwinism – BPH's approach

Def. EWt-dynamics for the subset $B_{S_t} = (B_j)_{j \in S_t}$: $\Phi^{B_{S_t}} := \text{Tr}_{\mathcal{E} \setminus B_{S_t}} \circ \Phi$



Quantum Darwinism – BPH's approach

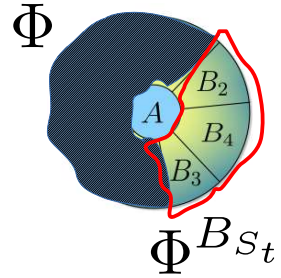
Def. EWt-dynamics for the subset $B_{S_t} = (B_j)_{j \in S_t}$: $\Phi^{B_{S_t}} := \text{Tr}_{\mathcal{E} \setminus B_{S_t}} \circ \Phi$



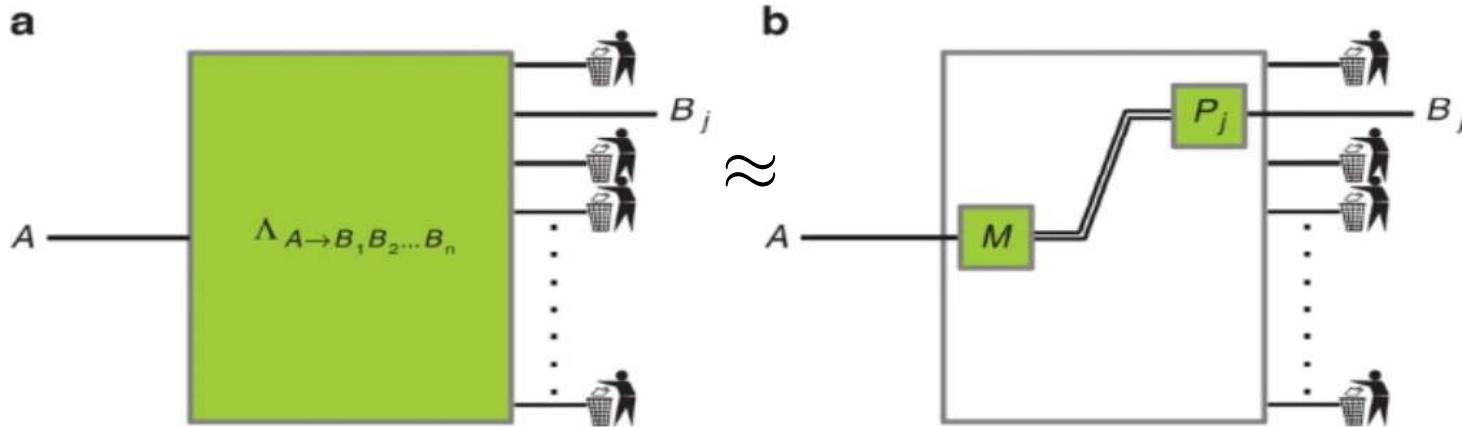
Theorem (Thm 2 of ref*)

Quantum Darwinism – BPH's approach

Def. EWt-dynamics for the subset $B_{S_t} = (B_j)_{j \in S_t}$: $\Phi^{B_{S_t}} := \text{Tr}_{\mathcal{E} \setminus B_{S_t}} \circ \Phi$

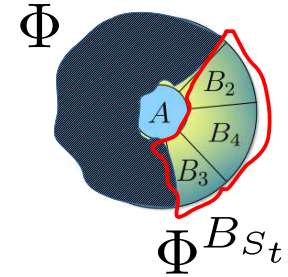


Theorem (Thm 2 of ref*)



*F. Brandão, M. Piani, P. Horodecki, Nat. Comm. 6 7908(2015)

Quantum Darwinism – BPH's approach



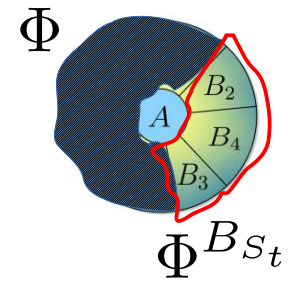
Def. EWt-dynamics for the subset $B_{S_t} = (B_j)_{j \in S_t}$: $\Phi^{B_{S_t}} := \text{Tr}_{\mathcal{E} \setminus B_{S_t}} \circ \Phi$

Theorem (Thm 2 of ref*)

Consider portion of size t . If $N \gg t$, then for most choices of B_{S_t} there exists $\{\tilde{E}_k\}$ and $\{\sigma_k^{B_{S_t}}\}$ s.t.

$$\Phi^{B_{S_t}}(\rho^A) \approx \sum_k \text{Tr}\{\tilde{E}_k \rho^A\} \sigma_k^{B_{S_t}}$$

Quantum Darwinism – BPH's approach



Def. EWt-dynamics for the subset $B_{S_t} = (B_j)_{j \in S_t}$: $\Phi^{B_{S_t}} := \text{Tr}_{\mathcal{E} \setminus B_{S_t}} \circ \Phi$

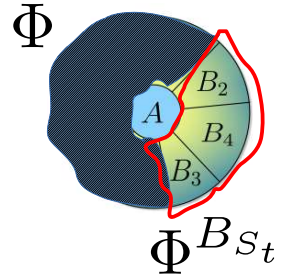
Theorem (Thm 2 of ref*)

Consider portion of size t . If $N \gg t$, then for most choices of B_{S_t} there exists $\{\tilde{E}_k\}$ and $\{\sigma_k^{B_{S_t}}\}$ s.t.

$$\Phi^{B_{S_t}}(\rho^A) \approx \sum_k \text{Tr}\{\tilde{E}_k \rho^A\} \sigma_k^{B_{S_t}}$$

*F. Brandão, M. Piani, P. Horodecki, Nat. Comm. 6 7908(2015)

Quantum Darwinism – BPH's approach



Def. EWt-dynamics for the subset $B_{S_t} = (B_j)_{j \in S_t}$: $\Phi^{B_{S_t}} := \text{Tr}_{\mathcal{E} \setminus B_{S_t}} \circ \Phi$

Theorem (Thm 2 of ref*)

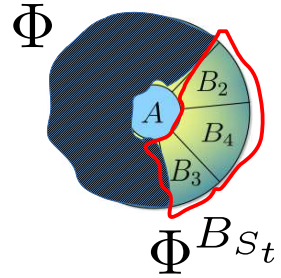
Consider portion of size t . If $N \gg t$, then for most choices of B_{S_t} there exists $\{\tilde{E}_k\}$ and $\{\sigma_k^{B_{S_t}}\}$ s.t.

$$\Phi^{B_{S_t}}(\rho^A) \approx \sum_k \text{Tr}\{\tilde{E}_k \rho^A\} \sigma_k^{B_{S_t}}$$

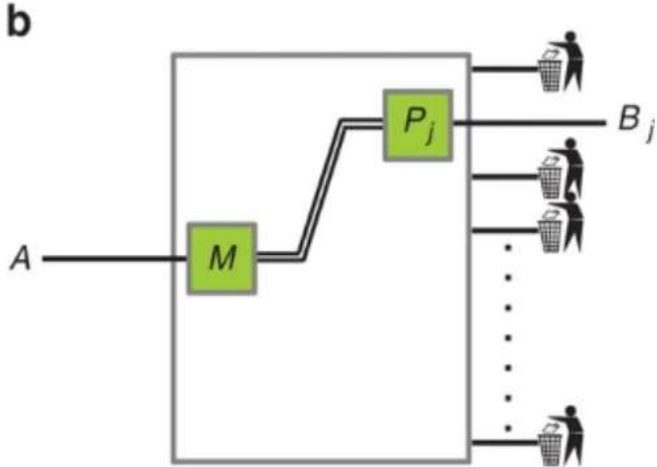
*F. Brandão, M. Piani, P. Horodecki, Nat. Comm. 6 7908(2015)

Quantum Darwinism – BPH's approach

Def. EWt-dynamics for the subset $B_{S_t} = (B_j)_{j \in S_t}$: $\Phi^{B_{S_t}} := \text{Tr}_{\mathcal{E} \setminus B_{S_t}} \circ \Phi$



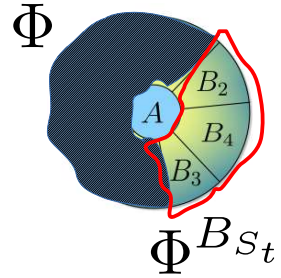
Theorem (Thm 2 of ref*)



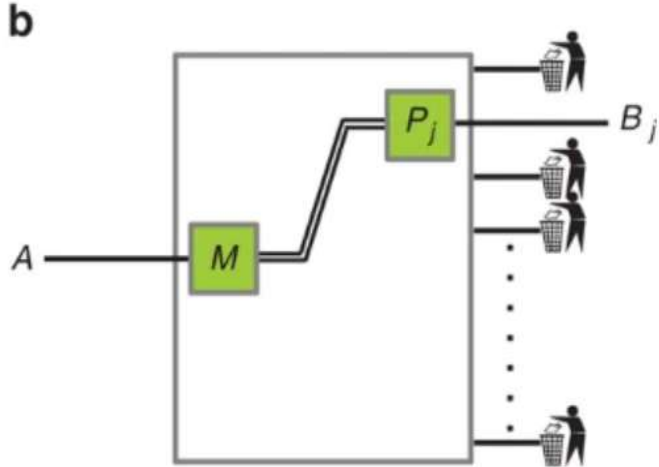
This Thm tells us that EWt-dynamics already leads to emergence of some objectivity: **objectivity of observables**

Quantum Darwinism – BPH's approach

Def. EWt-dynamics for the subset $B_{S_t} = (B_j)_{j \in S_t}$: $\Phi^{B_{S_t}} := \text{Tr}_{\mathcal{E} \setminus B_{S_t}} \circ \Phi$



Theorem (Thm 2 of ref*)



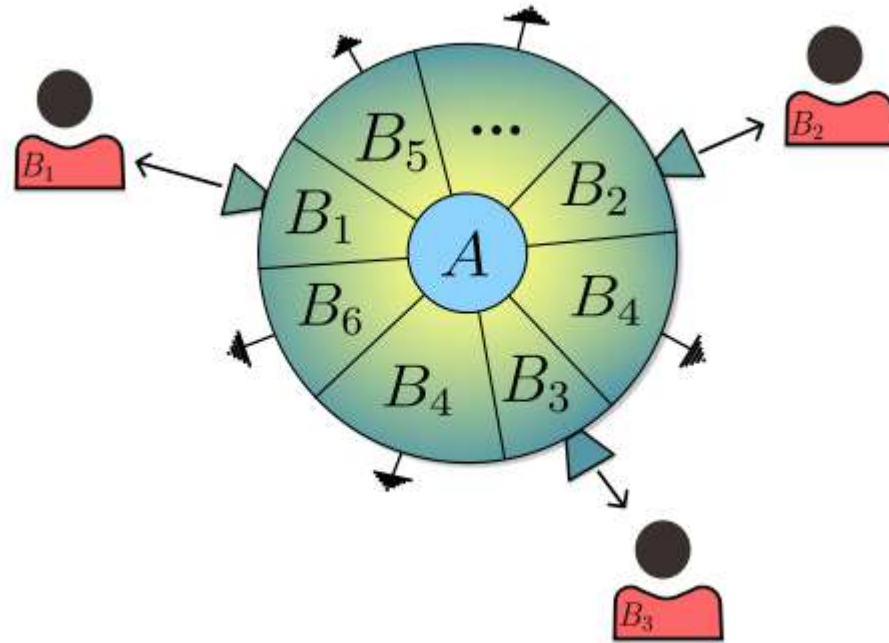
Assumptions:

- We consider portions in which the approximation is valid
- Infinite environment $\approx \rightarrow =$

$$\Phi^{B_{S_t}} = \sum_k \text{Tr}\{\tilde{E}_k \rho\} \sigma_k^{B_{S_t}}$$

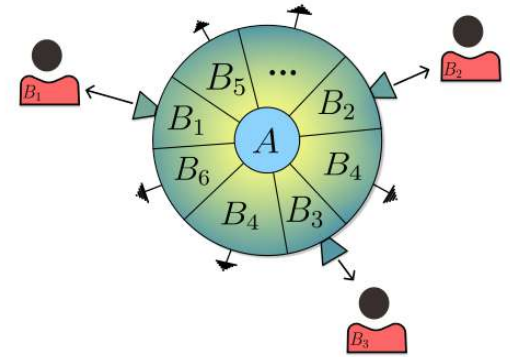
*F. Brandão, M. Piani, P. Horodecki, Nat. Comm. 6 7908(2015)

Quantum Darwinism – BPH's approach



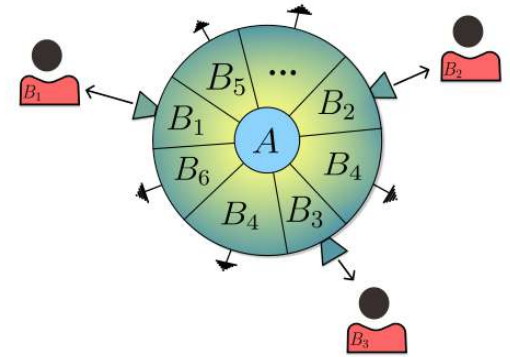
Quantum Darwinism – BPH's approach

$$\Phi^{B_{S_t}}(\rho^A) = \sum_k \text{Tr}\{\tilde{E}_k \rho^A\} \sigma_k^{B_{S_t}}$$



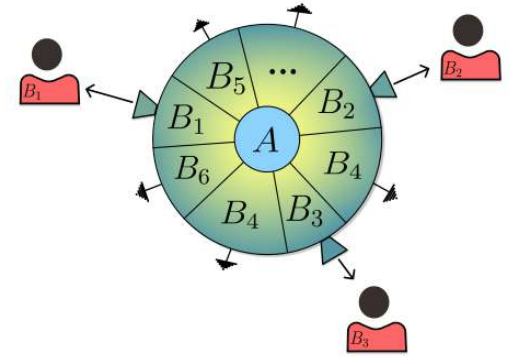
Quantum Darwinism – BPH's approach

$$\Phi^{B_j}(\rho^A) = \sum_k \text{Tr}\{\tilde{E}_k \rho^A\} \sigma_k^{B_j}$$



Quantum Darwinism – BPH's approach

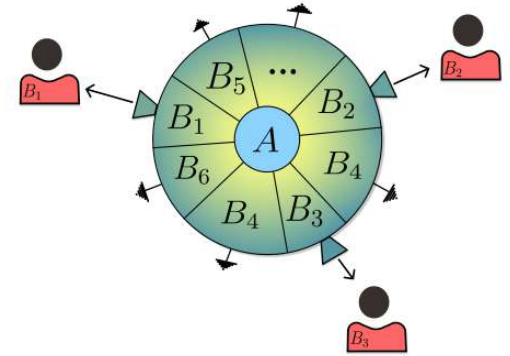
$$\Phi^{B_j}(\rho^A) = \sum_k \text{Tr}\{\tilde{E}_k \rho^A\} \sigma_k^{B_j}$$



QD assumption: $(\sigma_k)_k$ encode well enough the labels $(k)_k$!

Quantum Darwinism – BPH's approach

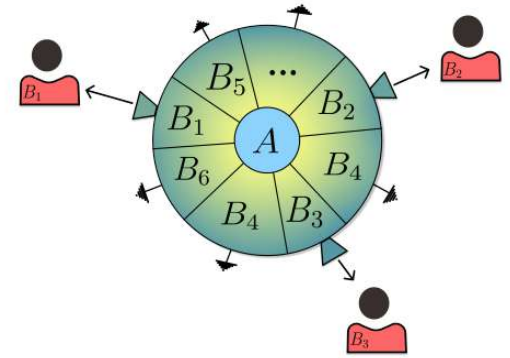
$$\Phi^{B_j}(\rho^A) = \sum_k \text{Tr}\{\tilde{E}_k \rho^A\} \sigma_k^{B_j}$$



QD assumption: $(\sigma_k)_k$ are sufficiently distinguishable!

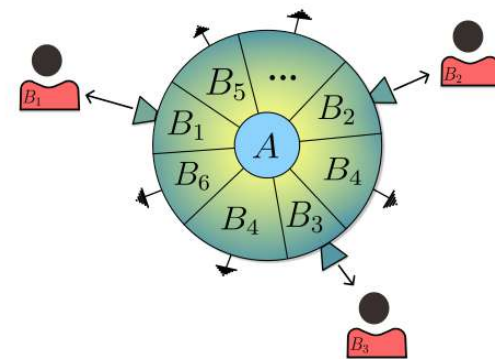
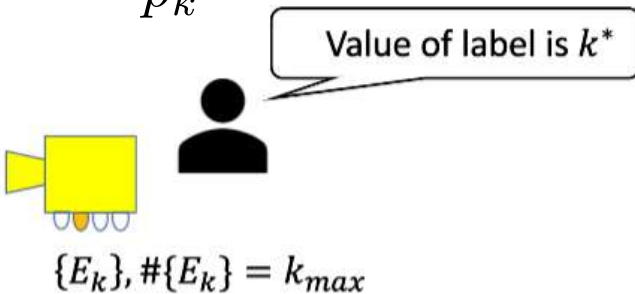
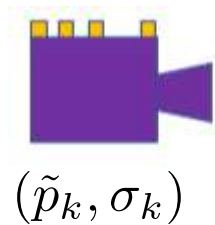
Quantum Darwinism – BPH's approach

$$\Phi^{B_j}(\rho^A) = \sum_k \underbrace{\text{Tr}\{\tilde{E}_k \rho^A\}}_{\tilde{p}_k} \sigma_k^{B_j}$$



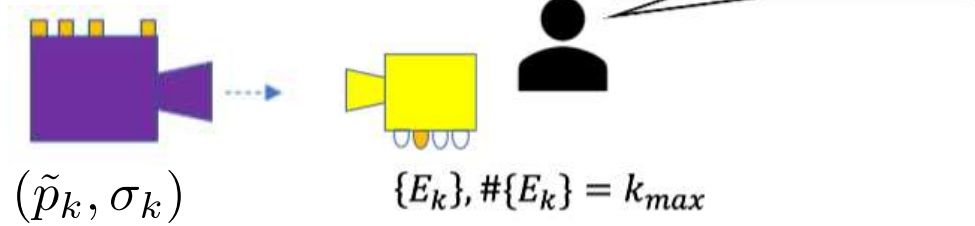
Quantum Darwinism – BPH's approach

$$\Phi^{B_j}(\rho^A) = \sum_k \underbrace{\text{Tr}\{\tilde{E}_k \rho^A\}}_{\tilde{p}_k} \sigma_k^{B_j}$$

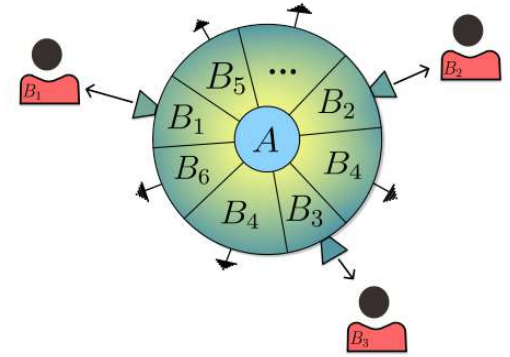


Quantum Darwinism – BPH's approach

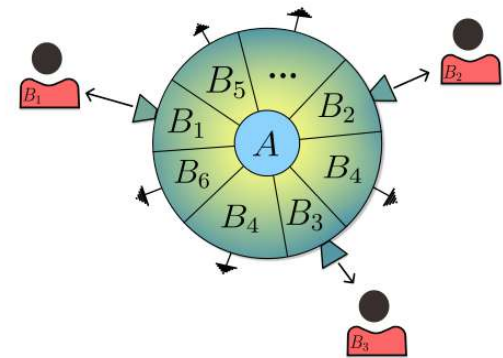
$$\Phi^{B_j}(\rho^A) = \sum_k \underbrace{\text{Tr}\{\tilde{E}_k \rho^A\}}_{\tilde{p}_k} \sigma_k^{B_j}$$



$$p_{\text{guess}}[(\tilde{p}_k, \sigma_k^{B_j})_k] := \max_{\{F_k\}} \sum_k \tilde{p}_k \text{Tr}\{F_k \sigma_k^{B_j}\}.$$



Quantum Darwinism – BPH's approach

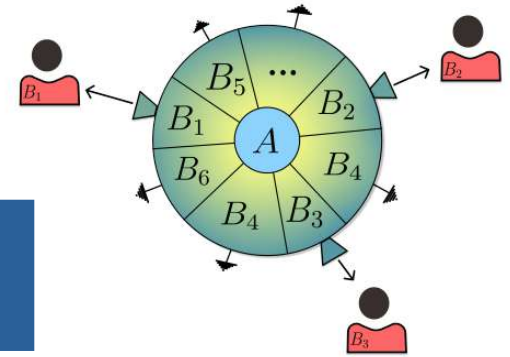


Def. Consider an EWt-dynamics $\Phi^{B_{S_t}}$ and t partial traces Φ^{B_j} .

We say that a QD_η process occurs if, for all $B_j \in B_{S_t}$, $\min_{\rho^A} p_{\text{guess}} \geq \eta$.

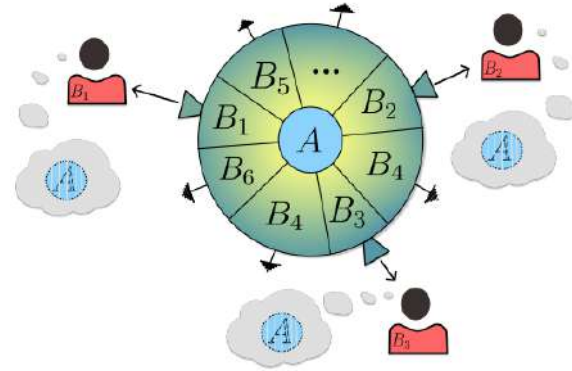
Quantum Darwinism – BPH's approach

Quantum Darwinism: the environment encodes information regarding system A



Def. Consider an EWt-dynamics $\Phi^{B_{S_t}}$ and t partial traces Φ^{B_j} .
We say that a QD_η process occurs if, for all $B_j \in B_{S_t}$, $\min_{\rho^A} p_{\text{guess}} \geq \eta$.

Quantum Darwinism – BPH's approach



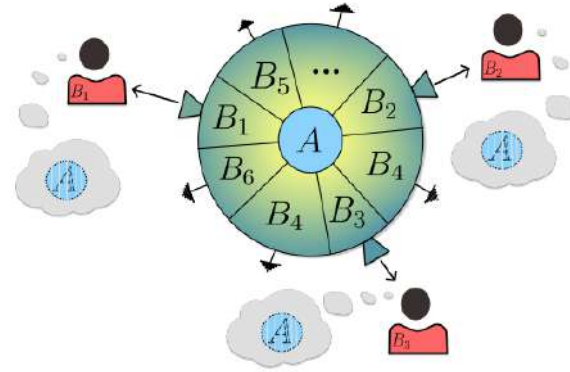
Prop.(adapt. from ref.*)

$$QD_{\eta}$$

EWt-dynamics + good enough encoding \Rightarrow Bobs are likely to agree!

*F. Brandão, M. Piani, P. Horodecki, Nat. Comm. 6 7908(2015)

Quantum Darwinism – BPH's approach



Prop.(adapt. from ref.*)

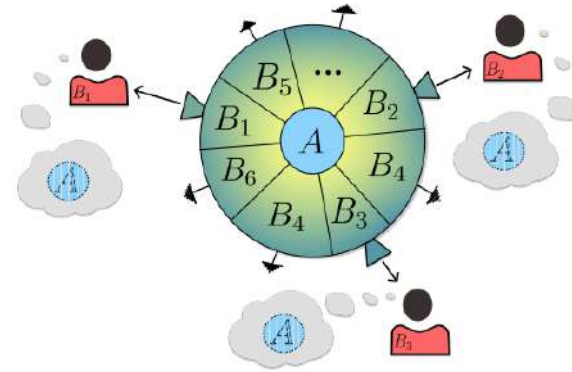
QD_η

EWt-dynamics + good enough encoding \Rightarrow **objectivity of outcomes**

Quantum Darwinism – BPH's approach

EW dynamics describes our generic experience

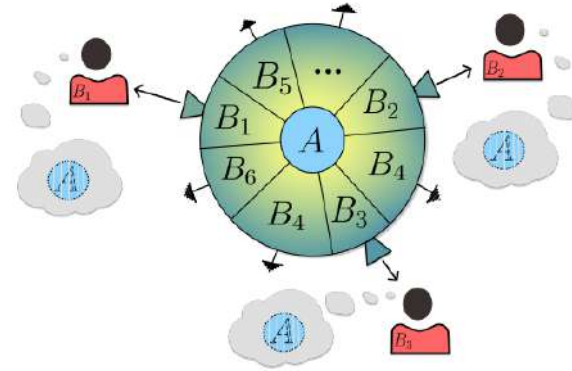
Specifics of the interaction do not matter!



Quantum Darwinism – BPH's approach

EW dynamics describes our generic experience

Specifics of the interaction do not matter!



There is no restriction on the observable
selected by the dynamics

(\tilde{E}_k) can be an IC-POVM!

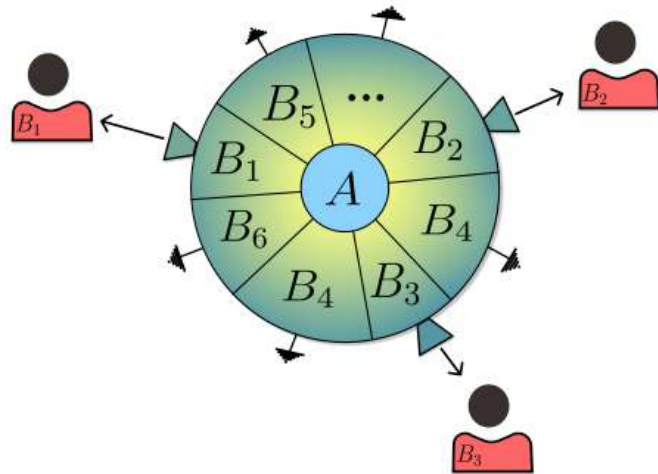


Results

Taking Spekkens' notion
of (non)contextuality

Results

Taking Spekkens' notion
of (non)contextuality



Taking Spekkens' notion
of (non)contextuality

Lemma 1

Consider an EWt-dynamics and the reduced $(\Phi^{B_j})_j$.

For those $B_j \in B_{S_t}$ in which $(\sigma_k^{B_j})_k$ are affinely independent, there exists a noncontextual model.

Taking Spekkens' notion
of (non)contextuality

Lemma 1

Consider an EWt-dynamics and the reduced $(\Phi^{B_j})_j$.

For those $B_j \in B_{S_t}$ in which $(\sigma_k^{B_j})_k$ are affinely independent, there exists a noncontextual model.

Taking Spekkens' notion
of (non)contextuality

Lemma 1

Consider an EWt-dynamics and the reduced $(\Phi^{B_j})_j$.

For those $B_j \in B_{S_t}$ in which $(\sigma_k^{B_j})_k$ are affinely independent, there exists a noncontextual model.

EWt-dynamics + affine independence \Rightarrow noncontextuality

Results

Taking Spekkens' notion
of (non)contextuality

Lemma 1

Consider an EWt-dynamics and the reduced $(\Phi^{B_j})_j$.

For those $B_j \in B_{S_t}$ in which $(\sigma_k^{B_j})_k$ are affinely independent, there exists a noncontextual model.

EWt-dynamics + affine independence \Rightarrow noncontextuality

Not operational...

Taking Spekkens' notion
of (non)contextuality

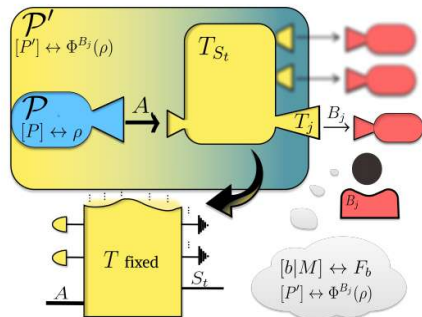
Lemma 2

Consider a reduced dynamics $(\tilde{E}_k, \sigma_k^{B_j})$. There exists $\hat{P}[(\tilde{E}_k)]$ s.t.,
if $p_{\text{guess}}[(\tilde{p}_k, \sigma_k^{B_j})] > \hat{P}[(\tilde{E}_k)] \forall \rho^A$, then $(\sigma_k^{B_j})$ are affinely independent.

Taking Spekkens' notion of (non)contextuality

Lemma 2

Consider a reduced dynamics $(\tilde{E}_k, \sigma_k^{B_j})$. There exists $\hat{P}[(\tilde{E}_k)]$ s.t., if $p_{\text{guess}}[(\tilde{p}_k, \sigma_k^{B_j})] > \hat{P}[(\tilde{E}_k)] \forall \rho^A$, then $(\sigma_k^{B_j})$ are affinely independent.

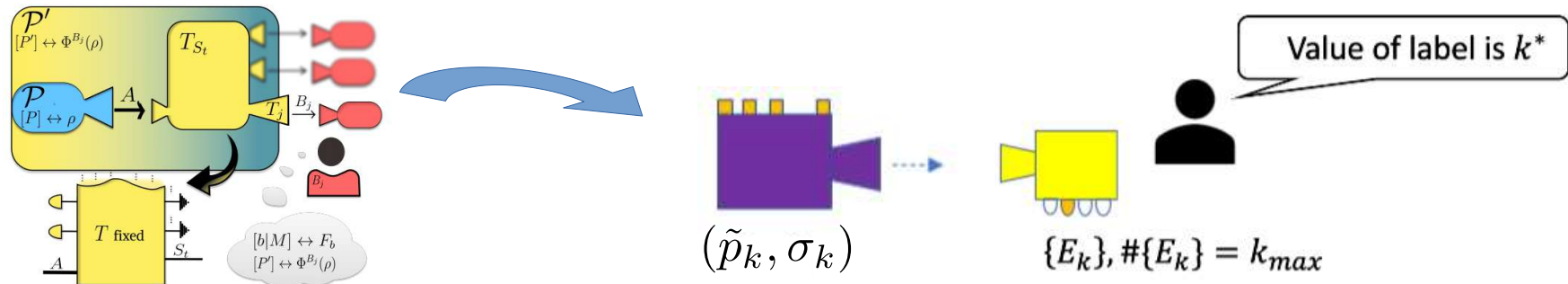


$$\sigma_P^{B_j} = \sum_k \tilde{p}_k \sigma_k^{B_j}$$

Taking Spekkens' notion of (non)contextuality

Lemma 2

Consider a reduced dynamics $(\tilde{E}_k, \sigma_k^{B_j})$. There exists $\hat{P}[(\tilde{E}_k)]$ s.t., if $p_{\text{guess}}[(\tilde{p}_k, \sigma_k^{B_j})] > \hat{P}[(\tilde{E}_k)] \forall \rho^A$, then $(\sigma_k^{B_j})$ are affinely independent.



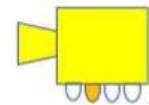
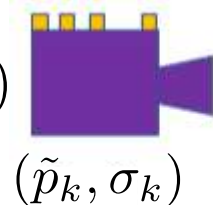
Results

Taking Spekkens' notion of (non)contextuality

Lemma 2

Consider a reduced dynamics $(\tilde{E}_k, \sigma_k^{B_j})$. There exists $\hat{P}[(\tilde{E}_k)]$ s.t., if $p_{\text{guess}}[(\tilde{p}_k, \sigma_k^{B_j})] > \hat{P}[(\tilde{E}_k)] \forall \rho^A$, then $(\sigma_k^{B_j})$ are affinely independent.

$p_{\text{guess}} > \hat{P} \Rightarrow$ affine independent $(\sigma_k^{B_j})$



$\{E_k\}, \#\{E_k\} = k_{\max}$



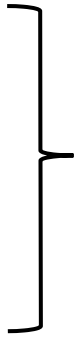
Value of label is k^*

Results

Taking Spekkens' notion
of (non)contextuality

Lemma 1

Lemma 2



Taking Spekkens' notion
of (non)contextuality

Lemma 1

Lemma 2

Theorem (Main result)

If a process QD_η occurs with $\eta > \hat{P}[\tilde{E}_k]$, every Bob $B_j \in B_{S_t}$ can construct a noncontextual ontological model

Taking Spekkens' notion
of (non)contextuality

Lemma 1

Lemma 2

Theorem (Main result)

If a process QD_η occurs with $\eta > \hat{P}[\tilde{E}_k]$, every Bob $B_j \in B_{S_t}$ can construct a noncontextual ontological model

$\text{QD}_\eta \Rightarrow$ noncontextuality for every Bob

Results – relaxing assumption

Infinite environment $\Phi^{B_{S_t}} = \sum_k \text{Tr}\{\tilde{E}_k \rho\} \sigma_k^{B_{S_t}}$

Finite environment $\Phi^{B_{S_t}} \approx \sum_k \text{Tr}\{\tilde{E}_k \rho\} \sigma_k^{B_{S_t}}$

Results – relaxing assumption

Infinite environment $\Phi^{B_{S_t}} = \sum_k \text{Tr}\{\tilde{E}_k \rho\} \sigma_k^{B_{S_t}}$

Finite environment $\Phi^{B_{S_t}} \approx \sum_k \text{Tr}\{\tilde{E}_k \rho\} \sigma_k^{B_{S_t}} \longrightarrow d(p) \leq \epsilon(N, t, d_A)$

Results – relaxing assumption

Infinite environment $\Phi^{B_{S_t}} = \sum_k \text{Tr}\{\tilde{E}_k \rho\} \sigma_k^{B_{S_t}}$

Finite environment $\Phi^{B_{S_t}} \approx \sum_k \text{Tr}\{\tilde{E}_k \rho\} \sigma_k^{B_{S_t}} \rightarrow d(p) \leq \overset{\approx 0}{\epsilon(N, t, d_A)}$

Conclusions and perspectives

- Noncontextuality generically emerges out of quantum Darwinism in infinite environments
-
-
-
-

Conclusions and perspectives

- Noncontextuality generically emerges out of quantum Darwinism in infinite environments
 - It can emerge even if Darwinism fails (EW+affine independence are enough);
 - Objectivity emerging from Darwinism can be considered classical, even in the broader case
-
-

Conclusions and perspectives

- Noncontextuality generically emerges out of quantum Darwinism in infinite environments
 - It can emerge even if Darwinism fails (EW+affine independence are enough);
 - Objectivity emerging from Darwinism can be considered classical, even in the broader case
- In finite environments, contextuality is strongly constrained;
-

Conclusions and perspectives

- Noncontextuality generically emerges out of quantum Darwinism in infinite environments
 - It can emerge even if Darwinism fails (EW+affine independence are enough);
 - Objectivity emerging from Darwinism can be considered classical, even in the broader case
- In finite environments, contextuality is strongly constrained;
- We have a bound on state discrimination tasks able to witness if the coding states are vertices of a simplex in the state space

Conclusions and perspectives

- Noncontextuality generically emerges out of quantum Darwinism in infinite environments
 - It can emerge even if Darwinism fails (EW+affine independence are enough);
 - Objectivity emerging from Darwinism can be considered classical, even in the broader case
- In finite environments, contextuality is strongly constrained;
- We have a bound on state discrimination tasks able to witness if the coding states are vertices of a simplex in the state space

Perspectives

- Do other classical limits (strong QD) allow for emergence of noncontextuality?
- Relation to non-Markovianity?
-

Conclusions and perspectives

- Noncontextuality generically emerges out of quantum Darwinism in infinite environments
 - It can emerge even if Darwinism fails (EW+affine independence are enough);
 - Objectivity emerging from Darwinism can be considered classical, even in the broader case
- In finite environments, contextuality is strongly constrained;
- We have a bound on state discrimination tasks able to witness if the coding states are vertices of a simplex in the state space

Perspectives

- Do other classical limits (strong QD) allow for emergence of noncontextuality?
- Relation to non-Markovianity?
- If we take the special case of the multiple observers with weak interaction, can we protect contextuality?

THANKS FOR LISTENING!



A word cloud of 'Thank You' in various languages and scripts, including:

- GRACIAS
- ARIGATO
- SHUKURIA
- GOZAIMASHITA
- EFCHARISTO
- JUSPAXAR
- DANKSCHEEN
- TASHAKKUR ATU
- YAQHANYELAY
- SUKSAMA
- EKHMET
- MEHRBANI
- MAAKE
- GRAZIE
- KOMAPSUNNIDA
- BIYAN
- SHUKRIA
- TINGKI
- THANK
- YOU
- BOLZIN
- MERCI

Appendix

Appendix 1

KS

Contextuality

KS Contextuality



Pic by Konrad
Jacobs, Erlangen



Pic from Princeton
website

KS Contextuality

Are there answers to all questions,
even those that were never asked?



Pic by Konrad
Jacobs, Erlangen



Pic from Princeton
website

KS Contextuality

Is Quantum Theory compatible with the idea of an underlying reality where measurements play only the role of revealing pre-determined values?



Pic by Konrad
Jacobs, Erlangen



Pic from Princeton
website

KS Contextuality

Is Quantum Theory compatible with the idea of an underlying reality where measurements play only the role of revealing pre-determined values?

Compatibility with an ontological model assigning deterministic values.

Probabilities arise as a result of our ignorance.



Pic by Konrad Jacobs, Erlangen



Pic from Princeton website

KS Contextuality

Is Quantum Theory compatible with the idea of an underlying reality where measurements play only the role of revealing pre-determined values?

Answer: Yes, but this assignment must be context-dependent.

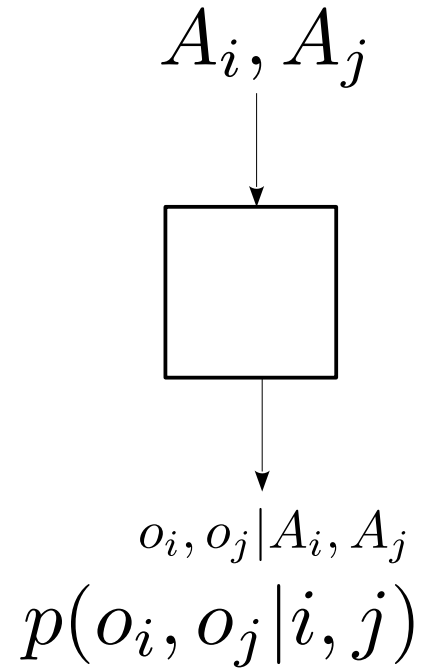


Pic by Konrad
Jacobs, Erlangen



Pic from Princeton
website

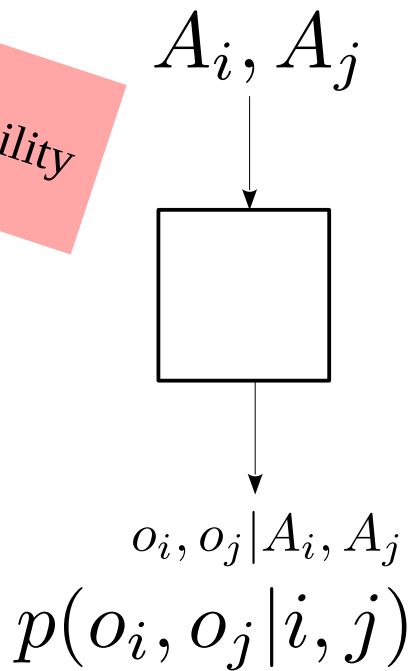
KS Contextuality



KS Contextuality

Specify
compatibility
structure!

$$\Gamma \equiv \{X, \mathcal{C}, O\}$$

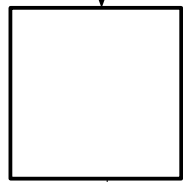


KS Contextuality

Specify
compatibility
structure!

$$\Gamma \equiv \{X, \mathcal{C}, O\}$$

A_i, A_j



$o_i, o_j | A_i, A_j$

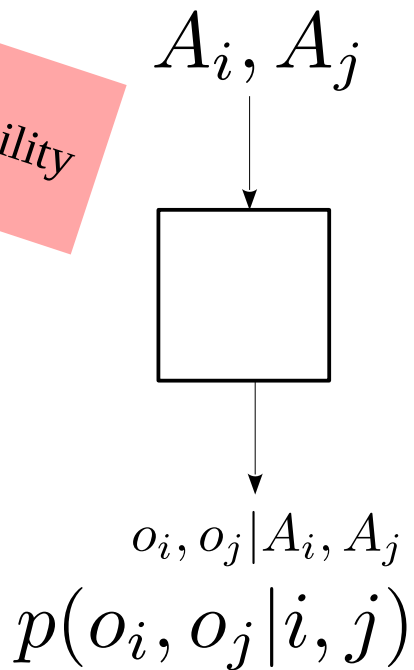
$$p(o_i, o_j | i, j)$$

$$p(o_i, o_j | i, j) \stackrel{?}{=} \sum_{\lambda} p(\lambda) p(o_i | i, \lambda) p(o_j | j, \lambda)$$

KS Contextuality

Specify
compatibility
structure!

$$\Gamma \equiv \{X, \mathcal{C}, O\}$$



$$p(o_i, o_j | i, j) \stackrel{?}{=} \sum_{\lambda} p(\lambda) p(o_i | i, \lambda) p(o_j | j, \lambda)$$
$$\iff \exists p(o_1, \dots, o_{|X|} | \lambda)$$

A. Fine Phys. Rev. Lett. 48, 291

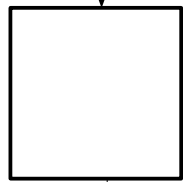
S. Abramsky and A. Brandenburger 2011 New J. Phys. 13 113036

KS Contextuality

Specify
compatibility
structure!

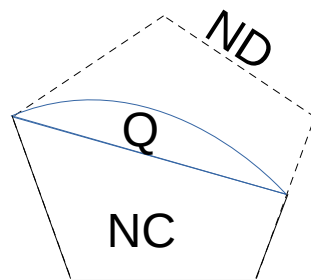
$$\Gamma \equiv \{X, \mathcal{C}, O\}$$

A_i, A_j



$o_i, o_j | A_i, A_j$

$$p(o_i, o_j | i, j)$$



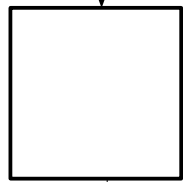
$$p(o_i, o_j | i, j) \stackrel{?}{=} \sum_{\lambda} p(\lambda) p(o_i | i, \lambda) p(o_j | j, \lambda)$$

KS Contextuality

Specify
compatibility
structure!

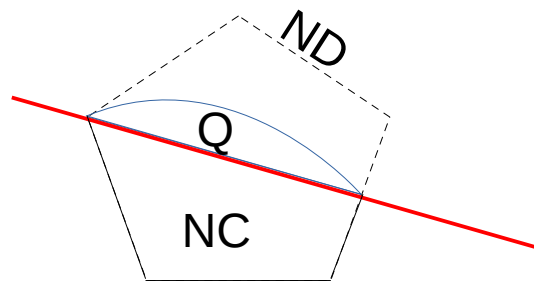
$$\Gamma \equiv \{X, \mathcal{C}, O\}$$

A_i, A_j



$o_i, o_j | A_i, A_j$

$$p(o_i, o_j | i, j)$$



KS noncontextuality
inequalities

$$p(o_i, o_j | i, j) \stackrel{?}{=} \sum_{\lambda} p(\lambda) p(o_i | i, \lambda) p(o_j | j, \lambda)$$

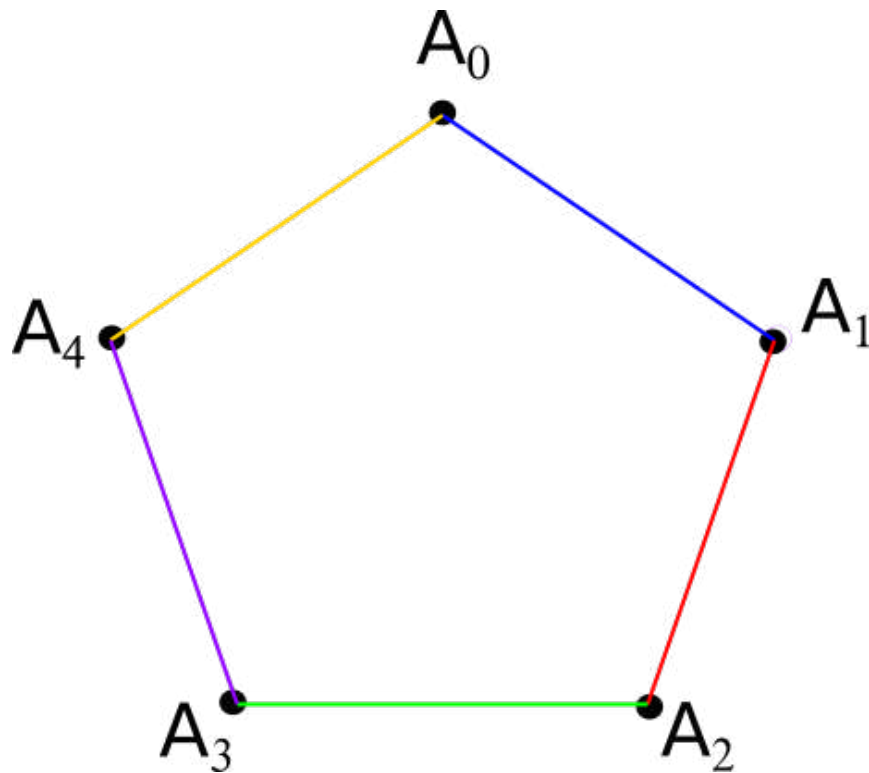
N-cycle
inequalities

Understanding the inequality

$$\sum_{i=0}^{N-1} \langle A_i A_{i+1} \rangle \geq -N + 2$$

-3

$$p(o_i, o_j | i, j) \stackrel{?}{=} \sum_{\lambda} p(\lambda) p(o_i | i, \lambda) p(o_j | j, \lambda)$$

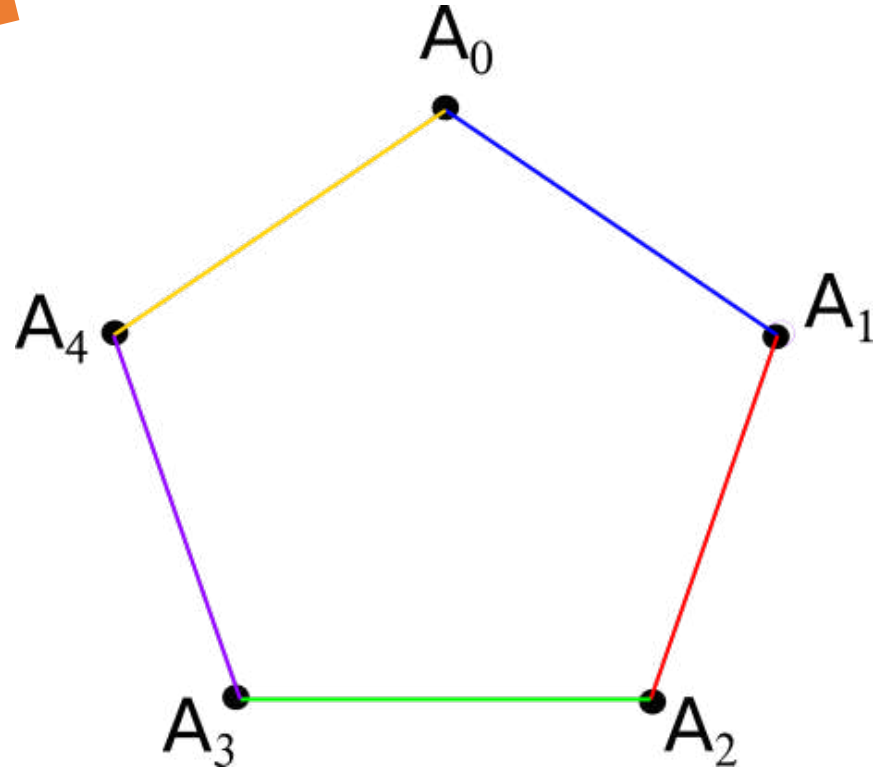


Foolish non-contextual assignment

Understanding the inequality

$$\sum_{i=0}^{N-1} \langle A_i A_{i+1} \rangle \geq -N + 2$$

-3

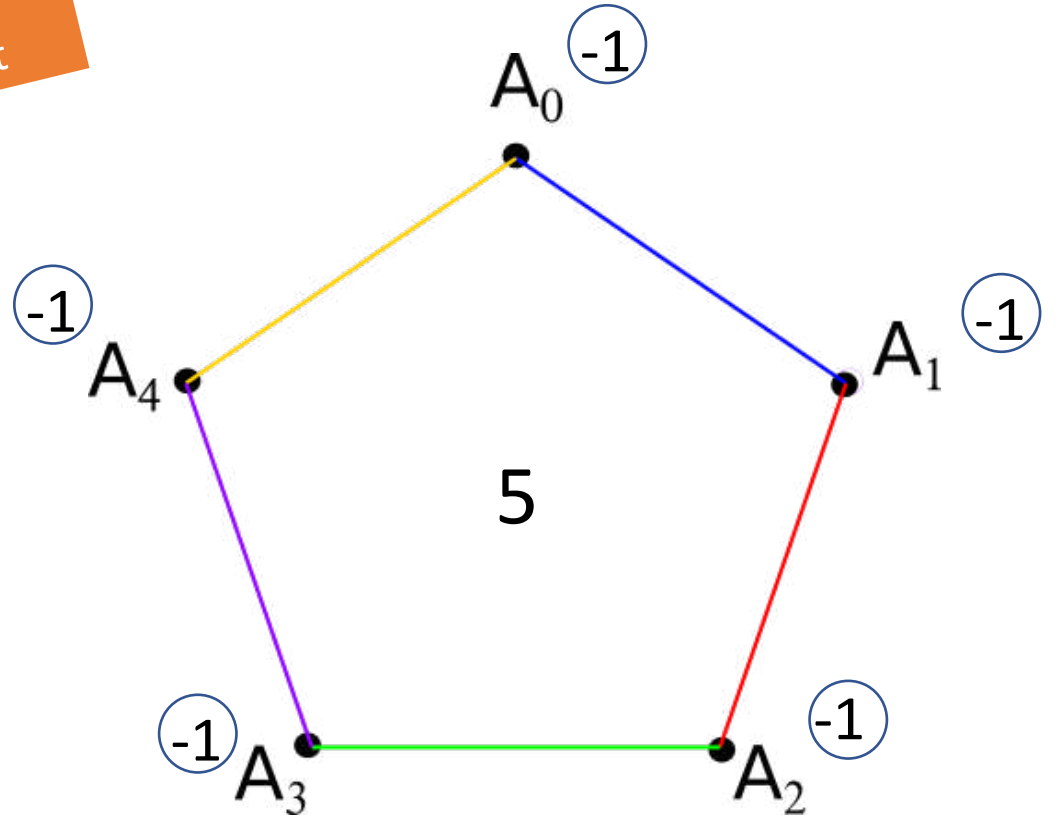


Foolish non-contextual assignment

Understanding the inequality

$$\sum_{i=0}^{N-1} \langle A_i A_{i+1} \rangle \geq -N + 2$$

-3



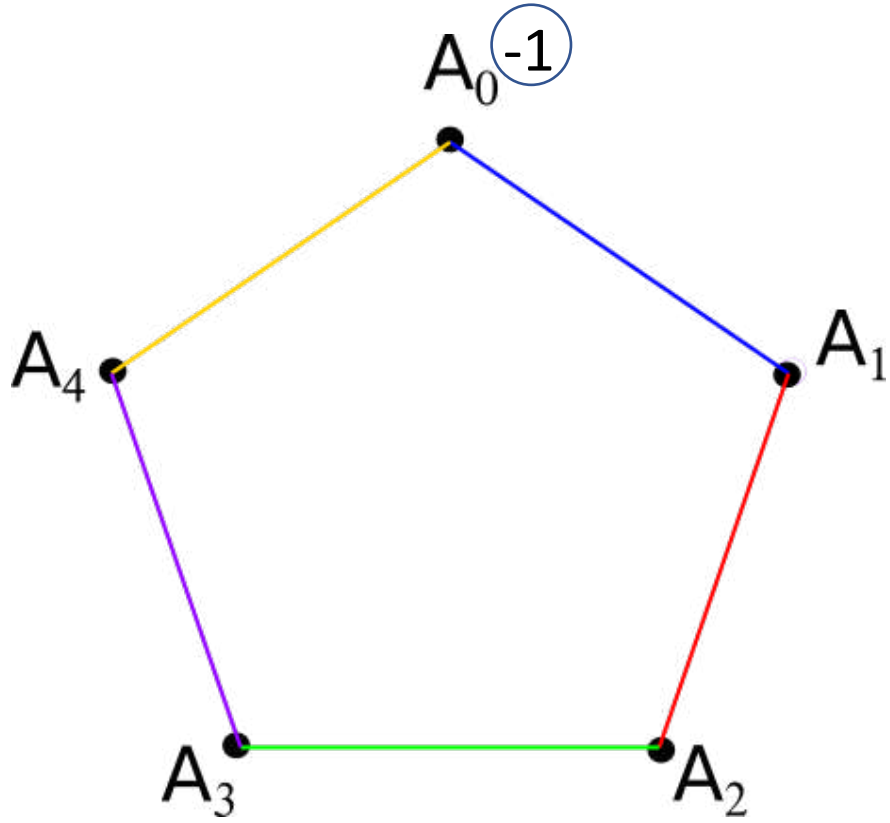
We need to alternate values...

Alternating values, non-contextual

Understanding the inequality

$$\sum_{i=0}^{N-1} \langle A_i A_{i+1} \rangle \geq -N + 2$$

-3

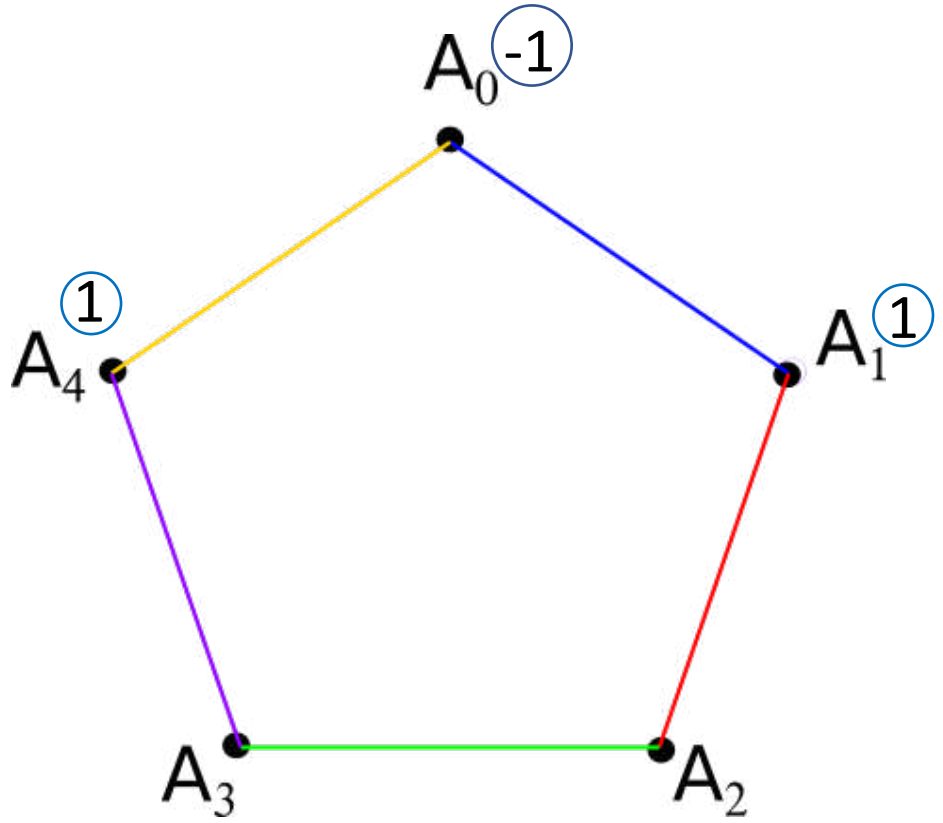


Alternating values, non-contextual

Understanding the inequality

$$\sum_{i=0}^{N-1} \langle A_i A_{i+1} \rangle \geq -N + 2$$

-3

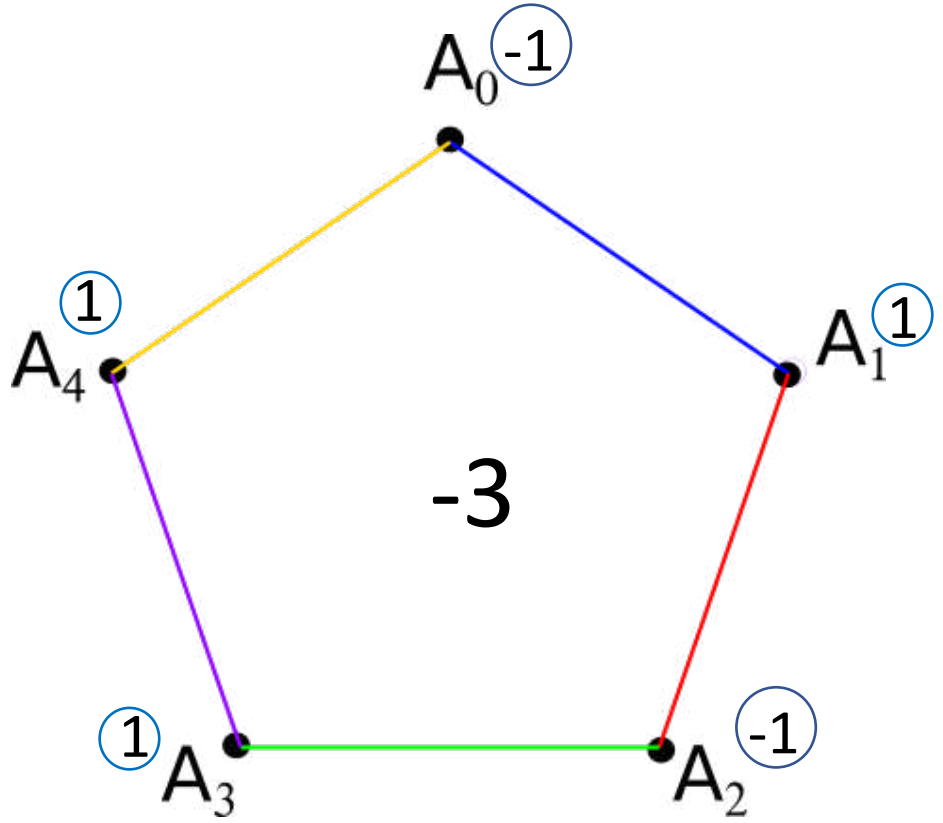


Alternating values, non-contextual

Understanding the inequality

$$\sum_{i=0}^{N-1} \langle A_i A_{i+1} \rangle \geq -N + 2$$

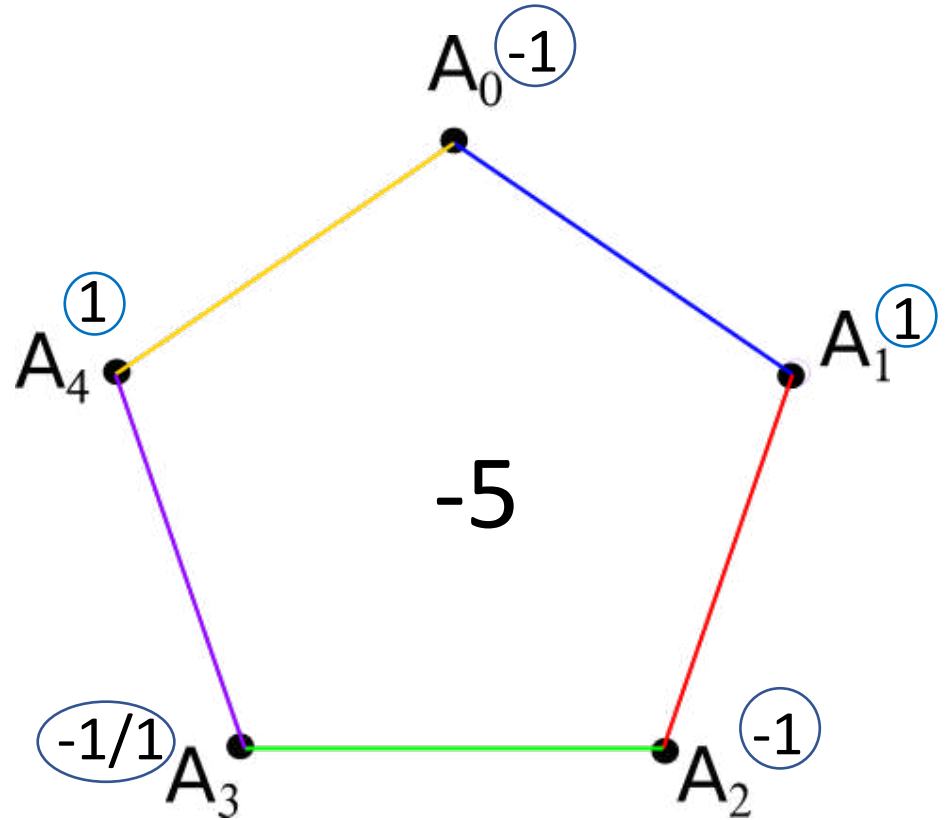
-3



Understanding the inequality

$$\sum_{i=0}^{N-1} \langle A_i A_{i+1} \rangle \geq -N + 2$$

-3



Contextual/Disturbing assignment

Odd N-cycle: new inequalities

$$\sum_{i=0}^{N-1} \langle A_i A_{i+1} \rangle \geq -N+2$$

$$\sum_{i=0}^{N-1} p(o_i = o_{i+1} | i, i+1) \geq 1$$

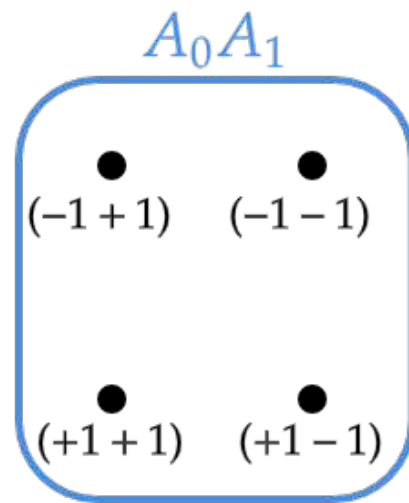
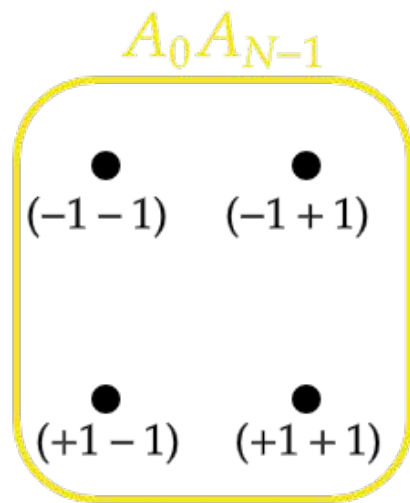
$$\sum_{i=0}^{N-1} p(o_i \neq o_{i+1} | i, i+1) \leq N-1$$

Odd N-cycle: new inequalities

- Non-Disturbance

$$\sum_{i=0}^{N-1} p(o_i = o_{i+1} | i, i+1) \geq 1$$

$$\sum_{i=0}^{N-1} p(o_i \neq o_{i+1} | i, i+1) \leq N - 1$$

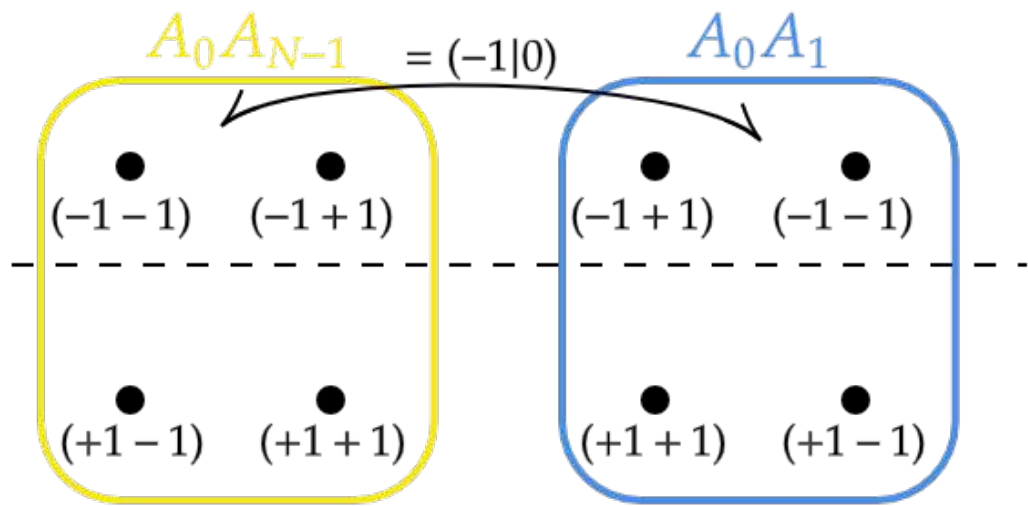


Odd N-cycle: new inequalities

- Non-Disturbance

$$\sum_{i=0}^{N-1} p(o_i = o_{i+1} | i, i+1) \geq 1$$

$$\sum_{i=0}^{N-1} p(o_i \neq o_{i+1} | i, i+1) \leq N - 1$$

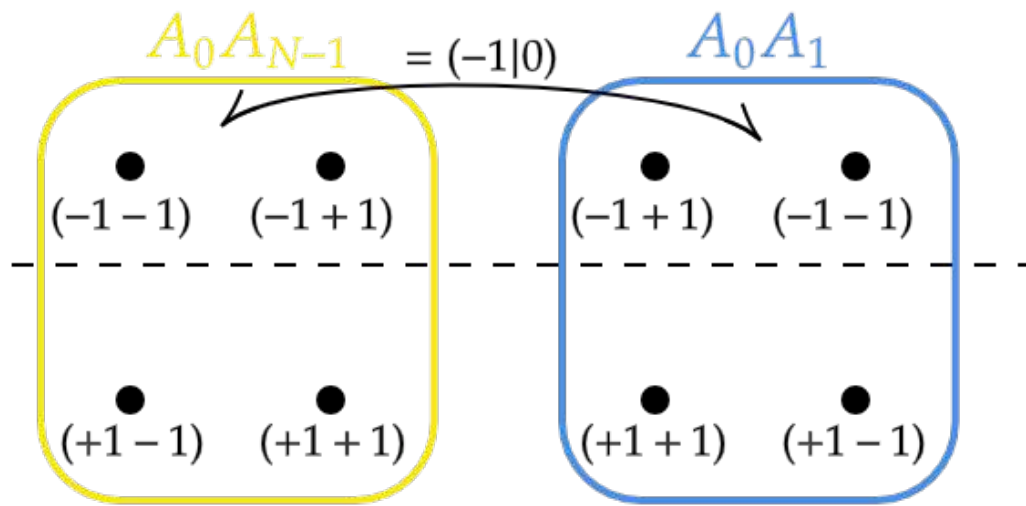


Odd N-cycle: new inequalities

- Non-Disturbance
- Exclusiveness: $p(-1, -1|i, i \pm 1) = 0$

$$\sum_{i=0}^{N-1} p(o_i = o_{i+1} | i, i+1) \geq 1$$

$$\sum_{i=0}^{N-1} p(o_i \neq o_{i+1} | i, i+1) \leq N - 1$$

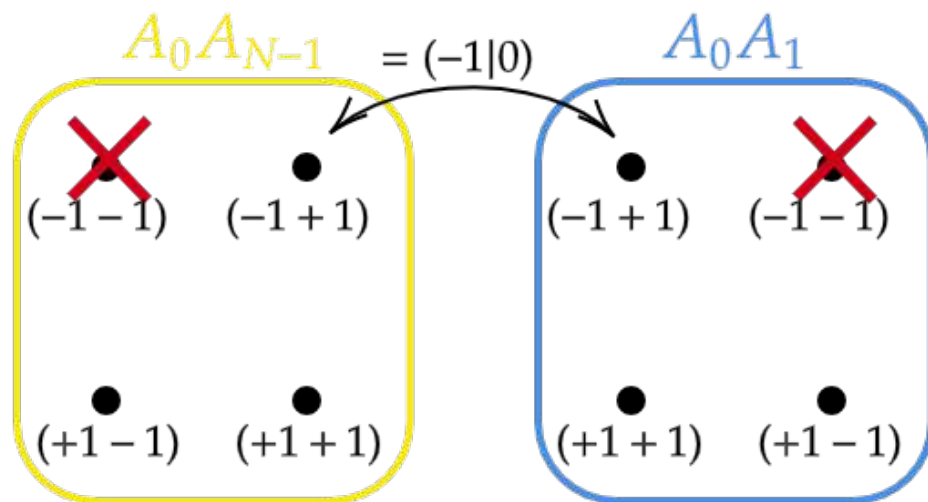


Odd N-cycle: new inequalities

- Non-Disturbance
- Exclusiveness: $p(-1, -1|i, i \pm 1) = 0$

$$\sum_{i=0}^{N-1} p(o_i = o_{i+1} | i, i+1) \geq 1$$

$$\sum_{i=0}^{N-1} p(o_i \neq o_{i+1} | i, i+1) \leq N-1$$

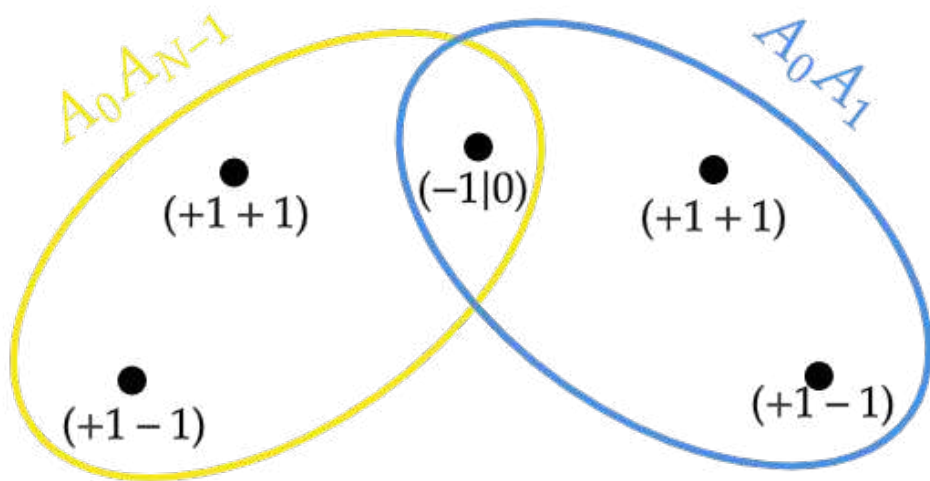


Odd N-cycle: new inequalities

- Non-Disturbance
- Exclusiveness: $p(-1, -1|i, i \pm 1) = 0$

$$\sum_{i=0}^{N-1} p(o_i = o_{i+1} | i, i+1) \geq 1$$

$$\sum_{i=0}^{N-1} p(o_i \neq o_{i+1} | i, i+1) \leq N-1$$



Odd N-cycle: new inequalities

- Non-Disturbance
- Exclusiveness: $p(-1, -1|i, i \pm 1) = 0$

$$\sum_{i=0}^{N-1} p(o_i = o_{i+1} | i, i+1) \geq 1$$

$$\sum_{i=0}^{N-1} p(o_i \neq o_{i+1} | i, i+1) \leq N - 1$$

Odd N-cycle: new inequalities

- Non-Disturbance
- Exclusiveness: $p(-1, -1|i, i \pm 1) = 0$

$$\sum_{i=0}^{N-1} p(o_i = o_{i+1} | i, i+1) \geq 1 \quad \longrightarrow \quad \sum_i p(1, 1 | i, i+1) \geq 1$$

$$\sum_{i=0}^{N-1} p(o_i \neq o_{i+1} | i, i+1) \leq N-1 \quad \longrightarrow \quad \sum_i p(-1 | i) \leq \frac{N-1}{2}$$

Odd N-cycle: new inequalities

- Non-Disturbance
- Exclusiveness: $p(-1, -1|i, i \pm 1) = 0$

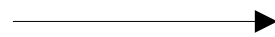
$$\sum_{i=0}^{N-1} p(o_i = o_{i+1} | i, i+1) \geq 1 \quad \longrightarrow \quad \sum_i p(\overset{b_i}{1, 1} | i, i+1) \geq 1$$

$$\sum_{i=0}^{N-1} p(o_i \neq o_{i+1} | i, i+1) \leq N-1 \quad \longrightarrow \quad \sum_i p(-1 | \underset{a_i}{i}) \leq \frac{N-1}{2}$$

Odd N-cycle: new inequalities

- Non-Disturbance
- Exclusiveness: $p(-1, -1|i, i \pm 1) = 0$

$$\sum_{i=0}^{N-1} p(o_i = o_{i+1} | i, i+1) \geq 1$$



$$\sum_i p(b_i) \geq 1$$

$$\sum_{i=0}^{N-1} p(o_i \neq o_{i+1} | i, i+1) \leq N - 1$$



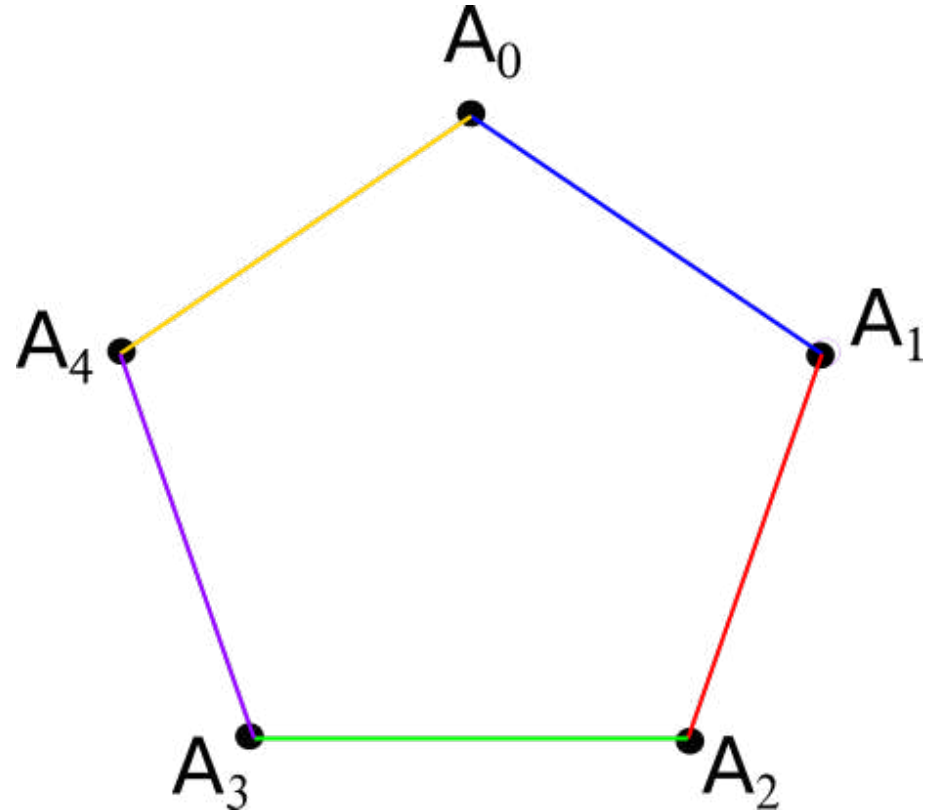
$$\sum_i p(a_i) \leq \frac{N - 1}{2}$$

Odd N-cycle: new inequalities

- Non-Disturbance
- Exclusiveness: $p(-1, -1|i, i \pm 1) = 0$

$$\sum_i p(b_i) \geq 1$$

$$\sum_i p(a_i) \leq \frac{N-1}{2}$$

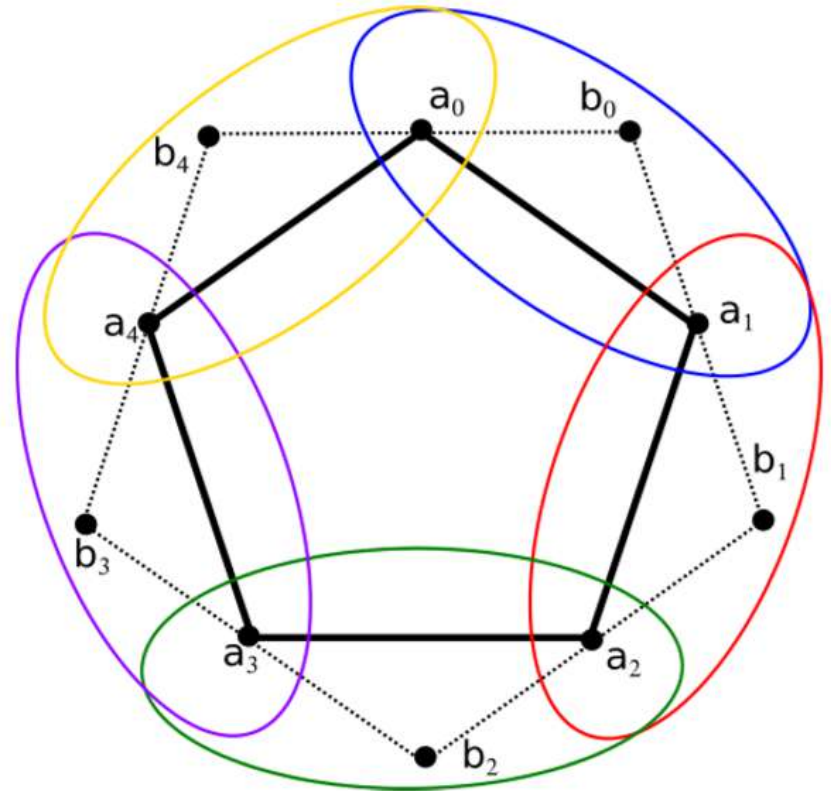


Odd N-cycle: new inequalities

- Non-Disturbance
- Exclusiveness: $p(-1, -1|i, i \pm 1) = 0$

$$\sum_i p(b_i) \geq 1$$

$$\sum_i p(a_i) \leq \frac{N-1}{2}$$

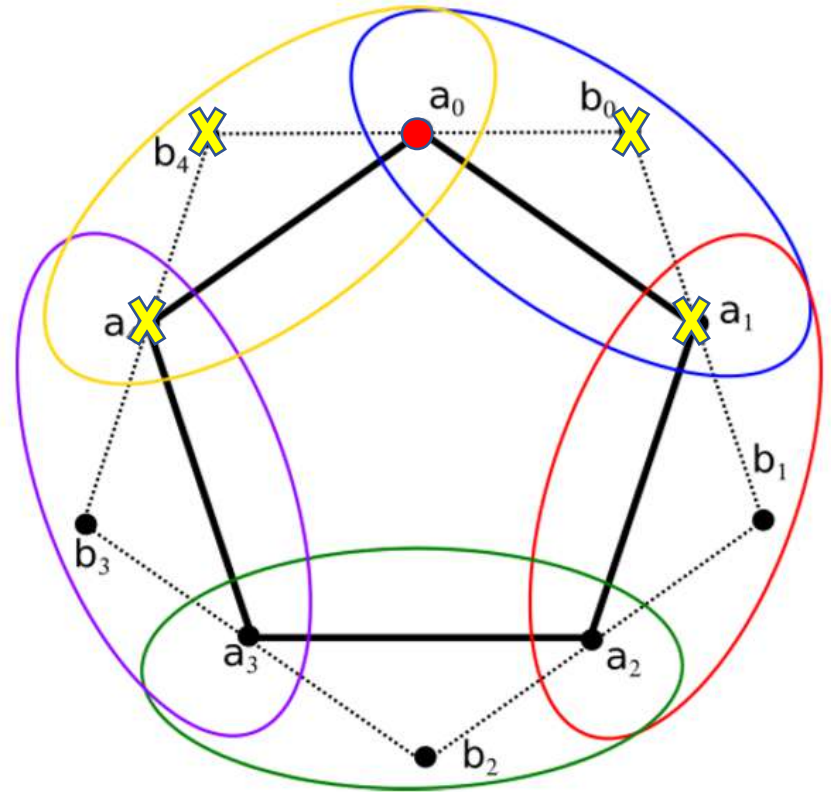


Odd N-cycle: new inequalities

- Non-Disturbance
- Exclusiveness: $p(-1, -1|i, i \pm 1) = 0$

$$\sum_i p(b_i) \geq 1$$

$$\sum_i p(a_i) \leq \frac{N-1}{2}$$

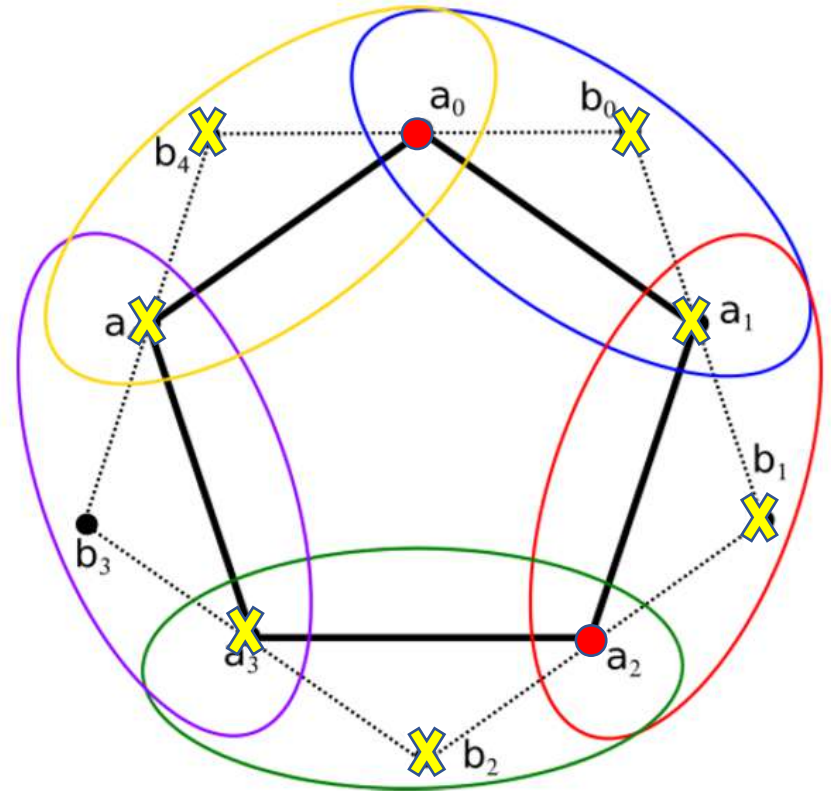


Odd N-cycle: new inequalities

- Non-Disturbance
- Exclusiveness: $p(-1, -1|i, i \pm 1) = 0$

$$\sum_i p(b_i) \geq 1$$

$$\sum_i p(a_i) \leq \frac{N-1}{2}$$



Odd N-cycle: new inequalities

- Non-Disturbance
- Exclusiveness: $p(-1, -1|i, i \pm 1) = 0$

$$\sum_i p(b_i) \geq 1$$

$$\sum_i p(a_i) \leq \frac{N-1}{2}$$

$$\sum_i \langle A_i A_{i+1} \rangle \geq -N + 2$$

The new inequalities can be violated while the original is obeyed!

Odd N-cycle: new inequalities

- Non-Disturbance
- Exclusiveness: $p(-1, -1|i, i \pm 1) = 0$

$$\sum_i p(b_i) \geq 1$$

$$\sum_i p(a_i) \leq \frac{N-1}{2}$$

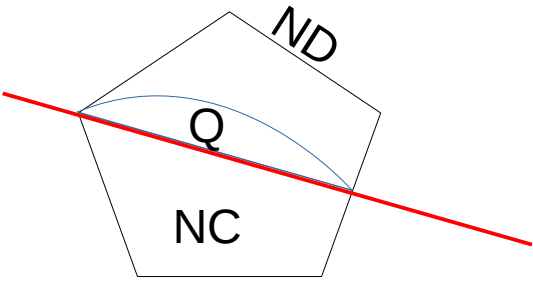
$$\sum_i \langle A_i A_{i+1} \rangle \geq -N + 2$$

The new inequalities can be violated while the original is obeyed!

Example: foolish noncontextual assignment...

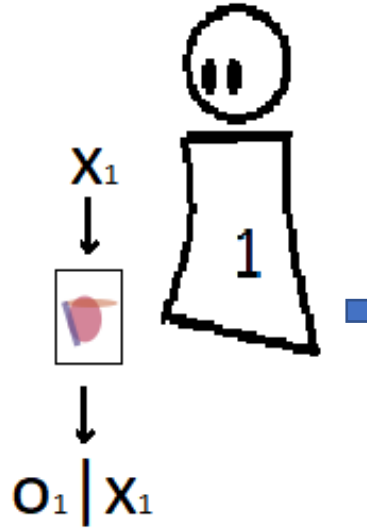
Sequential setup

KS Contextuality



KS noncontextuality
inequalities

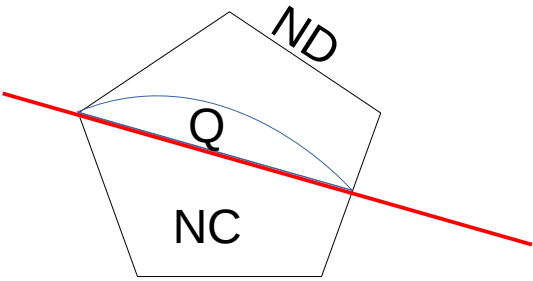
Colocar em
outro lugar?



Usual way to test the inequalities

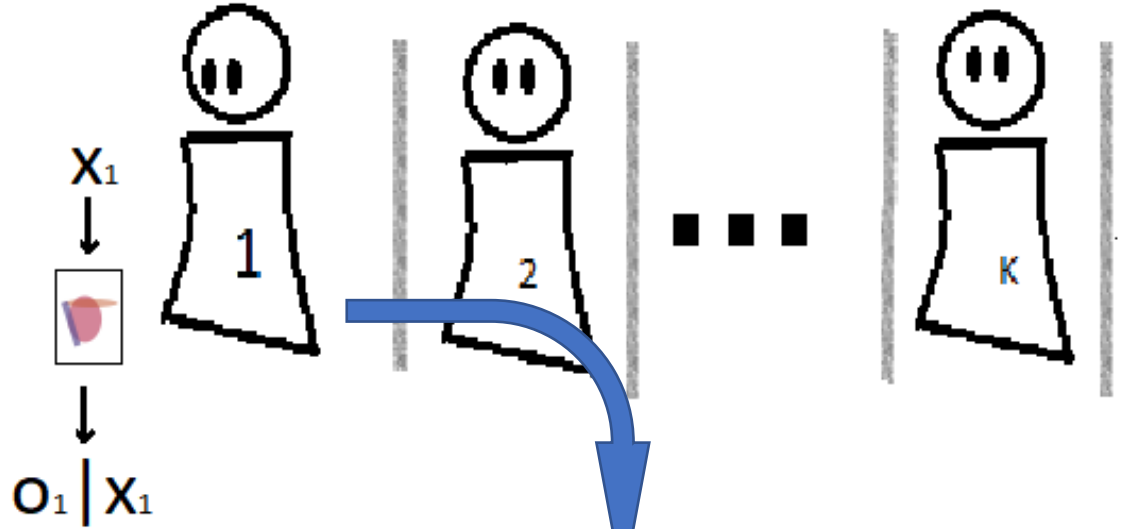


KS Contextuality



KS noncontextuality inequalities

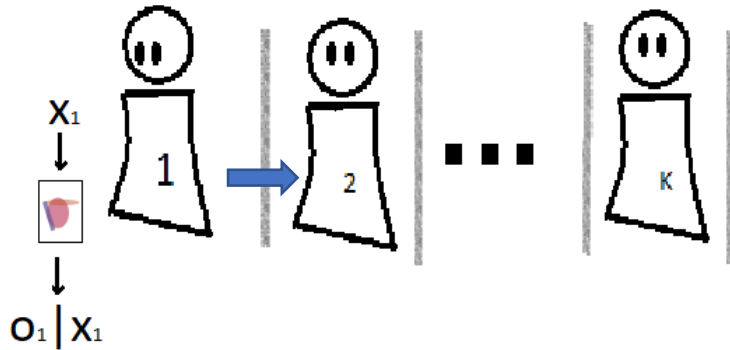
Colocar em outro lugar?



Our multiobserver setup proposal



Adapting to sequential observers



- Which observers still find violations?
 - Does this depend on N ?
 - Does this depend on the measurement protocol?

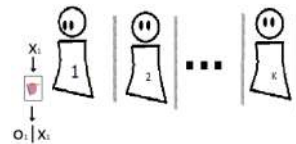
- Already analyzed, in different setups, for non-locality and steering, but lacking for Contextuality.

S. Sasmal et al, Phys. Rev. A98, 012305 (2018)

R. Silva et al, Phys. Rev. Lett.114, 250401 (2015)

D. Das et al, Phys. Rev. A99, 022305 (2019).

Results



$$1. \mathcal{M}_i = \{a_i, b_i, a_{i+1}\}$$

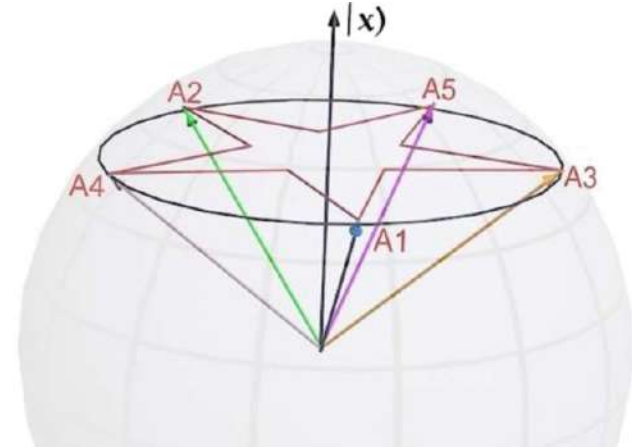
$$2. \mathcal{M}_i^a = \{a_i, \neg a_i\}$$

$$3. \mathcal{M}_i^b = \{b_i, \neg b_i\}$$

Quantum realizations

- Hilbert space of dimension 3

$$A_i = \mathcal{I} - 2|a_i\rangle\langle a_i|$$

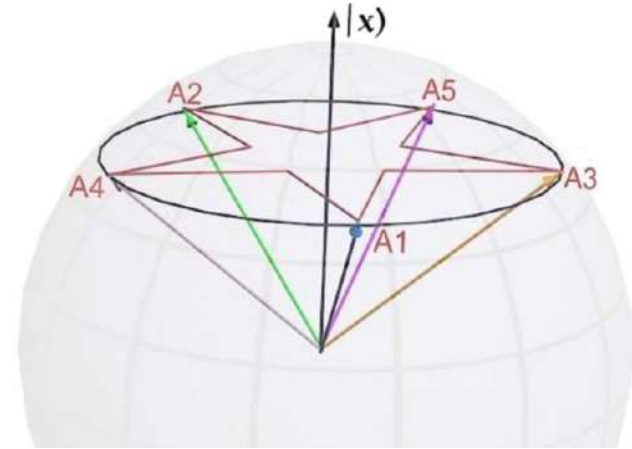


$$\langle a_i | a_{i+1} \rangle = 0$$

Quantum realizations

- Hilbert space of dimension 3

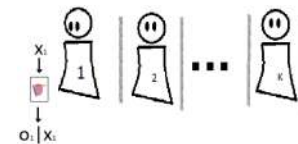
$$A_i = \mathcal{I} - 2|a_i\rangle\langle a_i|$$



$$|\psi_{handle}\rangle = (0, 0, 1)^T$$

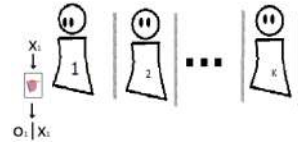
$$\langle a_i | a_{i+1} \rangle = 0$$

Results



- Protocol 1: Markovian process, preparing the post-measured state

Results



- Protocol 1: Markovian process, preparing the post-measured state

Best result by the handle!

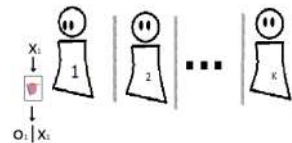


Asymptotic limit given by $\frac{N}{3}$



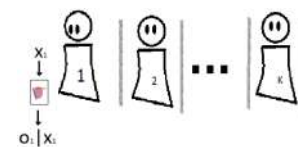
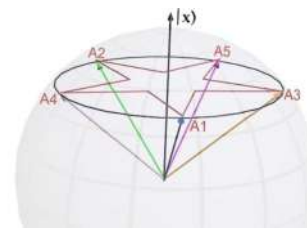
Results

- Measurement protocols 2 and 3
 - These do not determine completely the post-measured State
 - Using the symmetries of the N-cycle quantum realization...



Results

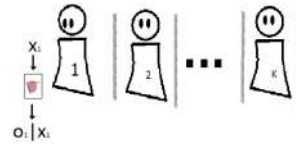
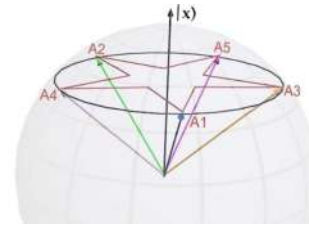
- Measurement protocols 2 and 3



$$\beta = \sum_{i=0}^{N-1} \langle |b_i\rangle \langle b_i| \rangle \geq 1$$

Results

- Measurement protocols 2 and 3



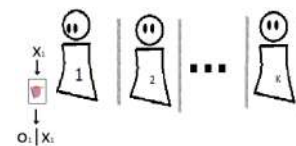
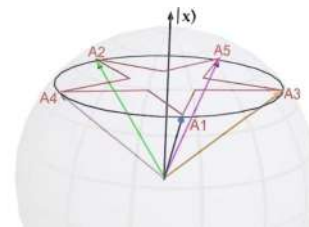
$$\beta = \sum_{i=0}^{N-1} \langle |b_i\rangle \langle b_i| \rangle \geq 1$$

$$\beta_Q^k = B_N \beta_Q^{k-1} + b_N$$

$$B = 1 - \frac{3b}{N}$$

Results

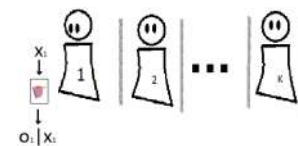
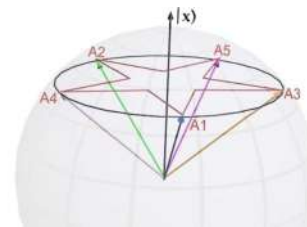
- Measurement protocols 2 and 3



$$\alpha_Q^k = C_N \alpha_Q^{k-1} + c_N \quad \beta_Q^k = B_N \beta_Q^{k-1} + b_N$$

Results

- Measurement protocols 2 and 3



$$\alpha_Q^k = C_N \alpha_Q^{k-1} + c_N \quad \beta_Q^k = B_N \beta_Q^{k-1} + b_N$$

1.

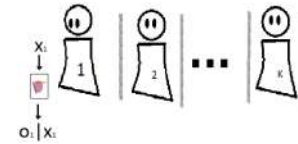
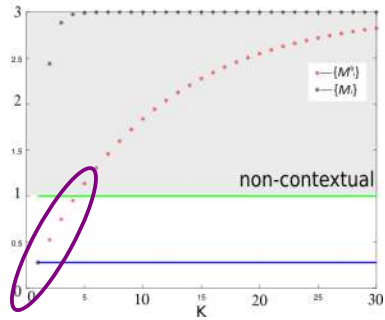
$$|\psi_{handle}\rangle = (0, 0, 1)^T \longrightarrow$$

Gives the best result!

2.

$$\alpha^\infty = \beta^\infty = \frac{N}{3}$$

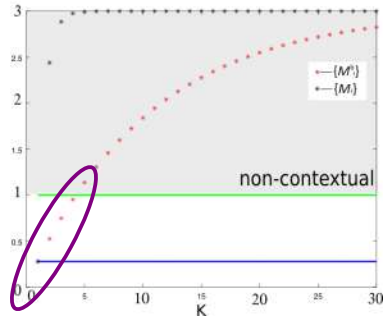
Results



| | K_{\max} - Predef. Order | | |
|-----|----------------------------|---------------|--------------|
| N | $M(\alpha, \beta)$ | $M^a(\alpha)$ | $M^b(\beta)$ |
| 5 | 1 | 1 | 2 |
| 7 | 1 | 1 | 3 |
| 9 | 1 | 1 | 4 |
| 11 | 1 | 1 | 5 |
| 13 | 1 | 1 | 5 |

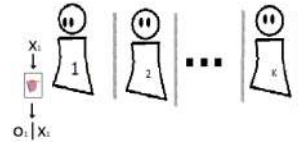
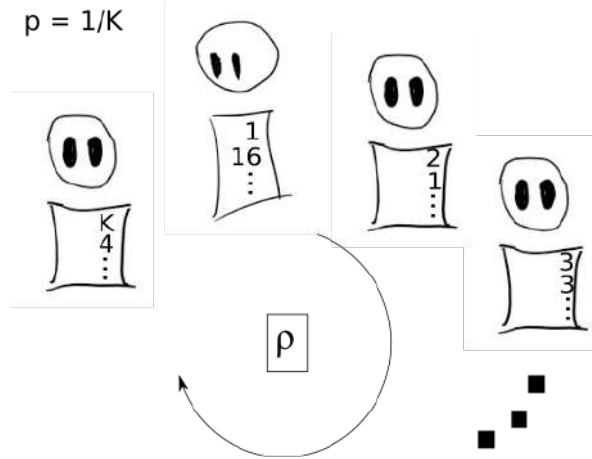
- Asymptotic limit without violation!
- Violation disappears quickly
- Protocols behave differently

Results



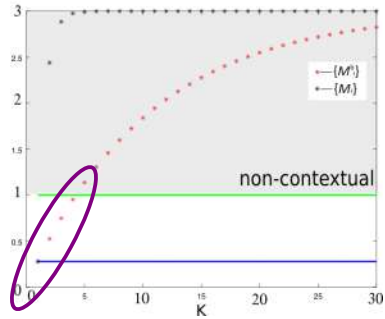
| | K_{\max} - Predef. Order | | |
|-----|----------------------------|---------------|--------------|
| N | $M(\alpha, \beta)$ | $M^a(\alpha)$ | $M^b(\beta)$ |
| 5 | 1 | 1 | 2 |
| 7 | 1 | 1 | 3 |
| 9 | 1 | 1 | 4 |
| 11 | 1 | 1 | 5 |
| 13 | 1 | 1 | 5 |

$$p = 1/K$$

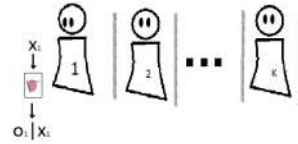
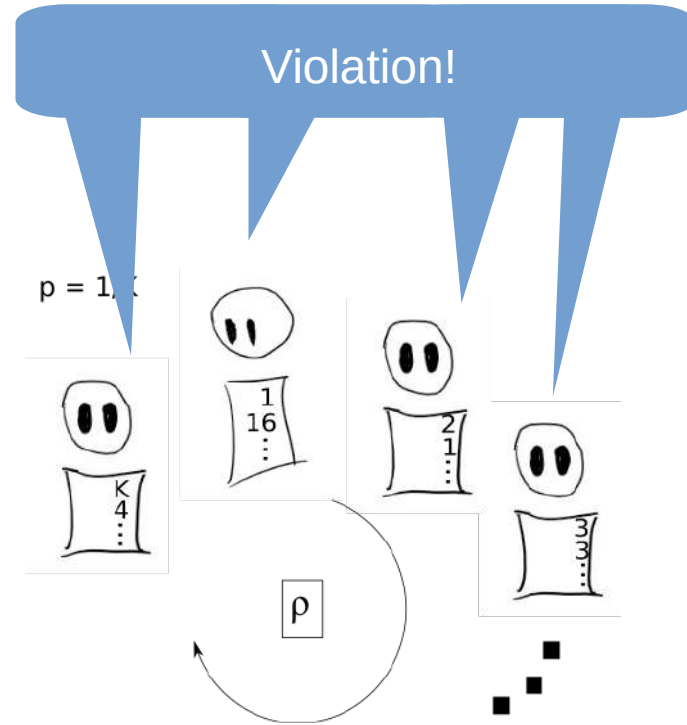


- Asymptotic limit without violation!
- Violation disappears quickly
- Protocols behave differently

Results

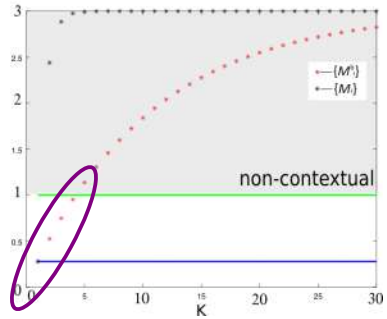


| | K_{\max} - Predef. Order | | |
|-----|----------------------------|---------------|--------------|
| N | $M(\alpha, \beta)$ | $M^a(\alpha)$ | $M^b(\beta)$ |
| 5 | 1 | 1 | 2 |
| 7 | 1 | 1 | 3 |
| 9 | 1 | 1 | 4 |
| 11 | 1 | 1 | 5 |
| 13 | 1 | 1 | 5 |

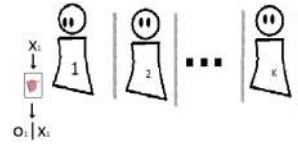
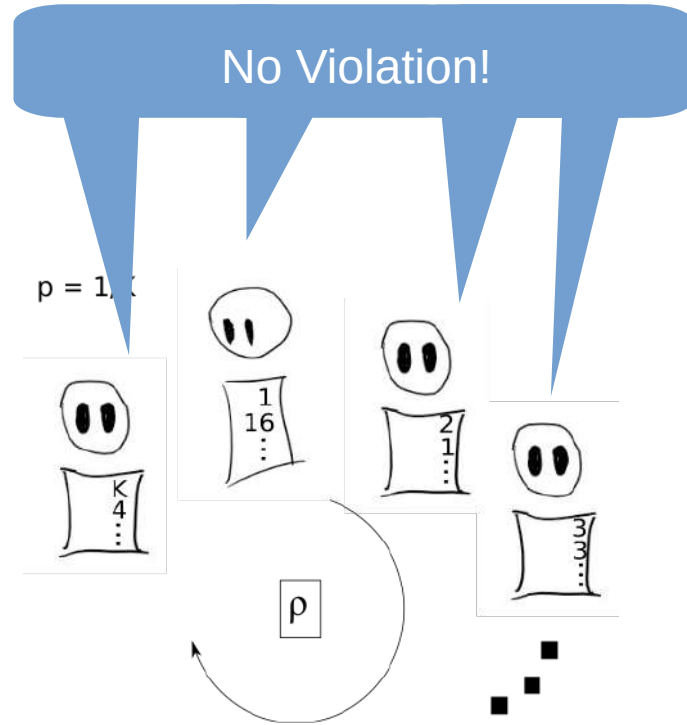


- Asymptotic limit without violation!
- Violation disappears quickly
- Protocols behave differently

Results

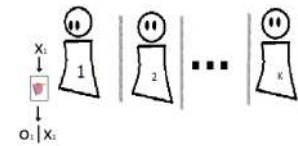
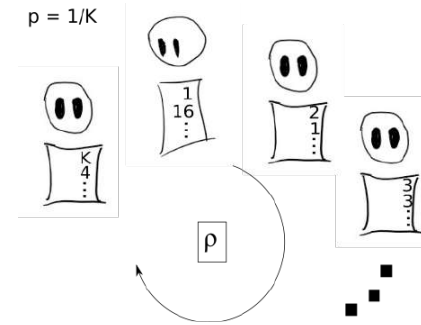
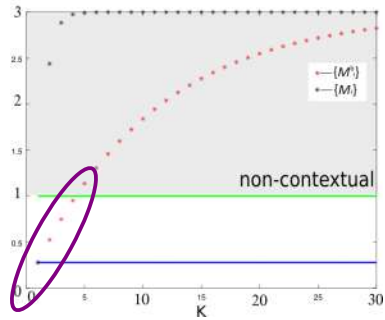


| | K_{\max} - Predef. Order | | |
|-----|----------------------------|---------------|--------------|
| N | $M(\alpha, \beta)$ | $M^a(\alpha)$ | $M^b(\beta)$ |
| 5 | 1 | 1 | 2 |
| 7 | 1 | 1 | 3 |
| 9 | 1 | 1 | 4 |
| 11 | 1 | 1 | 5 |
| 13 | 1 | 1 | 5 |



- Asymptotic limit without violation!
- Violation disappears quickly
- Protocols behave differently

Results

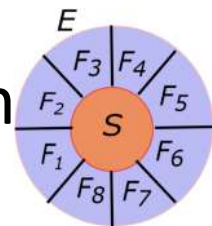
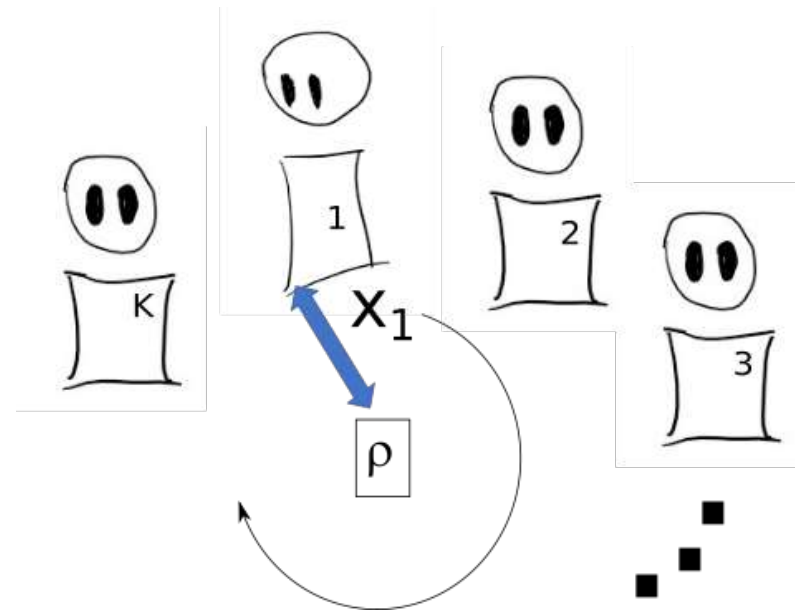
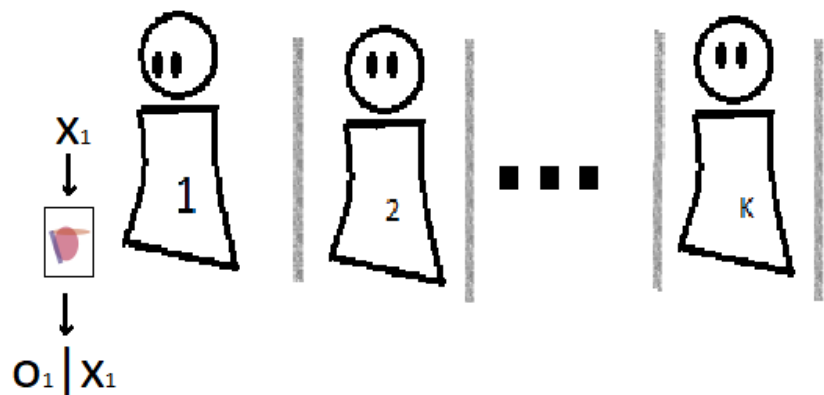


| N | K_{\max} - Predef. Order | | | K_{\max} - Unif. Distr. | | |
|-----|----------------------------|---------------|--------------|---------------------------|---------------|--------------|
| | $M(\alpha, \beta)$ | $M^a(\alpha)$ | $M^b(\beta)$ | $M(\alpha, \beta)$ | $M^a(\alpha)$ | $M^b(\beta)$ |
| 5 | 1 | 1 | 2 | 1 | 2 | 4 |
| 7 | 1 | 1 | 3 | 1 | 1 | 6 |
| 9 | 1 | 1 | 4 | 1 | 1 | 8 |
| 11 | 1 | 1 | 5 | 1 | 1 | 9 |
| 13 | 1 | 1 | 5 | 1 | 1 | 11 |

- Asymptotic limit without violation!
- Violation disappears quickly
- Protocols behave differently

Interpreting as a classical limit

- Multiplayers setup as collisional models and multisystem environment



Spekkens Contextuality

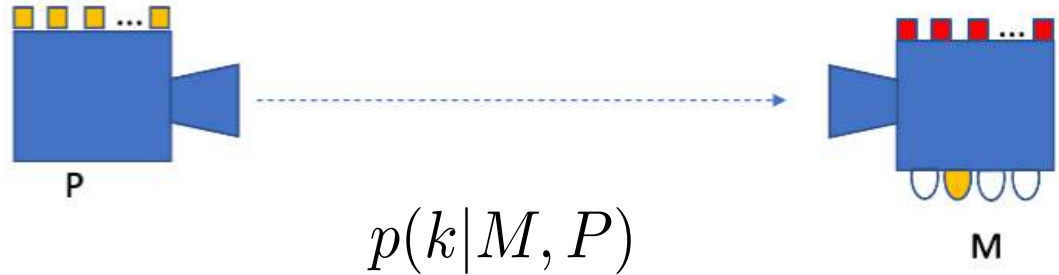
Spekkens Contextuality

Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;

Rule for obtaining behaviours;



Ontological Model

Spekkens Contextuality

Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;

Rule for obtaining behaviours;

Equivalences

$$P_1 \simeq P_2 \iff p(k|M, P_1) = p(k|M, P_2) \\ \forall k, M$$

Ontological Model

Spekkens Contextuality

Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;

Rule for obtaining behaviours;

Equivalences

$$P_1 \simeq P_2 \iff p(k|M, P_1) = p(k|M, P_2) \\ \forall k, M$$

$[P], [k|M]$

Ontological Model

Spekkens Contextuality

Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;

Rule for obtaining behaviours;

Equivalences

Quantum Theory

$$[P] \leftrightarrow [\rho_P]$$

$$[k|M] \leftrightarrow [E_k]$$

Ontological Model

Spekkens Contextuality

Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;

Rule for obtaining behaviours;

Equivalences



Ontological Model

Spekkens Contextuality

Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;

Rule for obtaining behaviours;

Equivalences



Ontological Model

$$P \mapsto \mu_P(\lambda)$$

$$k|M \mapsto \xi_M(k|\lambda)$$

Spekkens Contextuality

Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;

Rule for obtaining behaviours;

Equivalences



Ontological Model

$$P \mapsto \mu_P(\lambda)$$

$$k|M \mapsto \xi_M(k|\lambda)$$

$$\mu_P(\lambda) \geq 0 \quad \forall \lambda \quad \int \mu_P(\lambda) d\lambda = 1$$

$$\xi_M(k, \lambda) \geq 0 \quad \sum_k \xi_M(k|\lambda) = 1 \quad \forall \lambda$$

Spekkens Contextuality

Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;

Rule for obtaining behaviours;

Equivalences



Ontological Model

$$P \mapsto \mu_P(\lambda)$$

$$k|M \mapsto \xi_M(k|\lambda)$$

$$p(k|M, P) = \int \mu_P(\lambda) \xi_M(k|\lambda) d\lambda$$

Spekkens Contextuality

Definition: a context is a label distinction between elements of a given equivalence class

Spekkens Contextuality

Definition: a context is a label distinction between elements of a given equivalence class

Noncontextual ontological model



Operational Theory



Ontological Models



Spekkens Contextuality

Definition: a context is a label distinction between elements of a given equivalence class

Noncontextual ontological model



Operational Theory
 $P_1, P_2 \in [P]$



Ontological Models
 $\mu_{P_1} = \mu_{P_2} = \mu_{[P]}$



Spekkens Contextuality

Definition: a context is a label distinction between elements of a given equivalence class

Noncontextual ontological model



Operational Theory
 $P_1, P_2 \in [P]$



Quantum Theory
 $[\rho], [E_k]$



Ontological Models
 $\mu_{P_1} = \mu_{P_2} = \mu_{[P]}$



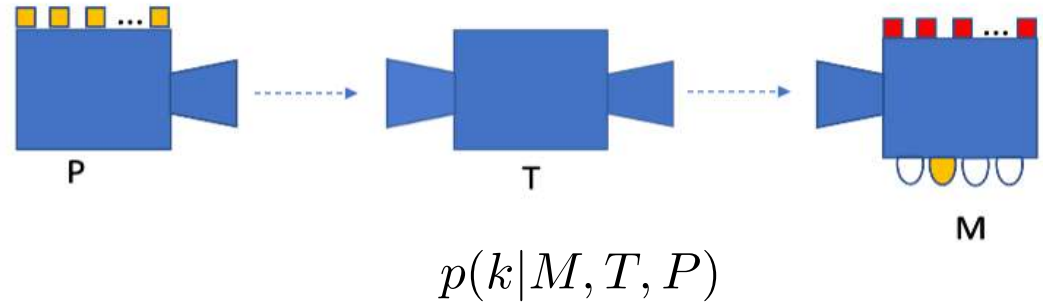
μ_ρ, ξ_{E_k}

Spekkens Contextuality

Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;



Ontological Model

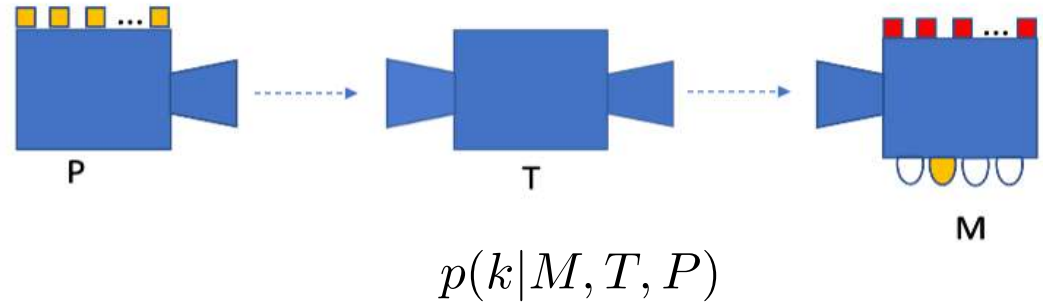
Spekkens Contextuality

Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;

Rule for obtaining behaviours;



Ontological Model

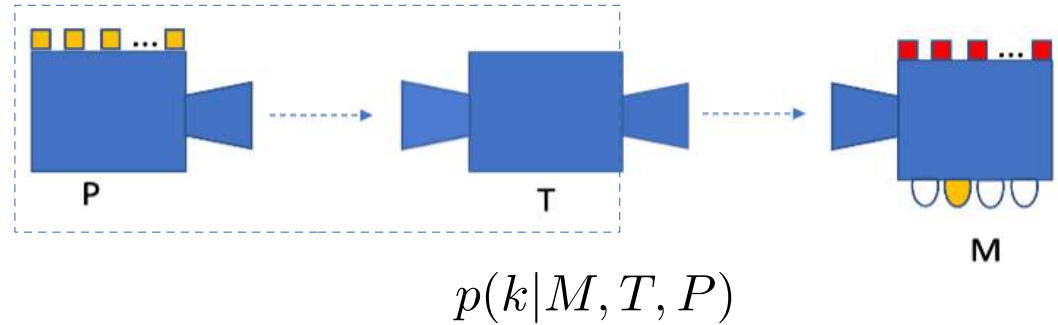
Spekkens Contextuality

Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;

Rule for obtaining behaviours;



Ontological Model

Spekkens Contextuality

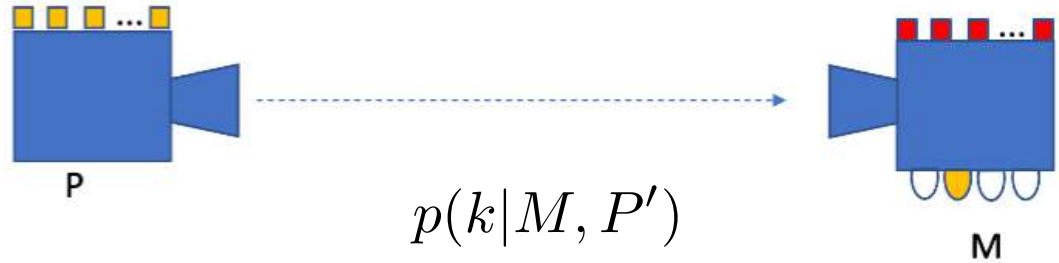
Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;

Rule for obtaining behaviours;

Equivalences



$$P'_1 \simeq P'_2 \iff p(k|M, P'_1) = p(k|M, P'_2) \forall k, M,$$

Ontological Model

Spekkens Contextuality

Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;

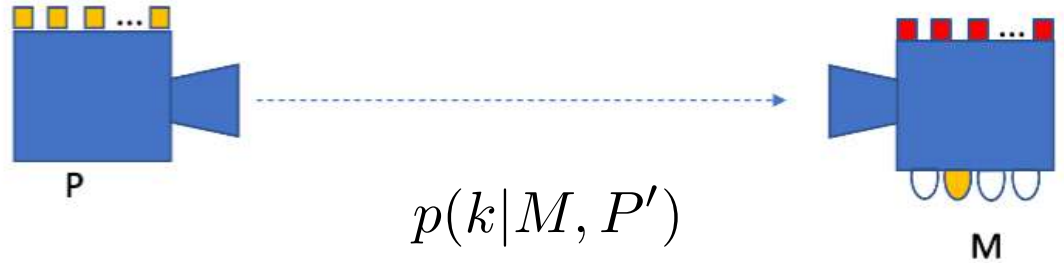
Rule for obtaining behaviours;

Equivalences

$$P'_1 \simeq P'_2 \iff p(k|M, P'_1) = p(k|M, P'_2) \forall k, M,$$

Equivalence classes $[P'], [k|M]$

Ontological Model



Spekkens Contextuality

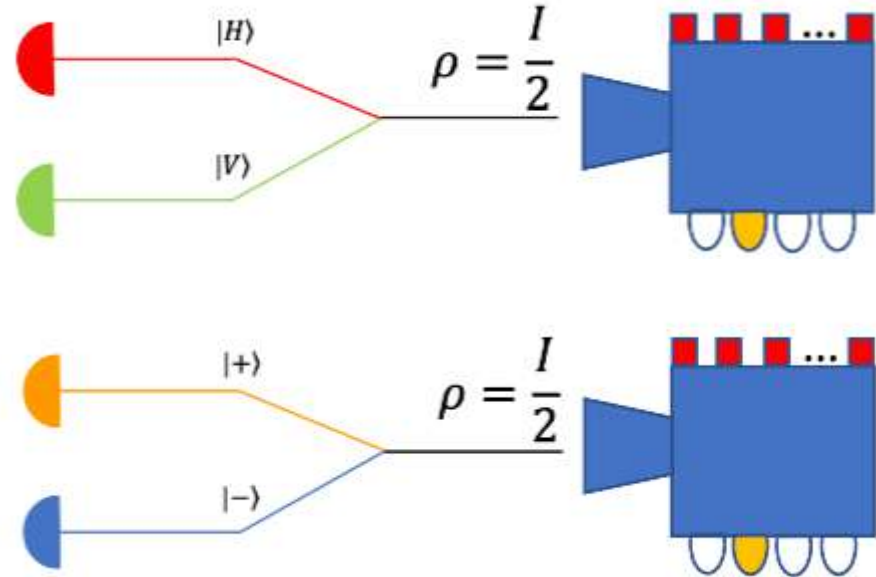
Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;

Rule for obtaining behaviours;
Equivalences

Ontological Model



$$P'_1 \simeq P'_2 \iff p(k|M, P'_1) = p(k|M, P'_2) \forall k, M,$$

Spekkens Contextuality

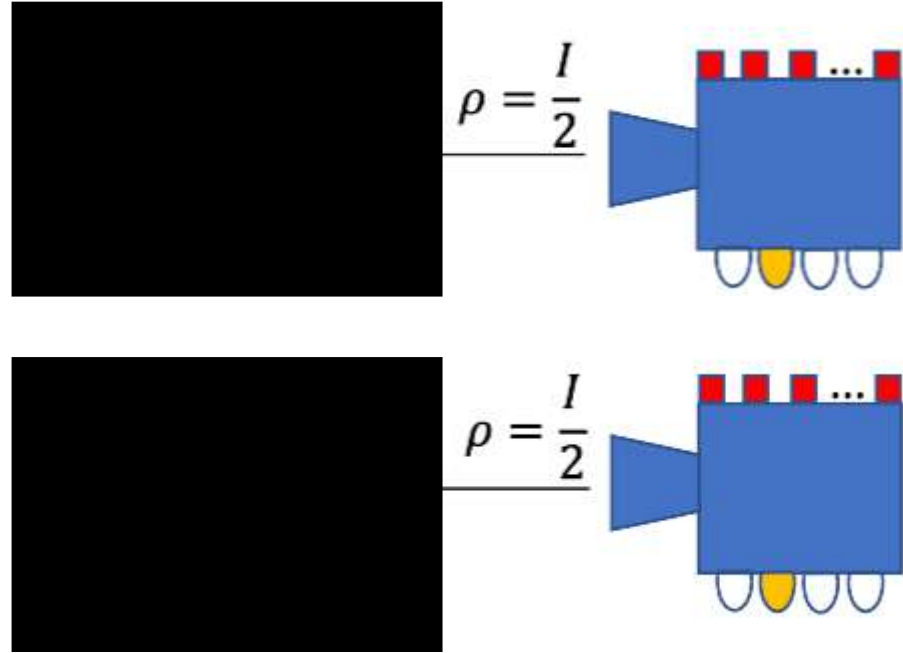
Operational Theory

Procedures

- Preparations,
- Transformations,
- Measurements;

Rule for obtaining behaviours;
Equivalences

Ontological Model



$$P'_1 \simeq P'_2 \iff p(k|M, P'_1) = p(k|M, P'_2) \forall k, M,$$

Appendix 2

BPH

BPH formal

Theorem 5 (Theorem 2 in ref. [15]). *Let $\Phi^{S_t} : \mathcal{D}(\mathcal{H}_A) \rightarrow \mathcal{D}(\bigotimes_{j \in S_t} \mathcal{H}_{B_j})$ be an EWt-dynamics, where $S_t \subset \{1, \dots, N\}$. For every $0 < \delta < 1$ there exists a POVM $\{\tilde{E}_k\}_k$ such that for more than a $(1 - \delta)$ fraction of the subsets S_t ,*

$$\left\| \Phi^{S_t} - \Phi_{obs}^{S_t} \right\|_{\diamond} \leq \left(\frac{27 \ln(2) d_A^6 \log(d_A) t}{N \delta^3} \right)^{\frac{1}{3}}, \quad (\text{A1})$$

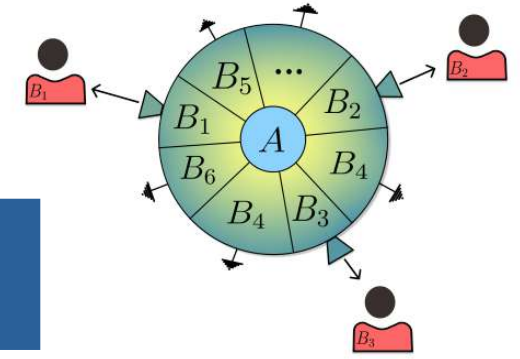
with $d_A \equiv \dim(\mathcal{H}_A)$, and where $\Phi_{obs}^{S_t}$ is a measure-and-prepare map with respect to the family of states $(\sigma_k^{S_t})_k$, meaning that for all $\rho \in \mathcal{D}(\mathcal{H}_A)$,

$$\Phi_{obs}^{S_t}(\rho) = \sum_k \text{Tr}\{\tilde{E}_k \rho\} \sigma_k^{S_t}. \quad (\text{A2})$$

Quantum Darwinism – BPH's approach

*F. Brandão, M. Piani, P. Horodecki, Nat. Comm. 6 7908(2015)

Quantum Darwinism: the environment encodes information regarding system A



Prop.(adapt. from ref.*)

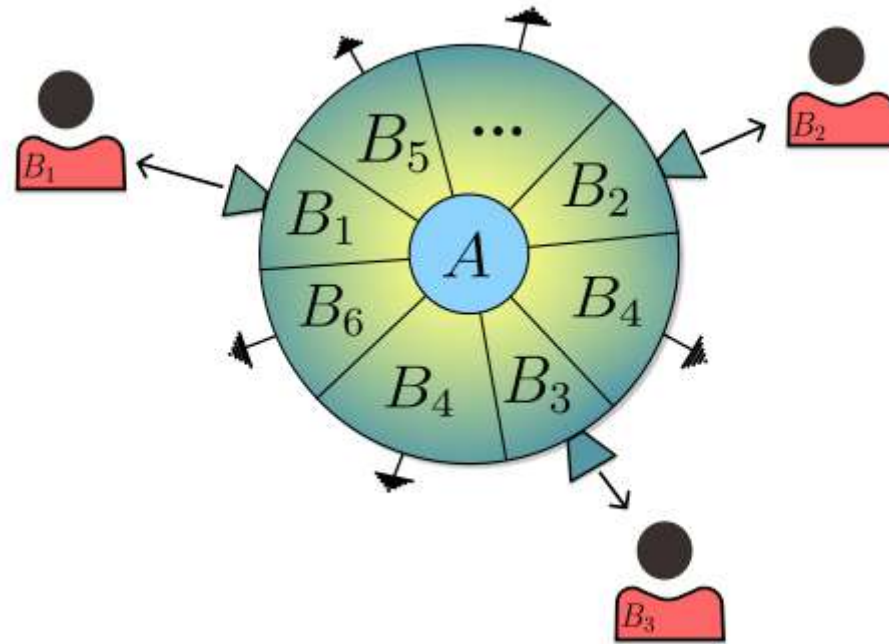
If a QD_η process occurs with high η , Bobs are likely to agree on the outcome they see, i.e.

$$\min_{\rho^A} \sum_k \tilde{p}_k \text{Tr} \left\{ \bigotimes_{j \in S_t} F_k^{B_j} \sigma_k^{B_{S_t}} \right\} \geq 1 - 6t\delta^{\frac{1}{4}},$$

where $\delta = 1 - \eta$

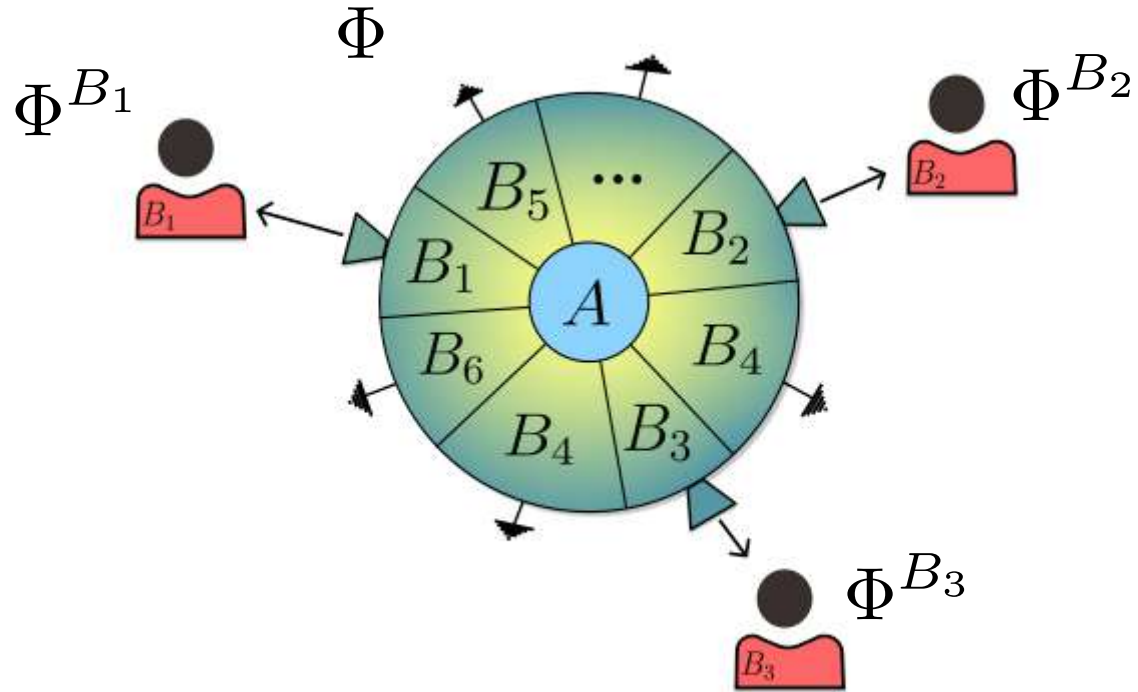
Emergence of noncontextuality in QD

Equivalences in QD settings



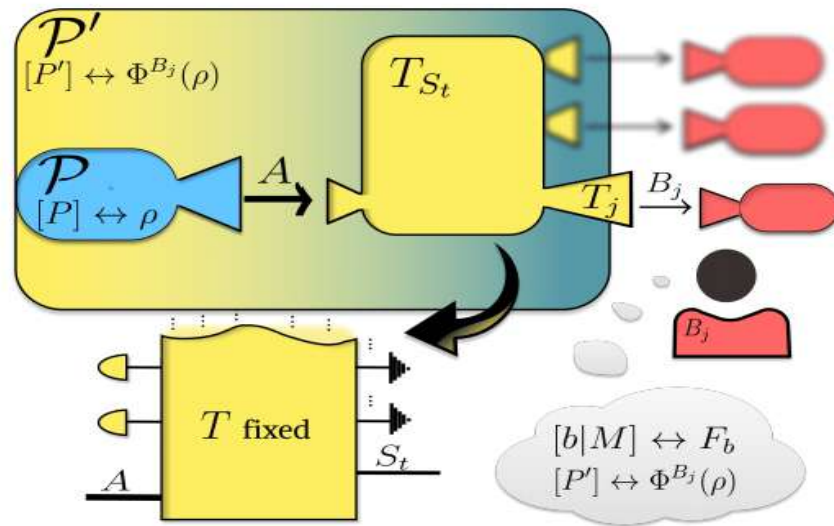
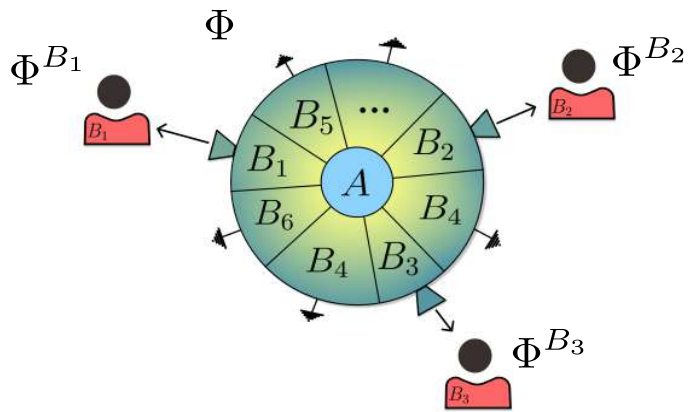
Equivalences in QD settings

Only Bobs' perspectives matter!



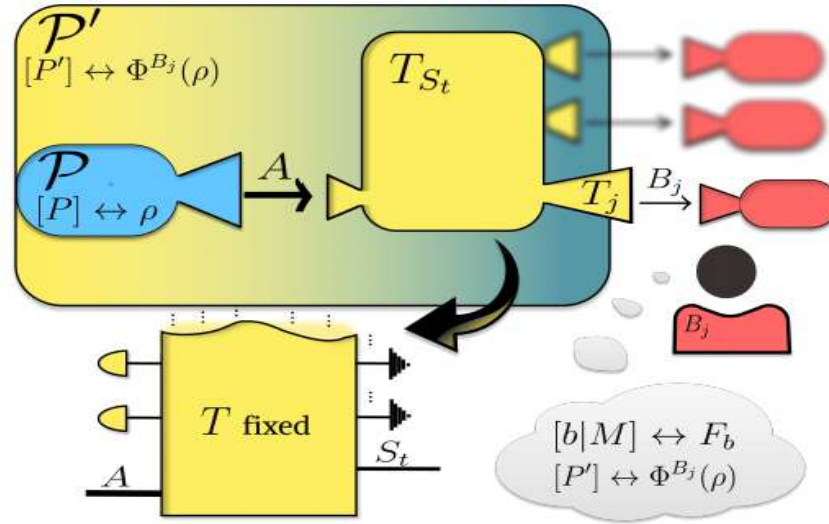
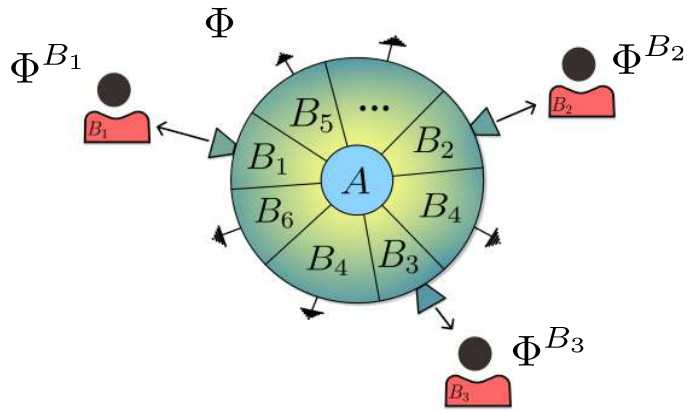
Equivalences in QD settings

Only Bobs' perspectives matter!



Equivalences in QD settings

Only Bobs' perspectives matter!



Equivalence classes for Bob B_j are $[\Phi^{B_j}(\rho^A)]$ and $[F_b]$

Results

Equivalence classes for Bob B_j are $[\Phi^{B_j}(\rho^A)]$ and $[F_b]$

Lemma 1

Consider an EWt-dynamics and the reduced $(\Phi^{B_j})_j$.

For those $B_j \in B_{S_t}$ in which $(\sigma_k^{B_j})_k$ are affinely independent, there exists a noncontextual model.

Results

Equivalence classes for Bob B_j are $[\Phi^{B_j}(\rho^A)]$ and $[F_b]$

Lemma 1

Consider an EWt-dynamics and the reduced $(\Phi^{B_j})_j$.

For those $B_j \in B_{S_t}$ in which $(\sigma_k^{B_j})_k$ are affinely independent, there exists a noncontextual model.

Proof.

$$P'_1 \simeq P'_2 \iff \sigma_{P'_1}^{B_j} = \sigma_{P'_2}^{B_j}$$

Results

Equivalence classes for Bob B_j are $[\Phi^{B_j}(\rho^A)]$ and $[F_b]$

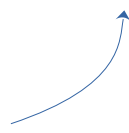
Lemma 1

Consider an EWt-dynamics and the reduced $(\Phi^{B_j})_j$.

For those $B_j \in B_{S_t}$ in which $(\sigma_k^{B_j})_k$ are affinely independent, there exists a noncontextual model.

Proof.

$$\begin{aligned} P'_1 \simeq P'_2 &\iff \sigma_{P'_1}^{B_j} = \sigma_{P'_2}^{B_j} \\ &\iff \text{Tr}\{\tilde{E}_k \rho_{P'_1}^A\} = \text{Tr}\{\tilde{E}_k \rho_{P'_2}^A\} \end{aligned}$$

$$\sigma_P^{B_j} = \sum_k \tilde{p}_k \sigma_k^{B_j}$$


Results

Equivalence classes for Bob B_j are $[\Phi^{B_j}(\rho^A)]$ and $[F_b]$

Lemma 1

Consider an EWt-dynamics and the reduced $(\Phi^{B_j})_j$.

For those $B_j \in B_{S_t}$ in which $(\sigma_k^{B_j})_k$ are affinely independent, there exists a noncontextual model.

Proof.

$$P'_1 \simeq P'_2 \iff \sigma_{P'_1}^{B_j} = \sigma_{P'_2}^{B_j}$$
$$\iff \text{Tr}\{\tilde{E}_k \rho_{P'_1}^A\} = \text{Tr}\{\tilde{E}_k \rho_{P'_2}^A\}$$

$$\Lambda := \{k\}_k;$$

$$\mu_{P'}(k) := \text{Tr}\{\tilde{E}_k \rho_P^A\};$$

$$\xi_M(b|k) := \text{Tr}\{F_b \sigma_k\}.$$

Results

Equivalence classes for Bob B_j are $[\Phi^{B_j}(\rho^A)]$ and $[F_b]$

Lemma 1

Consider an EWt-dynamics and the reduced $(\Phi^{B_j})_j$.

For those $B_j \in B_{S_t}$ in which $(\sigma_k^{B_j})_k$ are affinely independent, there exists a noncontextual model.

Proof.

$$P'_1 \simeq P'_2 \iff \sigma_{P'_1}^{B_j} = \sigma_{P'_2}^{B_j}$$
$$\iff \text{Tr}\{\tilde{E}_k \rho_{P'_1}^A\} = \text{Tr}\{\tilde{E}_k \rho_{P'_2}^A\}$$

$$\Lambda := \{k\}_k;$$

$$\mu_{P'}(k) := \text{Tr}\{\tilde{E}_k \rho_P^A\};$$

$$\xi_M(b|k) := \text{Tr}\{F_b \sigma_k\}.$$

$$p(b|M, P') = \sum_k \mu_{P'}(k) \xi_M(b|k)$$

Results

Equivalence classes for Bob B_j are $[\Phi^{B_j}(\rho^A)]$ and $[F_b]$

Lemma 1

Consider an EWt-dynamics and the reduced $(\Phi^{B_j})_j$.

For those $B_j \in B_{S_t}$ in which $(\sigma_k^{B_j})_k$ are affinely independent, there exists a noncontextual model.

Proof.

$$P'_1 \simeq P'_2 \iff \sigma_{P'_1}^{B_j} = \sigma_{P'_2}^{B_j}$$
$$\iff \text{Tr}\{\tilde{E}_k \rho_{P'_1}^A\} = \text{Tr}\{\tilde{E}_k \rho_{P'_2}^A\}$$

$$\Lambda := \{k\}_k;$$

$$\mu_{P'}(k) := \text{Tr}\{\tilde{E}_k \rho_P^A\};$$

$$\xi_M(b|k) := \text{Tr}\{F_b \sigma_k\}.$$

→ Noncontextual OM!

Results

Proof of Lemma 2

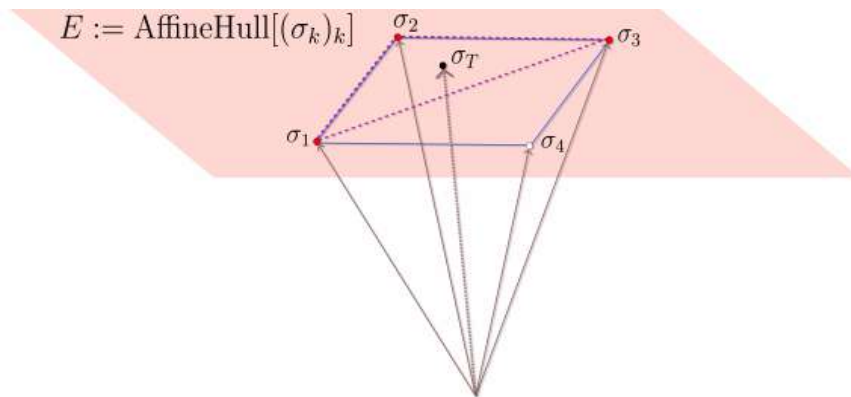
Definition 4 (Distinguishability Bound). *Consider a measure-and-prepare channel defined by $(\{\tilde{E}_k, \sigma_k\})_k$, with $\tilde{E}_k \neq 0$ for all k . Then, there will be states $\rho^A \in \mathcal{D}(\mathcal{H}_A)$ such that $\text{Tr}\{\tilde{E}_k \rho^A\} \neq 0$ for all k . Denote the set of such states by \mathcal{S} . Now, assume (w.l.g.) that $\text{Tr}\{\tilde{E}_1 \rho^A\} \geq \text{Tr}\{\tilde{E}_2 \rho^A\} \geq \dots \geq \text{Tr}\{\tilde{E}_{k_{\max}} \rho^A\} > 0$ (otherwise, relabel (\tilde{E}_k) so that it does). We define the *distinguishability bound* \hat{P} as*

$$\begin{aligned} \hat{P}[(\tilde{E}_k)_k] &:= \min_{\rho^A \in \mathcal{S}} \left[\sum_{i=1}^{k_{\max}-1} \text{Tr}\{\tilde{E}_i \rho^A\} + \frac{\text{Tr}\{\tilde{E}_{k_{\max}} \rho^A\}}{2} \right] \\ &= 1 - \frac{1}{2} \max_{\rho^A \in \mathcal{S}} \text{Tr}\{\tilde{E}_{k_{\max}} \rho^A\}. \end{aligned} \quad (\text{B1})$$

Results

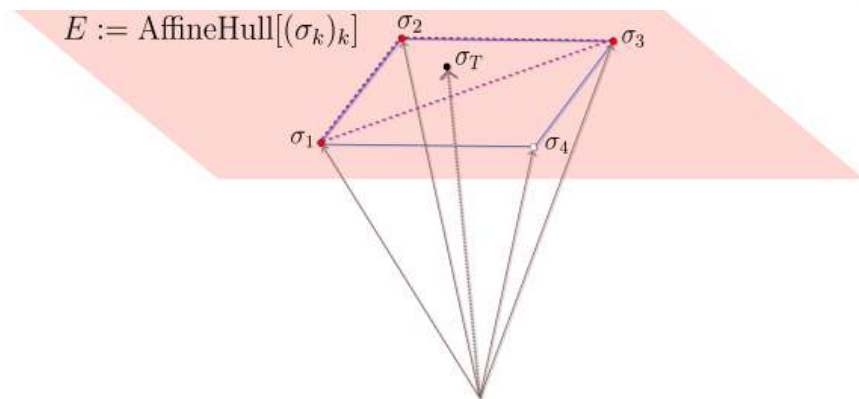
Proof of Lemma 2

Theorem 8 (Carathéodory, adapted from version of ref. [29]). *Given any affine space E of dimension n , for any (non-void) family $f = (\sigma_k)_{k=1}^{k_{\max}}$ in E , the set $\text{ConvHull}(f)$ is equal to the set of convex combinations of families of $n + 1$ points of f .*



Results

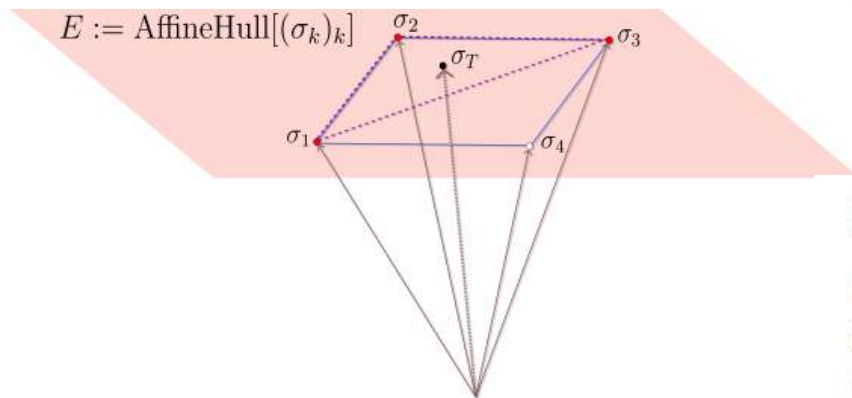
Proof of Lemma 2



Corollary 9. Consider any state of the form $\sigma_T = \sum_k \tilde{p}_k \sigma_k$ with $\tilde{p}_k > 0$ for $k \in \{1, \dots, k_{\max}\}$. Then, if $(\sigma_k)_k$ is an affinely dependent set, there exists a set of convex coefficients $\{q_k\}_k$, with $q_k = 0$ for at least one value of k , such that $\sigma_T = \sum_k q_k \sigma_k$.

Results

Proof of Lemma 2



Corollary 9. Consider any state of the form $\sigma_T = \sum_k \tilde{p}_k \sigma_k$ with $\tilde{p}_k > 0$ for $k \in \{1, \dots, k_{\max}\}$. Then, if $(\sigma_k)_k$ is an affinely dependent set, there exists a set of convex coefficients $\{q_k\}_k$, with $q_k = 0$ for at least one value of k , such that $\sigma_T = \sum_k q_k \sigma_k$.

Remark 1. Assume $(\sigma_k)_k$ to be an affinely dependent set. Consider a state $\sigma_T = \sum_k \tilde{p}_k \sigma_k = \sum_k q_k \sigma_k$ with $\tilde{p}_1 \geq \tilde{p}_2 \geq \dots \geq \tilde{p}_{k_{\max}} > 0$ and $q_k = 0$ for some k . Define the statistical distance between $\{q_k\}_k$ and $\{\tilde{p}_k\}_k$ as

$$D(\{q_k\}) = \frac{1}{2} \sum_k |q_k - \tilde{p}_k|. \quad (\text{B2})$$

Then, $D(\{\tilde{p}_k\}, \{q_k\}) \geq \tilde{p}_{k_{\max}}$.

Results

Proof of Lemma 2

Lemma 10 (Violation of $p_{\text{guess}} \leq \hat{P}$ for all ρ^A implies $(\sigma_k)_k$ is an affinely independent set). Consider an EW dynamics defined by the measure-and-prepare channel $(\tilde{E}_k, \sigma_k)_k$, with $\tilde{E}_k \neq 0$ for all k . Consider the associated distinguishability bound $\hat{P}[(\tilde{E}_k)_k]$ as per definition 4. Then, if for all $\rho^A \in \mathcal{D}(\mathcal{H}_A)$ the inequality

$$p_{\text{guess}}[(\text{Tr}\{\tilde{E}_k \rho^A\}, \sigma_k)_k] > \hat{P}[(\tilde{E}_k)_k] \quad (\text{B3})$$

holds, the states $(\sigma_k)_k$ are affinely independent.

$$\begin{aligned} \text{Tr}\{F_k^* \sigma_k\} &> \frac{1}{2} \quad \forall k \\ p_{\text{guess}}[(\tilde{p}_k, \sigma_k)_k] &= \sum_k \tilde{p}_k \text{Tr}\{F_k^* \sigma_k\} \\ &\leq \sum_{k \neq b} \tilde{p}_k \overbrace{\text{Tr}\{F_k^* \sigma_k\}}^{\leq 1} + \frac{\tilde{p}_b}{2} \\ &\leq \sum_{k \neq b} \tilde{p}_k + \frac{\tilde{p}_b}{2} \leq \sum_{k \neq k_{\max}} \tilde{p}_k + \frac{\tilde{p}_{k_{\max}}}{2} = \hat{P}, \end{aligned}$$

Results

Proof of Lemma 2

Lemma 10 (Violation of $p_{\text{guess}} \leq \hat{P}$ for all ρ^A implies $(\sigma_k)_k$ is an affinely independent set). Consider an EW dynamics defined by the measure-and-prepare channel $(\tilde{E}_k, \sigma_k)_k$, with $\tilde{E}_k \neq 0$ for all k . Consider the associated distinguishability bound $\hat{P}[(\tilde{E}_k)_k]$ as per definition 4. Then, if for all $\rho^A \in \mathcal{D}(\mathcal{H}_A)$ the inequality

$$p_{\text{guess}}[(\text{Tr}\{\tilde{E}_k \rho^A\}, \sigma_k)_k] > \hat{P}[(\tilde{E}_k)_k] \quad (\text{B3})$$

holds, the states $(\sigma_k)_k$ are affinely independent.

$$\begin{aligned} P_{\text{err}}^* &= \sum_k \sum_{b \neq k} \tilde{p}_b \text{Tr}\{F_k^* \sigma_b\} \\ P_{\text{err}} &= \sum_k \overbrace{\text{Tr}\{F_k^* (\sigma_T - \tilde{p}_k \sigma_k)\}}^{\geq 0} \\ &\geq \sum_{k \in K} \text{Tr}\{F_k^* (\sigma_T - \tilde{p}_k \sigma_k)\} \\ &= \sum_{k \in K} \text{Tr}\{F_k^* (\sum_i q_i \sigma_i - \tilde{p}_k \sigma_k)\} \\ &\geq \sum_{k \in K} \text{Tr}\{F_k^* (q_k \sigma_k - \tilde{p}_k \sigma_k)\}. \end{aligned}$$

Results

Proof of Lemma 2

Lemma 10 (Violation of $p_{\text{guess}} \leq \hat{P}$ for all ρ^A implies $(\sigma_k)_k$ is an affinely independent set). Consider an EW dynamics defined by the measure-and-prepare channel $(\tilde{E}_k, \sigma_k)_k$, with $\tilde{E}_k \neq 0$ for all k . Consider the associated distinguishability bound $\hat{P}[(\tilde{E}_k)_k]$ as per definition 4. Then, if for all $\rho^A \in \mathcal{D}(\mathcal{H}_A)$ the inequality

$$p_{\text{guess}}[(\text{Tr}\{\tilde{E}_k \rho^A\}, \sigma_k)_k] > \hat{P}[(\tilde{E}_k)_k] \quad (\text{B3})$$

holds, the states $(\sigma_k)_k$ are affinely independent.

$$\begin{aligned} P_{\text{err}} &\geq \sum_{k \in K} (q_k - \tilde{p}_k) \overbrace{\text{Tr}\{F_k^* \sigma_k\}}^{> \frac{1}{2}} > \frac{1}{2} \sum_{k \in K} (q_k - \tilde{p}_k) \\ &= \frac{1}{2} D(\{\tilde{p}_k\}, \{q_k\}) \geq \frac{p_{k_{\text{max}}}}{2}, \end{aligned} \quad (\text{B9})$$

Results

Proof of relaxation

$$\mathbf{d}(p) := \min_{q \in \mathcal{NC}(\mathfrak{S}_j)} \max_{\substack{M \in \mathcal{M}^{B_j} \\ P' \in \mathcal{P}'}} \sum_b |p(b|M, P') - q(b|M, P')|.$$

$$\mathbf{d}(p) \leq \max_{\substack{M \in \mathcal{M}^{B_j} \\ P' \in \mathcal{P}'}} \sum_b |p(b|M, P') - q_*(b|M, P')|.$$

$$\mathbf{d}(p) \leq \max_{\substack{M \in \mathcal{M}^{B_j} \\ P \in \mathcal{P}}} \sum_b \left| \text{Tr}\{F_b^M \Phi^{B_j}(\rho_P)\} - \text{Tr}\{F_b^M \Phi_{obs}^{B_j}(\rho_P)\} \right|$$

$$= \max_{\substack{M \in \mathcal{M}^{B_j} \\ P \in \mathcal{P}}} \sum_b \left| \text{Tr}\{F_b^M (\Phi^{B_j} - \Phi_{obs}^{B_j})(\rho_P)\} \right|$$

$$= \max_{\substack{M \in \mathcal{M}^{B_j} \\ P \in \mathcal{P}}} \sum_b \left| \text{Tr}\{F_b^M \mathcal{R}_j(\rho_P)\} \right|.$$

Results

Proof of relaxation

Lemma 6 (Hölder's inequality). *Let A, B be any $n \times n$ complex matrices. Then,*

$$|\mathrm{Tr} A^\dagger B| \leq (\mathrm{Tr} |A|^l)^{\frac{1}{l}} (\mathrm{Tr} |B|^s)^{\frac{1}{s}} \quad (\text{A4})$$

such that $1 \leq l, s \leq \infty$ with $\frac{1}{l} + \frac{1}{s} = 1$.

$$|\mathrm{Tr}\{F_b^M \mathcal{R}_j(\rho_P)\}| \leq \sqrt{\mathrm{Tr}\{|\mathcal{R}_j(\rho_P)|^2\}} \sqrt{\mathrm{Tr}\{|F_b^M|^2\}}$$

$$\begin{aligned} \mathbf{d}(p) &\leq \max_{\substack{M \in \mathcal{M}^{B_j} \\ P \in \mathcal{P}}} \sum_b |\mathrm{Tr}\{F_b^M \mathcal{R}_j(\rho_P)\}| \\ &\leq \max_{\substack{M \in \mathcal{M}^{B_j} \\ P \in \mathcal{P}}} \sum_b \sqrt{\mathrm{Tr}\{|\mathcal{R}_j(\rho_P)|^2\}} \sqrt{\mathrm{Tr}\{|F_b^M|^2\}} \\ &= \max_{M \in \mathcal{M}^{B_j}} \sum_b \sqrt{\mathrm{Tr}\{|F_b^M|^2\}} \max_{P \in \mathcal{P}} \sqrt{\mathrm{Tr}\{|\mathcal{R}_j(\rho_P)|^2\}} \\ &= C \max_{P \in \mathcal{P}} \sqrt{\mathrm{Tr}\{|\mathcal{R}_j(\rho_P)|^2\}} \end{aligned}$$

Results

Proof of relaxation

After relating the 2-norm to the 1-norm,...

$$\begin{aligned} \mathbf{d}(p) &\leq \frac{C}{d_A} \max_{\mathbf{1}_A \otimes \rho \in \mathcal{D}(\mathcal{H}_A^{\otimes 2})} \left\| \text{id}_A \otimes \left(\Phi^{B_j} - \Phi_{obs}^{B_j} \right) (\mathbf{1}_A \otimes \rho) \right\|_1 \\ &\leq \frac{C}{d_A} \max_{\sigma \in \mathcal{D}(\mathcal{H}_A^{\otimes 2})} \left\| \text{id}_A \otimes \left(\Phi^{B_j} - \Phi_{obs}^{B_j} \right) (\sigma) \right\|_1 \\ &= \frac{C}{d_A} \left\| \Phi^{B_j} - \Phi_{obs}^{B_j} \right\|_{\diamond}. \end{aligned}$$