

Causal contextuality and adaptive MBQC

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(joint work with Cihan Okay)

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Joint work with Cihan Okay



Bilkent University



Funded by
the European Union



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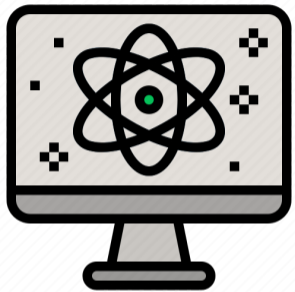
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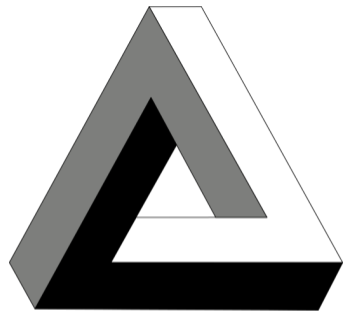
- ▶ Related to talks by Samson & Amy, but only using a particular type of models.
- ▶ May have some relation to upcoming talk by Sivert.



Introduction



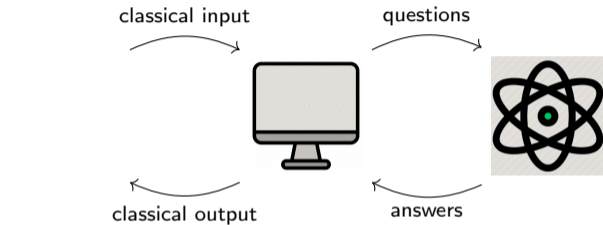
Quantum advantage



Contextuality / Nonclassicality

Contextuality in MBQC

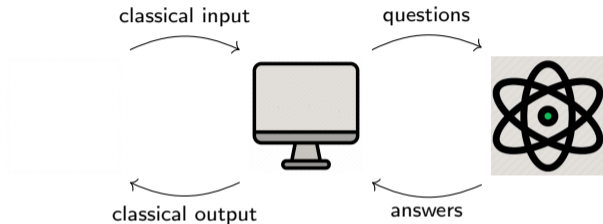
'Contextuality in measurement-based quantum computation', Raussendorf, PRA 2013.



MBQC: Classical control computer with access to quantum resources

Contextuality in MBQC

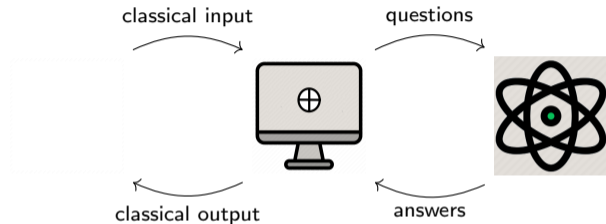
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l_2 -MBQC: Classical control computer with access to quantum resources

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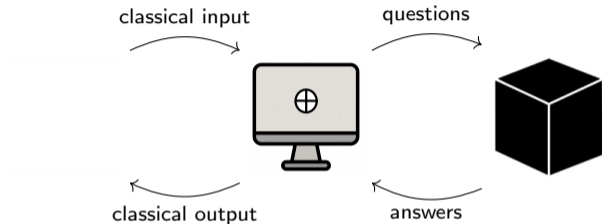


l_2 -MBQC: Classical control computer with access to quantum resources

- ▶ Classical control restricted to \mathbb{Z}_2 -linear computation

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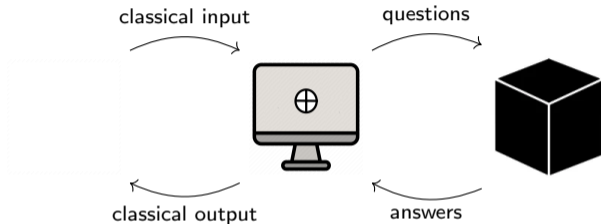


ℓ_2 -MBQC: Classical control computer with access to quantum resources

- ▶ Classical control restricted to \mathbb{Z}_2 -**linear** computation
- ▶ Resource treated as a **black box**, described by its **behaviour**

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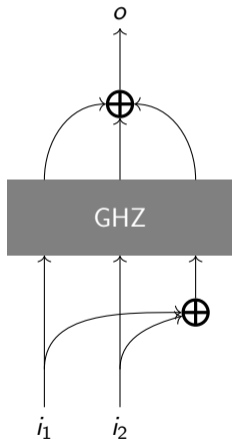
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Theorem

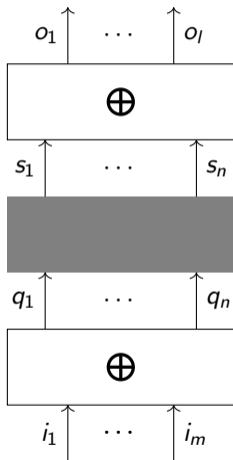
If an ℓ_2 -MBQC **deterministically** computes a **nonlinear** Boolean function then the resource is **strongly contextual**.

The AND function

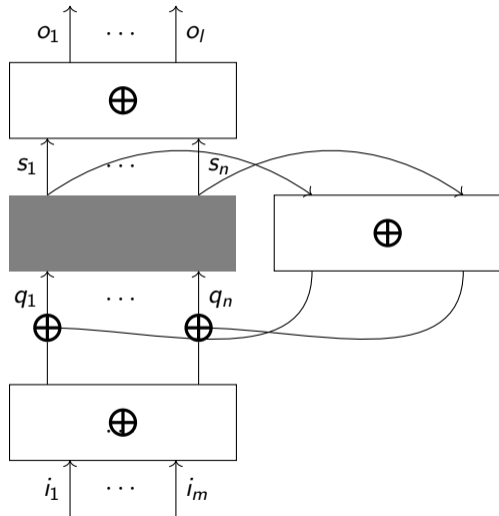
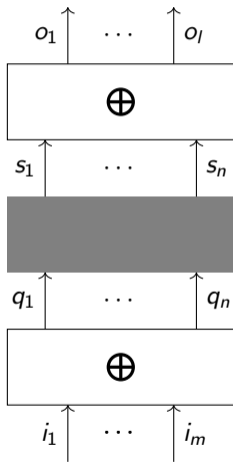
'Computational power of correlations', Anders & Browne, PRL 2009.



Adaptive MBQC



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- ▶ For a given computation, the black box is used in a given (partial) order.

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Question

In adaptive MBQC:

- ▶ For a given computation, the black box is used in a given (partial) order.
- ▶ Why should the classical benchmark be so restrictive?
- ▶ We could think of a classical model that exploits this (causal) knowledge.

Can we find conditions on the computed functions that exclude even such classical HV models?

Non-locality

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Deterministic local models

A **deterministic local** model is given by a family of functions

$$f_\omega : \mathcal{Q}_\omega \longrightarrow \mathcal{A}_\omega \quad (\omega \in \Omega).$$

E.g. bipartite scenario: $(\mathcal{Q}_A \longrightarrow \mathcal{A}_A) \times (\mathcal{Q}_B \longrightarrow \mathcal{A}_B)$.

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Locality and no-signalling

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Causal contextuality

Causal scenarios

'*The sheaf-theoretic structure of definite causality*', Gogioso & Pinzani, QPL 2021.



- ▶ A causal (partial) order between sites

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- ▶ A causal (partial) order between sites
- ▶ Classical models are allowed to use information from the causal past
- ▶ i.e. the answer at a given site may depend on the questions asked at sites in its past.
- ▶ Correspondingly, no-signalling gets relaxed, permitting signalling to the future.

NB: a special class of scenarios within the formalism presented by Samson & Amy.

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Notation: $\downarrow \omega := \{\omega' \in \Omega \mid \omega' \leq \omega\}$ $\downarrow S := \bigcup_{\omega \in S} \downarrow \omega = \{\omega' \in \Omega \mid \exists \omega \in S. \omega' \leq \omega\}$

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A **deterministic causally classical** model is given by a family of functions

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E.g. bipartite scenario with $A \leq B$: $(\mathcal{Q}_A \longrightarrow \mathcal{A}_A) \times (\mathcal{Q}_A \times \mathcal{Q}_B \longrightarrow \mathcal{A}_B)$.

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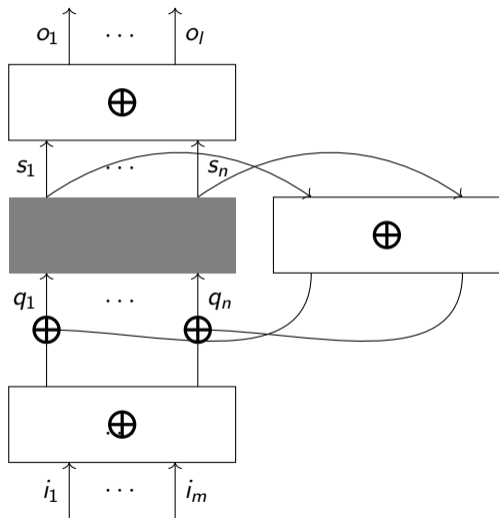
This yields models that are **no-signalling except from the past**.

$f : \mathcal{Q}_A \times \mathcal{Q}_B \longrightarrow D(\mathcal{A}_A \times \mathcal{A}_B)$ such that $P_f(a_A | q_A, q_B) = P_f(a_A | q_A)$ **but not for a_B** .

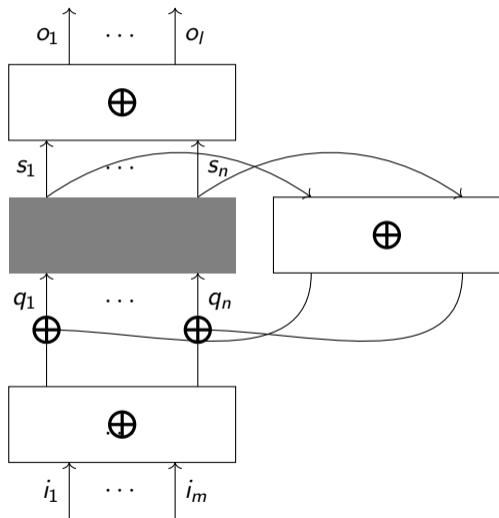
Measurement-based quantum computation

Adaptive ℓ_2 -MBQC

- ▶ input size m
- ▶ output size l
- ▶ adaptive structure (Ω, \leq) with $n = |\Omega|$



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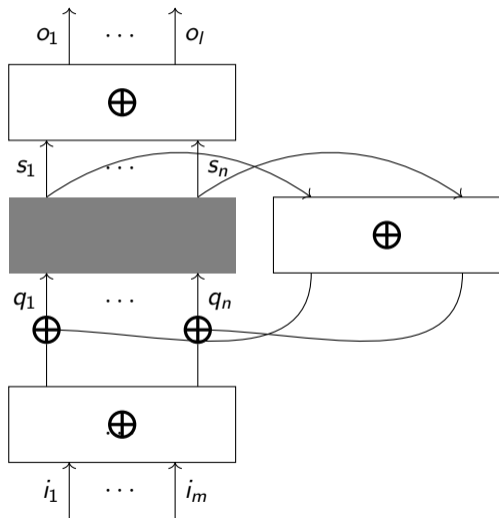
- ▶ $Q : \mathbb{Z}_2^m \rightarrow \mathbb{Z}_2^n$

- ▶ $T : \mathbb{Z}_2^n \rightarrow \mathbb{Z}_2^n$

- ▶ $Z : \mathbb{Z}_2^n \rightarrow \mathbb{Z}_2^l$

such that $T_{\omega, \omega'} \neq 0 \Rightarrow \omega \leq \omega'$

Adaptive ℓ_2 -MBQC



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such that $T_{\omega, \omega'} \neq 0 \Rightarrow \omega \leq \omega'$

$$\mathbf{q} = Q\mathbf{i} + T\mathbf{s}$$

$$\mathbf{s} \leftarrow e(\mathbf{q})$$

$$\mathbf{o} = Z\mathbf{s}$$

implements a function $\mathbb{Z}_2^m \rightarrow D(\mathbb{Z}_2^l)$.

Causal contextuality and adaptive MBQC

Main result

- ▶ Functions $g : \mathbb{Z}_2^m \rightarrow \mathbb{Z}_2$ can be represented as m -variable polynomials in \mathbb{Z}_2 , $\pi(g)$.
- ▶ Functions $g : \mathbb{Z}_2^m \rightarrow \mathbb{Z}_2^l$ are represented by l -tuples of m -variable polynomials $\pi(g) = \langle \pi(g)_1, \dots, \pi(g)_l \rangle$.

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Theorem

Let (e, Q, T, Z) be an Ω -adaptive ℓ_2 -MBQC protocol that **deterministically** computes a function $g : \mathbb{Z}_2^m \rightarrow \mathbb{Z}_2^l$. If e is **causally classical** then each $\pi(g)_j$ is a polynomial with degree **at most the height of Ω** , where the height of a poset is the maximum length of a chain in it.

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NB: If Ω is flat, i.e. has height 1, one recovers Raussendorf's result about nonlinear functions.

Questions...

?